A Business User Model-Driven Engineering Method for Developing Information Systems

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22 February 2015

This Thesis is submitted for the degree of Doctor of Philosophy

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Abstract

With the rapid development of general-purpose programming languages and platform technologies, software engineers have faced more various challenges in software development to those that occurred in the past decades. Requirements Change might cause several project management and technical conflicts associated with requirement elicitation, intercommunication and later changes of specifications.

In the real-world, there is a demand to adopt an accurate information system that satisfies the requirements and is used effectively for the business. However, having vague or misinterpreted requirements causes errors and extra costs. Therefore, domain experts, who clearly understand the business logic and goals, and are aware of what exactly they need inside an organisation without professional software developing skills, should play key roles in the development lifecycle using high level tools. Model-Driven Engineering, is a software engineering discipline that aims at raising the level of abstraction by capturing system specifications through the employment of Models. MDE supports integration and interoperability, improves software quality, and reduces development cost by supporting the automatic creation of software systems using appropriate toolsets.

This thesis is all about raising the level of abstraction at which information systems are built, using business end-users knowledge and MDE to achieve the result. The work introduces, first, Micro-Modelling Language (μML) , a lightweight modelling language that is used to express basic structural and behavioural aspects of information systems using effectively business-users knowledge of their desired system. Throughout the work, graphical notation and semantics for the language concepts are identified, providing a simpler and semantically cleaned modelling language than standard UML and other UML-based languages.

The work also proposes BUILD (Business-User Information-Led Development), an End-User MDE approach to support the construction of information systems using high-level specifications and accelerate the development process using layered model transformation and code generation. Throughout the thesis, a number of development phases and model transformation steps are identified to allow the low-level technical detail be introduced and developed automatically by rules, with less end-users engagement. Domain-Specific code generators, for generating executable Java Swing Applications code and MySQL script, are used to demonstrate the validity of the research.

Acknowledgements

I would like to use this opportunity to express my gratitude and love to all the people in my life, who were there for me and were a part of the process during which this thesis would not have happened without the generous support and encouragement of many people.

The first person I would like to thank is my supervisor Dr. Anthony J.H. Simons, as this PhD research is the synthesis of five years of work whereby I have been accompanied and guided by him. During these years I have known Anthony (Tony) as both a clever researcher, and a genuine and friendly person. Together with him, I would like to thank Prof. Georg Struth for his valuable advice, ideas and suggestions during the design of the formal logic, which surely improved the quality of the thesis. Moreover, many thanks to Dr. Simon Foster for his worthy assistance in the theoritical side of Graph Transformation at the early stage of the research.

I would like to mention all members of the Verification and Testing (vt) group. To those I have spent most of these years together with in the Department of Computer Science, University of Sheffield. Many thanks to all of them. Moreover, huge thanks go to my dear wife and all friends for putting up with me and for their support and patience that were fundamental in concluding my PhD.

I feel a deep sense of gratitude and I owe sincere and earnest thankfulness to my valuable Irish friends for their special caring, Dr. Hayder Ahmed, his lovely family, Dr. A Almanea, and Mrs L Ahmed, You are a part of my life who were here for me when I needed any piece of advice and encouragement. My special thanks and love to the special one who was never-ending support and inspiration in every conceivable way. Without the patience, love and endless support this thesis wouldn't have been written.

List of Acronyms

List of Publications

- Ahmad F. Subahi, Anthony .J.H. Simons. A Multi-level Transformation from Conceptual Data Models to Database Scripts Using Java Agents, Proceedings in the Workshop on Composition and Evaluation of Model Transformations, King's College London, 2011.
- Ahmad F. Subahi and Anthony J.H. Simons, A model transformation approach for translating conceptual database schemas into executable database systems, Technical Report, Department of Computer Science, University of Sheffield, 2011, 1-10.
- Ahmad F. Subahi and Anthony J.H. Simons, Domain-Specific Language for Enabling End-Users Model-Driven Information System Engineering, World Academy of Science, Engineering and Technology, International Science Index 79, International Journal of Computer, Information, Systems and Control Engineering, 7(7), 2013, 318 - 321.
- Anthony J.H. Simons, Ahmad F. Subahi and Stephen M.T. Eyre, Practical model-tocode transformation in four object-oriented programming languages, Technical Report, Department of Computer Science, University of Sheffield, 2011, 1-25.

Declaration

I hereby declare that this thesis is my own work and effort and that it has not been submitted anywhere for any award. Where other sources of information have been used, they have been acknowledged.

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1 Introduction

"Walking on water and developing software from a specification are easy if both are frozen" Edward V. Berard

1.1 Context

The thesis is about raising the level of abstraction at which information systems are built, and this requires a user-friendly modelling language and an appropriate Model-Driven Engineering (MDE) method to achieve the result. In this chapter an overview of the general area addressed by this thesis is presented, highlighting certain aspects such as the principles of MDE and some comparisons of different MDE concepts and techniques. The outline of the thesis is presented next, in the rest of the chapter, including the motivation, aims and a summary of contributions and objectives.

1.2 Grand Challenges in Software Development

With the rapid development of general-purpose programming languages and platform technologies, software engineers have in the current decade faced more varied challenges in software development than those encountered in earlier decades. According to [23], these challenges are a result of three major factors, which lead to the accidental increase in lines of code: Implementation Complexity, Platform Diversity, and Requirements Change. For instance, building a robust software system that is feasible, distributed, secure, and portable requires additional sophisticated requirements and more stakeholders to be involved in the development process. This might cause several project management and technical conflicts associated with requirement elicitation, inter-communication and later changes of specifications. Consequently, there will be a demand to expand the project size, which leads to an increase in the development and maintenance cost, extra testing and bug-fixing tasks, and a delay in time of market delivery [62].

Besides this, dealing with the variety of existing platforms brings other issues to developers and their knowledge. It is unattainable for a developer to be a specialist of every programming language, and domain technology. For example, there is no standard way, adopted in all programming languages and platform technologies, to deal with a particular requirement. Each domain has an appropriate approach to handle some types of requirements, which cannot be applied to other domains. Thus, developers must be capable of coping with a variety of domains, implementation technologies and languages to produce good-quality systems. Furthermore, as an on-going process, the clients might change the requirements during the development lifecycle. The major changes might require code rewriting, and disturbing other processes, which increases cost [23].

1.3 Recent Possible Solutions for the Problems

A number of software development approaches, such as Model-Driven Engineering (MDE), and Language-Driven Engineering (LDE) [23], have introduced the idea of raising the level of abstraction beyond the General-Purpose Programming Languages (GPLs) through the usage of Models and high-level programming languages tailored to a specific domain, known as Domain-Specific Languages (DSLs). These various strategies are commonly adopted in software development, as Domain-Specific Modelling Frameworks (DSMs) [62], to tackle the challenges stated above. The DSM approaches are based on the usage of DSL with an associated compiler to model and capture system specifications of a particular domain at a higher level of abstraction [62]. They bring a remarkable contribution for reducing software complexity, achieving the separation of concerns between the design and implementation phases, and providing a better view of the system. This enables developers to specify their system and capture its requirements and design problems independently from any platform-specific implementation, associated with a particular technology. Then the mappings from design to code are defined in order to derive the executable software automatically from high-level specifications. As a result, the overall software quality and productivity will be increased [23].

1.4 Model-Driven Engineering (MDE)

In MDE, a model is considered a primary entity in the development lifecycle that provides an abstract representation of a system. It supports integration and interoperability, improves software quality, and reduces a development cost [6]. The general strategy of MDE aims to capture system specifications through the employment of models that are expressed at a very high level of abstraction. The models can be defined using a metamodel or an appropriate DSL for a particular domain [62]. In Model-Driven approaches, the metamodel, a model of a model, defines the abstract syntax and the relationship between concepts of the model. Figure 1.1 demonstrates the relationship between model and metamodels.

The Meta-Object Facility (MOF) [87] is a metamodelling strategy that is provided by the Object Management Group (OMG) [86] and supported in Model Driven Architecture (MDA). It has a hierarchical structure that consists of four layers from M0 to M3 (Figure 1.2), each one defines the level below to ensure the consistency and the correctness of the instance model syntax and semantics at each level of abstraction [22].

Figure 1.1: The Relationship between Models and Metamodels

Figure 1.2: MOF Levels

- M3 is the meta-meta-model layer and defines models (metamodels) in the layer below it, the M2 layer. For example UML [70] node and arc diagrams that provide meta language to define UML modelling language [70].
- M2 is the metamodel layer and is a language that describes the models in the M1 layer. For example, UML [70] elements, namely, Class, Association, and Attribute are defined at M2, as illustrated in a snapshot of the UML [70] metamodel (Figure 1.3).

Figure 1.3: An example of M2 level models (metamodels)

• M1 is the model layer. This is where instances of UML [70] diagrams are categorised and is a level where user designed elements are defined. It contains the application model, such as the classes of an object-oriented system, or the table definitions of a relational database. Figure 1.4 illustrates an example of models at M1 level.

Figure 1.4: An example of M1 model

• M0 is the runtime layer that contains data about the objects that have been created from the definitions in the model to represent an instance of the system at run time. Figure 1.5 shows an example of M0 models of the relationship, called lives, between Person and Address classes.

Figure 1.5: An Example of models at M0 level

By using models and metamodels, the automation can be fulfilled via the utilisation of model transformation approaches or tools. In MDE for example, a series of model transformations, that are defined by rules, is applied to an abstract model conforming to a metamodel to form a more concrete target model conforming to the same or different metamodel. The model transformations can be categorised into two types: horizontal and vertical transformations. A horizontal one resides at the same abstraction level in which it might present different views of the system or performs a refactoring process [62]. It might be a model in-place optimisation or Refactoring as in Figure 1.6 (A, E) . On the other hand, a vertical transformation, Figure 1.6 (B, C, D) not only presents a distinct view of the system, but also it moves models between different abstraction levels [62]. It might be a *refinement* or *abstraction* (Figure 1.6 (C)), *code generation* or *visualisation* (Figure 1.6 (B) , and (D)) respectively.

The model transformations, shown in Figure 1.6, can also be classified, based on the direction, into two types, namely, Unidirectional, and Bidirectional transformation. Forward, and Reverse Engineering, Figure 1.6 (B) and (D) respectively, are examples of the Unidirectional transformations. Whereas Figure 1.6 (C) illustrates a Bidirectional transformation.

Figure 1.6: Types of model transformations

1.5 Motivation

When developing an information system to satisfy a business need, the expectation is that such a system will accurately satisfy its well-specified requirements. However, having vague or misinterpreted requirements causes errors and extra costs. Therefore, domain experts, who clearly understand the business logic and goals, and are aware of what exactly they need inside an organisation without professional software developing skills, should play key roles in the development lifecycle using high level tools. There are many approaches that aim to tackle these issues and reduce the gap between initial requirements and implementation and accelerating the development process, such as Model-Driven Engineering (MDE), Domain-Specific Languages (DSL) and End-user Development (EUD).

Although the Unified Modelling Language, UML [70], is commonly used to express the structure and behaviour of a system within MDE approaches, it suffers from semantic ambiguity and complexity issues. While it is possible to create designs in UML 2.x, from which complete code can be generated, this can only be done at the expense of supplying complete low-level detail in the models; that is, UML cannot at the same time be both abstract and sufficient for code generation. These fully-detailed UML models must necessarily contain far more detailed technical information than a typical end-user would be expected to understand; indeed if they have to create such models, there is a risk that these will be inconsistent, incomplete or simply difficult to understand.

To sum up, we find that current MDE tools require a degree of technical skills when dealing with complex models that are syntactically and semantically unclear. This prevents end-users from expressing directly their functional requirements. Researchers focus on developing appropriate modelling languages and DSLs to accelerate and enhance the development process led by skilful developers. As we will see in the literature review (Chapter 2), there is less attention on the role of business end-users during the MDE development lifecycle, than we would hope to see.. Even though some MDE approaches allow the end-user, at some point of development, to work side by side with developers just for customising system artefacts, it is worth exploring and improving the role of business users by allowing them to lead the development process using their conceptual knowledge of their business.

1.6 Research Problems

In MDE, a model has to be designed in such away that is able to express all critical aspects of the system. This requires skilful developers to construct these artefacts. This need for technical skills prevents the MDE approaches from being suitable for end-users [91]. As in the reviewed approaches (Chapter 2), end-users require a degree of technical knowledge in order to be qualified to model their system at a higher level of specification. For example, in MDA [89], end-users might be overwhelmed by the need to learn various OMG [86] standards, such as UML [70], OCL [85], and QVT [88], to be able to construct a rich Platform-Independent-Model (PIM) model that consists of adequate details for enabling a full code generation.

Adopting a DSM approach might solve this issue partially. Although the DSM approach succeeds in raising the level of abstraction, focusing on the problem domain instead of the implementation details, end-users still need to learn a new limited language for a particular domain, even if they were familiar with its concepts. For instance, end-users in the WebML [124] approach should specify all composition and navigation features of their web application using a number of design languages. The modelling process starts with constructing the data model and ends up with designing the hypertext and presentation view. According to that, end-users must act as a web designer to design each part of the system. From that, it can be argued that these approaches are appropriate for designers or engineers rather than business end-users. The highest level of abstraction requires, to some extent, technical skills to express constraints and behaviours.

The premise for this thesis is that UML [70] is too unwieldy to serve as the basis for model-driven engineering. The models in UML [70] are too complex and eclectic to be given a single, clear interpretation, while paradoxically not covering all of the views that are needed to completely specify a software system. We propose some ideas for a simpler notation, with a cleaner semantics, in which the iconography is more consistent. Individual models are smaller and more restricted, but there are more kinds of model to cover the different interlinking views of a system. As a result, it is possible to specify partial and total transformations between different kinds of model. The result is known as μ ML, or the Micro-Modelling Language.

 μ ML aims at raising the level of abstraction to suit business end-users, enabling them to construct their system easily. Furthermore, it reduces the ambiguity of requirements and troubles that occur during client-designer communication.

1.7 Contribution of the Research

The research contribution here aims at raising the level of abstraction to suit business users, enabling them to construct their system easily by themselves, using less technical knowledge in a more efficient way than existing approaches. It can be said, this tackles some software engineering issues in requirement elicitations to accelerate the development process and meet end-user requirements. For instance, allowing end-users to act as system designers to express, model, and customise their functional requirements might reduce the ambiguity of requirements and troubles that occur during client-designer communication. The investigation aims at exploring to what extent the captured functional requirements, supplied by end-users, can generate a complete information system that meets their need. This raises a number of initial research questions as follows:

- What is the required interaction mechanism between the end-user and the transformation approach to capture their desired system specifications?
- What kind of abstract system views do we need to capture in order to have comprehensive knowledge and behaviour of a system?
- What transformation rules are required to fold and optimise high-level views and introduce extra detailed design information of a system?
- What transformation rules are required to refine high-level models and introduce richer design information to a system?
- At which level of development is end-user engagement required to supply new knowledge in order to be considered by model translators at the next translation step?
- To what extent are end-users able to generate a complete system for their demand?
- With respect to code generation, to what extent are we able to construct and link the information system layers from a generic and less technical specifications?

1.8 Objective of the Research

In technical terms, we are aiming to introduce a Model-Driven Engineering framework that supports various model transformation mechanisms to generate information systems (ISs) with a backend relational databases. Although, there are several MDE/MDA and DSL approaches, our work, to some extent, is distinct in its modelling and transformation approach, in which:

- Our transformations implement folding strategies between different views of ISs rather than translating each high-level view into a target code to represent a part of the system (user interfaces, or business classes). This means, new concepts might appear in models during the transformation to be used at the code generation phase.
- End-users are able to customise, modify and then generate their updated systems from intermediate level, rather than executing the whole transformation from the beginning (highest level), as it occurs in forward engineering.

Besides this, we are following such a liberal modelling strategy of the construction of highlevel models using simple Java approach for constructing models and rules, which is more accessible than having to learn transformation languages such as ATL [33], Kermeta [52] or other OMG standards (OCL [85] and QVT [88]). Models in the proposed approach are expressed using simpler and more constrained models associated with underlying XML [122] representation in place of rich and unconstrained models in traditional UML [70]. The intermediate models and a final system are derived by model transformations supplied with end-user customisations or some decisions.

1.9 Organisation of the Thesis

The rest of this thesis is organised as follows:

Chapter 2 - Literature Review discusses relevant literature in the field of Model-Driven Engineering (MDE) and End-User Development (EUD). Different paradigms of model-transformation are presented and exemplified using commonly used model-transformation languages. Furthermore, several development approaches of information systems are explored with respect to a metamodel that are used, system's veiws that are considered, development stages, and the suitability to end-users.

Chapter 3 - Framework Overview and Analysis introduces a general overview of the proposed MDE method (BUILD) for developing information systems, including its four development phases. In addition to this, a lightweight modelling language (μML) and the foundation of the concepts of its metamodel are presented formally using First-Order Predicate Logic (FOPL). This chapter is considered a key chapter of the thesis.

Chapter 4 - μ ML Concepts and Notations in the Requirement Sketching Phase describes, in detail, the specification of Micro Modelling Language (μML) models appearing in the Requirement Sketching Phase. For each model, the definition of its syntax and semantics is provided graphically, and formally using FOPL.

Chapter 5 - μ ML Concepts and Notations in the Analysis Phase presents, in depth, the specification of (μML) models appearing in the *Analysis Phase*. For each model, concept specification and notation is discussed and formalised using FOPL.

Chapter 6 - μ ML Concepts and Notations in the Design Phase defines, in detail, the syntax and semantics of Micro Modelling Language (μ ML) artefacts appearing in the Design Phase. For each model, concept and notation is discussed and defined formally using FOPL.

Chapter 7 - The Architecture of Model Transformation Approach discusses the overall structure and mechanism of model transformation approach. The design of the top level framework and the architecture of the concrete model transformation frameworks, via the various development stages of BUILD, are discussed with in detail. Moreover, the architecture of the Code Generation Framework is discussed, including the internal structure of the two domainspecific generators, namely, *Java Swing* and *MySQL* code generators.

Chapter 8 - The Rules of Model Transformation explains and discusses, in depth, mapping rules that are used to derived each element in the intermediate and design models. For each translation step, First-Order Predicate Logic statements are utilised to express formally every rule in that step.

Chapter 9 - Case Studies illustrates the adoption of the designed framework (BUILD) via its associated μ ML models to develop a real-world information system using the University Administration System case study. The chapter includes three main running experiments with a demonstration of their generated intermediate models and final executable code.

Chapter 10 - Evaluation and Testing presents criteria to evaluate the method. Completeness and correctness of results and transformations, as well as the simplicity and the expressiveness of μ ML are checked, traced and inspected using the final output produced from both JDBC Java Swing Applications and MySQL databases generators. Limitations of the proposed framework are also highlighted.

Chapter 11 - Conclusion and Future Work concludes the overall work presented in the thesis, including a summary of findings, contributions with respect to four main dimensions, namely, modelling language, MDE development approach, model transformation strategy, business users participation in the development process. Moreover, possible avenues for future work are highlighted briefly.

2 Literature Review

"The difficulty of literature is not to write, but to write what you mean" Robert Louis Stevenson

2.1 Context

This chapter reviews the literature in the field of Model-Driven Information Systems Engineering and highlights some issues that need to be considered in order to enhance end-user model driven information systems engineering.

The chapter starts first by presenting a brief overview of some of the related model-based methodologies that are used to raise the level of abstraction in software development compared with writing code. Then a more detailed overview of model-transformation paradigms is introduced, including simple examples. This part of the chapter focuses on how each language is able to express, manage the order of execution, design and compose, and reuse, rules of transformation.

The third part of the literature survey discusses some of the software engineering approaches, such as MDE, OOD, and Functional approaches for web applications, enterprise and software systems development. It highlights the features, system views (aspects), semantics and notation of their adopted modelling language (UML Profile or DSL). Furthermore, the last part of the literature highlights end-user rules in some of current MDE approaches that allow business user participations.

2.2 Overview of Related Methodologies

This section provide a general review of the related Model-based development approach, including the type of modelling language, system artefacts (views), trasformation strategies, adopting technologies and related tools.

2.2.1 Model Driven Architecture (MDA)

The MDA standard specification explicitly aims to integrate many Object Management Group (OMG) standards, UML [70], Object Constraint Language (OCL) [85], Meta-Object Facility (MOF) [87], and Query/View/ Transformation (QVT) [88] to produce a coherent MDE approach for managing models, and provide executable systems that are automatically-generated from specifications [93]. According to [62], MDA has three types of models: Computational-Independent Model (CIM), Platform-Independent Model (PIM), and Platform-Specific Model (PSM) [89].

The CIM model acts as use cases and feature-oriented diagrams to represent the business requirements and features [40]. It also describes the system domain and requirements, from a high perspective, without including any computational implementation. Moreover, a *Platform*-Independent Model (PIM) can be defined as a model of a system, that is completely independent of the specific technological implementation. The third model in MDA [89], the Platform-Specific model (PSM), is a model of a system that holds technical implementation detail relating to a specific platform or environment [89].

The PIM model is always considered a source model of the transformation program within within the MDA approach [89] approach to derive PSMs [6]. On the other hand, the PSM model, the output of the transformation phase, is considered the lowest-level model in MDA [89], which contains ideal and adequate information about a targeted platform used [15].

2.2.2 Model-integrated Computing (MIC)

Model Integrated Computing (MIC) [54] is an example of the Model Integrated Development (MID) approach that has been developed over the past two decades as a software development methodology for building embedded software systems. It uses models as primary entities in the development lifecycle to synthesize, analyze, integrate, test, and operate real-time systems [61]. MIC [54] realises MDE using a different strategy than MDA [89], with fairly similar concepts.

MIC [54] emphasises the employment of Domain-Specific modelling techniques (DSM) using well-tailored Domain-Specific modelling Languages (DSMLs) and multiple views of the system to model comprehensively embedded system domains [61] rather than concentrating on the usage of UML [70] only for modelling tasks as in MDA [89]. Models in MIC [54] consist of logical and functional requirements of a system as well as the physical characteristic of its hardware such as power, time, fault, and size. These physical features are included as physical properties within the platform models and are mapped to the requirements of the application models [111].

The MIC framework [54] consists of a built-in metamodelling mechanism to support the creation of DSMLs, and a meta-generation facility to create DSM tools. Instance models, constructed using these tools, can be translated into other presentations using meta-generation features [61]. The graph-based model transformation tool, GReAT [53], is considered an example of existing tools that support the MIC [54] software development.

2.2.3 Aspect-oriented Software Development (AOSD)

AOSD [39] is another model-based software development methodology that provides more powerful localization and encapsulation mechanisms than traditional component technologies [39, 18, 24, 11]. It aims to describe a system using multiple views, such as a data view, a security view, and a business process view, in which the design concerns are separated as different aspects. It emphasises providing a clear separation of crosscutting concerns at the models level. These concerns, together with a primary (base) model, are composed to produce an integrated view of the logical architecture of a particular system in technology-independent terms.

The model transformation process then takes the responsibility for folding these associated aspects together to form the target system. Using AOSD [39] with MDE can contribute in maintaining this separation during the vertical transformation from a high level specification model into low level implementation ones via the designing of an appropriate DSL to demonstrate the localisation ability of AOSD [7].

The Theme [48] approach, for example, is an AOSD [39] method that covers various software development phases, including requirements, analysis, design and implementation. This approach uses a UML extension, called Theme/UML [11] for designing of mobile, context-aware applications, which supports, via an implemented toolset, the automatic generation of PSMs and executable code from a number of generic PIMs (instead of one large PIM) [11, 18].

2.2.4 Executable UML (xUML)

Constructing a complete PIM model, which is totally separated from any implementation details, and then deriving one or more PSMs from it via a series of transformations until generating the final code, are considered common steps in various MDA approaches [89, 24, 54]. Stephen Mellor and Marc Balcer introduced an approach called Executable UML (xUML) [2], implemented as an object-oriented analysis (OOA) tool, that jumps from PIM level to a direct code without a model-to-model transformation step (PIM-to-PSM). They consider the executable code a PSM model that conforms to a particular language definition (grammar). The xUML [2] approach promises to express a system in a platform-independent way using rich UML [70] diagrams to represent its structural and behavioural views. As a starting point of model transformation, UML Class and State Machine Diagrams are adopted to model the structure and behaviour parts of the system respectively.

The detailed procedure, including the execution algorithm and constraints, for each state in the behaviour model are expressed formally using an Action Specification Language (ASL) [1] to express the detailed control flows of a system. This enables and supports the direct transformation from this abstracted level into the target code using a model compiler, unlike other MDA approaches that tend to derive PSMs first, as an entry point for code generation.

The ASL [1] is expressed in a higher level of abstraction than General-Purpose Languages (GPLs), which support the mapping between ASL [1] and various programming languages, such as C++, and Java. However, it is a designer responsibility to write ASL code that is syntactically and semantically correct in order to generate the correct code at the end. Therefore, writing ASL [1] code at the modelling level requires a degree of knowledge about its correct syntax and semantics. From this, it can be argued that raising the level of abstraction at ASL [1] to express the detailed behaviour of a system is not adequate for end-users who do not have this technical knowledge.

2.2.5 Software Factories

Software Factories [47] is a software development methodology introduced by Microsoft, and it is motivated by two existing software development paradigms, namely, Model-Driven Development (MDD) and Component-Based Development (CBD). According to [47], Software Factories is a product line that configures extensible development tools such as the Microsoft Visual Studio Team System (VSTS) with packaged content and guidance, designed for building particular types of applications.

Assembling components is the general strategy of this methodology, in which system functionalities are distributed across these components or assets to form the product family. The Integrated Development Environment (IDE) is used then to collect product family members into a Software Factory template and deliver it as a plug-in to the IDE.

Abstraction, as a core dimension in *Software Factory* innovation, aims to reduce complexity, hide implementations, and specify product features. Its purpose is overlapping with the purpose of raising the level of abstraction in MDE. The MDE approaches have a number of platformspecific model compilers/generators to generate implementation for target platforms from a single specification model. Therefore MDE and DSL play a role in Software Factories such as maintaining the synchronisation between components and constructing the system from high level specifications [47].

2.2.6 Generative Software Development (GSD)

Generative Software Development (GSD) is a software development methodology that aims to achieve an automatic creation of *system-family members*. The system can be generated from its high level generic specifications. The essence of this approach is based on the usage of reusable assets such as models, and components for building systems that are developed using appropriate Domain-Specific Languages. According to several resources, GSD shares some principles of the Component-based, the Software Factory [47], and the MIC [54] approaches, in which all of them seek an ideal integration of associated members to form the final system [61, 24, 47].

Besides this, GSD is considered one of the related approaches to MDE and MDA [89] due to their trend to capture system specifications and to represent them using DSLs, or UML profiles [70]. The mapping from a PIM to the PSMs in MDA matches the mapping from a Problem Space to a Solution Space in GSD. Therefore, the technologies and techniques used to construct the problem space model might contribute to the process of creating the PIM in MDA [24].

2.2.7 Summary

- It is clear that MDA [89] and Software Factories [47] are distinct in their delivery. Software Factories delivers plug-ins that are embedded into the used IDE as an extension forming a new product, whereas MDE and MDA [89] aim to generate a complete software system, that can run in the target environment.
- MIC and *Software Factories* are somewhat related to the Component-Based Development (CBD), for example, a product of *Software Factories* is based on assembling the product family members or assets as plug-in components to form a target system [47]. MIC [54]

develops the run-time components and maps them to hardware resources to form the embedded system [61].

• Aspect-oriented Software Development (AOSD) and Model-Integrated Computing (MIC) [54] have a strong similarity in describing a system using multiple views in which they separate the design concerns in different aspects; the process of folding them together is considered a part of the model transformation process.

2.3 Overview of Model Transformations

Model transformation plays a significant role for supporting and achieving the automation of various Model-Driven Engineering (MDE) tasks such as refining, normalising and refactoring models, synchronizing and merging and weaving [15, 25, 81]. Model transformation is defined as a process of converting one model into another form using transformation rules acting on models at different levels of abstraction [79]. It consists of one or more defined rules for expressing the mapping between the source and target model. Tratt in [115] also defined model transformation as a program or compiler that converts one model into another. The following figure (2.1) demonstrates a general overview of the main concepts of model transformation, by showing the simplest scenario of mapping between an input (source) model and an output (target) model.

Figure 2.1: Concepts of Model Transformation

Transformation in general can be divided into three types, endogenous, exogenous, and inplace, based on the languages of the source and target models. On the one hand, the endogenous transformation is a transformation between models expressed in the same modelling language. Model *refactoring* and *normalising* are regarded as obvious examples of this type of transformation. On the other hand, the *Exogenous* transformation occurs between different languages as in code generation (from a model to a target code), and model translation (translate a model into an equivalent model) [79]. When having only one model involved in the transformation, it means that all modifications occur on this model, called in-place model transformation or modification [81].

Based on the definitions above, model transformations can be divided into two main categories: model-to-code and model-to-model transformation. The model-to-code mapping (code generation) is considered a special case of model-to-model transformation in which a detailed model that conforms to a metamodel is translated into a target programming language that defined using language grammar [81]. It is considered a final transformation stage in all MDE approaches for obtaining a complete executable code. However, model-to-model transformation is known as layered transformation steps for constructing models in the lowest level of abstraction. In Model-Driven Architecture (MDA) [89] for instance, the process of obtaining the

Platform-Specific Model (PSM) model from one or more Platform-Independent Model (PIM) models reflects the *model-to-model* transformation stage.

2.3.1 Model transformation Paradigms

Model transformations are classified into various paradigms based on the chosen presentation style of the concrete syntax of the transformation rules, which are the atomic units of the transformations considered in different approaches. The following is a brief survey of the existing model transformation classifications. The taxonomies proposed in [25, 81] and a survey conducted by [15] discusses the transformation mechanisms offered in a number of transformation paradigms and tools.

2.3.1.1 Direct Manipulation

In this approach, the rules of transformation are implemented directly using one of the wellknown General-Purpose Programming Language (GPLs), e.g. Java. This approach is developed as an Object-Oriented framework for facilitating model management tasks such as organising, scheduling, and controlling model transformations. The user is responsible for implementing the rules of transformation from scratch with little support from the tool [81]. Appropriate reading and writing, and possibly traversing, mechanisms are used to visit and read elements in the source model, apply transformation rules, and then attach translated elements into the output models [15]. Here we review two model transformation frameworks: SiTra [9] and JaMDA [105].

The Simple Model Transformations, SiTra [9], can be considered one example of the directmanipulation approach. It is regarded as a minimal Java-based model transformation framework that consists of a simple Java library for supporting the implementation of algorithms that executes a transformation based on rules. SiTra [9] is not introduced as an alternative to other model transformation paradigms; it only aims to introduce the concepts of transformation rules for experienced programmers without the need to learn any transformation language. It is not designed as a new transformation specification or full transformation framework [5].

SiTra [9] presents transformation rules by considering two interfaces (Rule and Transformer) and a class for implementing a transformation algorithm. In this regard, they provide a standard way of representing rules and transformations using pure Java code (classes) to mimic the behaviour of rule-based transformation system. The Transformer interface is implemented by transformation algorithm class [5], which is based on pattern-matching rules that implements the Rule interface.

The Rule interface is implemented using Java in which a transformation is broken up into a number of rules. This interface includes different methods for checking whether or not the rule is applicable to an input element from the source model, and constructing an output element in the target model. Moreover, the implementation consists of four main methods (two methods for invoking rules, and two methods for mapping between the objects in the source and target models). It takes a transformation rule and an object expected to be transformed as parameters and executes the algorithm explicitly [5]. Listing 2.1 demonstrates the definition of a rule that maps a class to a relational table [9].

Listing 2.1: An example of a transformation class written in Java

```
1 class Class2Table {
2 Table build(Class cls) {
3 Table tbl = new Table( cls.getName() );
4 for(Attribute att: cls.getAttribute()) {
5 tbl.addColumn(
6 new Column (\text{att.getName}), \text{att.getType})
7 . \sigmaetName())
8 );
9 }
10 return tbl;
11 }
12 }
```
As SiTra [9] is introduced as a framework to support the implementation of simple model transformations, its limitations appear when dealing with complex transformations. For instance, in the case of having multiple rules map to the same target object, SiTra [9] has no way to determine which rule is responsible for target object construction and which rule retrieves the object from the target model. Thus, manual contributions, by the designer, are required to decide the creation rule [5].

The second example reviewed here is the Java Model-Driven Architecture, JaMDA [105]. It is another object-oriented framework for generating Java code from UML diagrams [70] of core business classes using their underlying representations (XMI standards). According to [105], the developer usually writes a business logic code in the target programming language, which can be merged into the generated system. It is expected that the developers manually write about 20% of the total code.

JaMDA [105] is noted as a model compiler that follows the direct-manipulation technique of model transformations that consists of a library of independent transformers. Each transformer is designed to handle a special type of transformation [81, 105]. Thus a user can implement horizontal transformations that can optimise a model or vertical transformations that can refine a model towards code. Jamda implements transformers and generators using a visitor-based approach where the input model is traversed and each element has the appropriate transformer to apply the mapping rules [105].

Jamda can be said to be platform-independent, as it is written in Java and produces the source code in Java. However, it does not output any other executable programming languages. For instance, it is used for generating N-tier web applications based on Enterprise Java Beans technology (EJB). Complete UML models [70] that represent the whole application are used to generate each tier. Each tier has a collection of generated java classes that have been derived from the initial UML [70] input models and classes of the previous tier [105].

2.3.1.2 Model transformation approaches for Declarative (Relational) rules.

In this approach, the rules of transformation are expressed declaratively using a set of constraints to specify what the relations between source and target models are [110]. A number of languages and tools that support this relational approach are used, such as QVT Relations [95], Tefkat [73], and XMF-Mosaic. Here we review in detail the QVT Relations language [95].

The QVT Relations language [95] is considered a domain-specific model transformation language that expresses the mapping logic between two models as a set of relations. The relationships between source and target elements are encoded as pattern-matching expressions including predicates (preconditions) that must be satisfied by the input elements [95, 44]. The precondition normally consists of two expressions: an enabling (when clause) that returns a Boolean value, and an enforced (where clause) expression that causes side effects on the target model [95]. Unlike other declarative graph-based model transformation frameworks, namely, GReAT [53], VIATRA2 [32], and AGG [112], that provides unidirectional transformations only, the QVT Relations supports both unidirectional and bidirectional transformations [95]. Listing 2.2 illustrates snapshots of a transformation program writing in Tefkat [73], a model transformation language based on F-logic, to express the relation between persistent UML [70] classes and RDBMS tables.

Listing 2.2: An example of the declarative representation of rules using QVT-Relations

Listing 2.3 demonstrates a transformation rule for translating OO Class to a relational Table using QVT-Relations [95].

Listing 2.3: An example of the declarative representation of rules using Tefkat

```
1 RULE ClassAndTable(C, T)
2 FORALL Class C {
3 is_persistent: true;
4 name: N;
5 }
6 MAKE Table T {
7 Name: N:
8 }
9 LINKING ClsToTbl WITH
10 class = C, table = T;
```
In the case of having more than one transformation rule to derive a particular element in the target model, QVT-Relations takes into account the ability to check whether or not a target object is already created by one rule in the target model. As a result the second rule only updates the created element rather than constructing it twice. Unlike the direct-manipulation framework (SiTra [9]), QVT Relations [95] may offer a strategy to avoid any duplicate creation, or deletion task of target elements and replaces it by an appropriate update task. This is achieved by introducing the notion of identity to target elements [95, 44]. According to the UML-to-RDBMS case
study presented in [95], QVT-Relations adopts the concept of a key from relational databases [95] to handle this issue. This new concept is described as a set of properties, or fields, that uniquely identify an object instance of the class in a model.

Model transformation for Graph-based rules

As a formal approach of the declarative model transformation strategy, Graph Transformation has been widely utilised for expressing model transformations. It provides a theoretical foundation and a formal flavour to transformation frameworks, so, in general usage, source and target models can be expressed visually and textually using the notion of a graph (graph-like structure). The graph consists of a set of vertices $V = v_1, v_2, v_n$ and edges $E = e_1, e_2, e_n$ in which each edge e in E has a source(s) and target(t), vertex $s(e)$ and $t(e)$ in V, respectively. Several types of graphs, such as typed graphs and instance graphs, are specially designed to represent the abstract syntax of diagrammatic class, and the instance models respectively, e.g. a metamodel of the UML Class diagram and an Object diagram [25, 12, 114]. Figure 2.2 illustrates and example of the representation of a system using the visual notation of a graph [110].

Figure 2.2: A graph representation of an Object Diagram

A graph transformation rule (T) between models is formally considered a function f: $L \longrightarrow R$ that maps elements in a source graph SG to elements in the target one TG. The transformation typically consists of a pair of graph patterns [25], left-hand side and right-hand side graph $GT = (L, R)$ in which the union $L \cup R$ is defined. For instance, edges that appear in both L and R are connected to the same vertices in both graphs. The left-hand side pattern is considered a sub-graph of the source model that describes the precondition of a particular rule, whereas the right-hand side pattern is regarded as a sub-graph of the target model that describes the post-condition of the rule [15, 12].

The above technique is successful in *one-to-one* and *many-to-one* mappings between source and target graphs. However, in the case of many-to-many, the transformation (function) has to be reconsidered as a *relation* (R) in order to allow this kind of mappings. The Triple Graph Grammars (TGG) [63] technique introduces a solution of this issue to represent the concrete syntax of transformation rules as a graph-like structure. Each mapping rule consists of two graphs (left and right), as well as a Correspondence Graph $TGG = (L, C, R)$ to enable the double push out strategy between the three graphs. This strategy has been adopted in many graphbased model transformation tools such as AGG [112]. For example, in the basic mapping between Attributes and Fields in the Class-to-RDBMS transformation, the many-to-many mapping can be found in translating *Attributes* to either columns or key columns. Figure $2.3(A)$ shows how many-to-many mapping is splitting up into two many-to-ones using correspondence graphs $(Attr2Fld, and Attr2KFld)$, whereas this is not obvious in the (B) side of Figure 2.3.

(a) Triple Graph Grammar

Figure 2.3: Example to M-to-M mapping using TGG

There are a number of frameworks and tools for graph-based model transformation, such as Fujaba [114], VIATRA2 [32], AGG [112] and GReAT [53]. Fujaba [114] is considered one of the most popular general-purpose graph transformation tools [98] for generating Java code from the UML specification [70] and for round-trip engineering. It has the ability to transform many source models to many target models. It provides an endogenous transformation only [61, 80, 121] in which the source and the target models conform to the same metamodel. Fujaba [114] uses the graph technology to represent the underlying representation of UML [70] diagrams with formal foundations. The transformation rules that are declared using the TGG [63] technique are implemented as method bodies with a control structure [114].

Fujaba [114] is able to generate code from the UML Class Diagram, a collection of Story Diagrams, and rules of transformation expressed using TGG [63]. The Story Diagrams are used to model the system behaviour. A story is used to specify the body of methods related to a class. Whereas, the UML Class Diagrams are used to model the structure part of the system. Consequently, a code generator can generate Java code for these story and class diagrams [13].

2.3.1.3 Model transformation approaches for Imperative (Operational) rules

KerMeta

KerMeta [52] is considered an Object-Oriented meta language that is used to specify the structure and concrete syntax of models with the ability to define their behaviour using operational semantics. It is fully integrated into Eclipse, in which it offers an EMF-based metamodel, and is used as a MDA [89] model transformation language or tool to specify the mappings between source and target models [52, 57]. It is an extension to Essential MOF (EMOF) with an action language. The language is designed based on two existing languages, namely Xion, and MTL. Therefore many features of these two languages, such as the action language (Xion), and multiple inheritance (MTL), can be re-expressed in KerMeta [84].

This action language is an imperative OO language used for defining operations and constraints [55] in metamodels. It is also used for implementing executable metamodels to provide all model management tasks [57, 65], and for performing transformations [15]. KerMeta [52] supports many OO language features such as static typing, and multiple inheritance, and also it supports more specific concepts to be suitable for modelling and model transformation tasks. For example, many OCL [85] operations such as *collect, select*, and each are available in KerMeta [52] to apply such a condition and dealing with a certain collection (Listing 2.4).

Listing 2.4: An example of the representation of the transformation rules using KerMeta

```
1 getAllClasses(inputModel)
2 .select{c| c.is_persistent}
3 .each{c| var table:Table init Table.new
4 Table.name:= c.name<br>5 Class2Table.storeTr
            5 Class2Table.storeTrace(c,table)
6 Result.table.add{table}
7 }
```
2.3.1.4 Model transformation approaches for Hybrid rules

AtlanMod Transformation Language (ATL)

ATL [33] is a MOF-based model transformation language that supports both imperative and declarative representation (hybrid) of mapping models (Figure 2.5). This allows users to express a complex transformation that is hard represent declaratively using an ATL [33] imperative structure. Its representation of mapping rules forms a transformation model (program) that describes how to create the output model of transformation [15]. ATL transformation models are implemented as ATL Modules that consist of a number of transformation rule definitions (rules), and some helper methods. Listing 2.5 demonstrates a snapshot of an ATL module for defining the transformation between OO Classes to Relational Tables.

The ATL [33] approach delivers only part of what MDE promises, in that it supports only the model-to-model transformation [33]. ATL [33] is not directly based on the mathematical foundation of graph transformation techniques [59]. ATL [33] has abstract and concrete syntax that are expressed textually [15, 41], and a particular compiler and virtual machine [110]. It is essential to make a clear distinction between source and target models, since ATL [33] cannot use the same model for both. The input model here is considered a read-only model. The navigation task on the input model is achieved using ATL query units. These queries define the navigation over one or more source models to produce a value. Conversely, the output model is not navigated and considered a write-only model [123].

Listing 2.5: An example of the representation of the transformation rules using ATL

```
1 module Class2Relational;
2 create OUT : Relational from IN : Class;
3
4 rule ClassAttribute2Column {
5 from
6 a \cdot Class k<sup>1</sup> \uparrow k<sup>1</sup> \uparrow \uparrow \uparrow7 a.type.oclIsKindOf(Class!Class) and not
8 a.multiValued)
9 to
10 out : Relational!Column
11 name \leq -a.name + 'Id',
12 type <- thisModule.objectIdType)
13 }
```
Epsilon Transformation Language (ETL)

Epsilon [34] provides a framework called the Epsilon Model Management Infrastructure that consists of a family of task-specific modelling languages such as the Epsilon Merging Language (EML), Epsilon Comparison Language (ECL), and Epsilon Transformation Language (ETL) [34]. The ETL [66] is another hybrid model-to-model transformation language that is designed in such a way to be integrated smoothly with other task-specific languages within Epsilon [34, 100]. Unlike ATL [33] that is designed to perform many model management tasks such as model validation and merging, ETL [66] is a DSL for implementing transformations [67].

Listing 2.6: An example of the representation of the transformation rules using ETL

```
1
2 rule Tree2Node
3 transform t : Tree!Tree
4 to n : Graph!Node {
5
6 \mid n.name = t.label;
7
8 // If t is not the top tree create an edge connecting n
9 // with the Node created from t's parent
10
11 if (t.parent.isDefined()) {
12 var e : new Graph!Edge;
13 e.source ::= t.parent;
14 e.target = n;
15 }
16 }
```
Epsilon has an OCL-based core language, Epsilon Object Language (EOL) [64], that is used in developing the Epsilon family of task-specific modelling languages. The ETL [66] uses the navigational features and other imperative features of EOL [64], such as accessibility to multiple models, model modifications, programming constructs (e.g. loops, branches) to express complex transformations [66]. Similar to ATL [33], transformations in ETL [66] are implemented as a number of modules, containing a set of *rules* and *operations* to support the imperative side of the rule [66]. Listing 2.6 above illustrates a snapshot of an ETL [66] module for defining the transformation from a Tree node to a Graph node.

2.3.1.5 Rule Scheduling Mechanisms

The order of rules execution can be specified using two styles: *implicit* and *explicit* forms. In the implicit style, rules call another rule without referencing explicitly the rule name; that is, in which the calling mechanism is based on rule dependencies. The transformation engine has a scheduling algorithm that determines the order. In general, a rule A is executed when another rule B is waiting for the result produced by the rule A. This mechanism is common in declarative model transformation languages such as QVT-Relations discussed above [69].

In contrast, the explicit style might use control flow structure within a rule including explicit rule invocations (internal). It also can be external by using logic to separate transformation rules without a direct invoking of each other. This style is usually found in imperative and hybrid model transformation languages such as ATL [69] and ETL [66].

2.3.1.6 Template-based Approach For Code Generation

According to [15, 25], a template-based approach is widely adopted in the majority of MDA [89] tools such as AndroMDA [8], OptimalJ [19], and Acceleo [3], for supporting the model-to-text transformation. A template of a target model that contains a meta-code is accessed, mapped, and filled from the source model. The structure of the template is closely related to the generated code [26].

AndroMDA

AndroMDA [8] is an open-source Model-Driven Architecture (MDA) framework. It takes one or more graphical UML models [70] and produces a target component in one of the J2EE technologies such as Java, PhP, and EJB. In the MDA [89] process, AndroMDA [8] takes the Platform-Independent Model (PIM), resulting from the analysis phase, and refines it to construct a Platform-Specific Model (PSM) for a target technology with a template for producing the final code.

AndroMDA is regarded as a template-based approach for code generation that is able to generate different code based on the language of the template [8] using language-specific cartridges. The user can customise these template files to produce a source code in any programming language. The metamodel in AndroMDA [8] is based on MOF [87]. The cartridges are designed to get the information necessary to generate a target code from MOF models inside a MOF repository. The AndroMDA Cartridges, primary plug-ins in the framework, are designed to perform this process by parsing the underlying XMI representation of UML diagrams collaborating with transformation libraries. These cartridges consist of code templates that are written using the Apache Velocity template engine [8].

Acceleo

Acceleo [3] is a template-based MDA [89] code generator approach for generating code for various platforms such as Java, $C++$, $C#$, PHP, JEE, and Python. It is fully integrated with Eclipse and the EMF framework (GMP) in which it offers a user-friendly template editor for organising the template of target code. The framework consists of a number of separate technology-specific modules for code generation. Each module is created from templates that express the information required to generate executable code from a metamodel. Within each template, several scripts allows the user to customise the generator accurately [3]. The basic idea of this approach is that the generation program (script) is applied to a source model, which conforms to an EMF based metamodel, to fill the target template.

The modularity supported in the Acceleo architecture enables adding new generators as plug-ins without any problem, which accelerates the implementation process in general [3]. It could be argued that Acceleo can deliver a complete MDA [89] solution or implementation when it is integrated with other Eclipse-based tools such as ATL that supports model-to-model transformation.

ZOOM

Z-based Object-Oriented Modelling notation (ZOOM) is a model-driven engineering framework that is based on a formal modelling notation. It aims at improving the quality and productivity of software development processes by handling existing drawbacks, such as incomplete modelling notations, and the lack of an effective model transformation mechanism. It is a template-based model transformation approach that is supported by a CASE tool, called Hierarchical Relational Metamodel Transformation (HRMT) [56]. Transformation templates (cartridges) are considered a collection of transformation rules that define the mapping between source and target models. The framework consists of a set of templates; each one is used for performing a specific transformation task on a target platform [76]. End users are able to add their preference only at user interface generation [56].

ZOOM has a simplified metamodel called Hierarchical Relational Metamodel (HRM) that maintains a tree structure and relationships representation between PIM elements. The modelling notation, which is consistent with UML 2 [70], has a textual syntax defined by BNF and a similar internal representation of models to programming languages, which is an Abstract Syntax Tree (AST). It uses mathematical collections to depict such complicated modelling language constructs as associations [76].

The functional requirements in the framework derive the structural, behavioural and UI models. ZOOM provides a pre-defined event model, which is processed by an event-driven framework, to bind the structural, behavioural, and UI models together. The integrated ZOOM model will be processed by the Knowledge-based Model Compilation Tools resulting in different implementations of the software system based on the specific platform and knowledge base. To start the MDE process, the developer needs to build a platform-independent model and other models are mostly generated using several steps. Firstly, the textual representation of ZOOM notation is parsed into an abstract syntax tree (AST). Then a post-order traversing mechanism is used to visit all elements in the AST and applies mapping rules to generate the code [56].

2.3.2 Composition of Model Transformations

Model transformation plays a significant role for supporting and achieving the automation of various Model-Driven Engineering (MDE) tasks, such as creating, refining and refactoring models, as well as synchronizing and merging and weaving. As a result of its critical roles within various MDE tasks that are focused on the mapping, checking and validation of models, the transformation units are becoming more complex and harder maintain. Therefore, the composition of model transformations has emerged as an important research topic in its own right, allied to model transformation in general. Tasks include designing a suitable orchestrating mechanism for controlling the execution of the decomposed transformations, in order to produce a single consistent result.

Many composition techniques have appeared to handle this issue. For instance, one technique is achieved by simply linking several pre-designed model transformations (external) whether they are expressed using one or different languages and executed by one or several tools. On the other hand, the other technique is based on decomposing the rules of transformation into independent units (internal) to collaborate as a big transformation unit to perform a complete transformation [46]. Here, we discuss some of the limitations that exist in the current model transformation approaches, particularly, the rule-based and direct manipulation approach, and review an example of a composition technique in the graph-based and rule-based, model transformation paradigm.

2.3.2.1 Composition Techniques

Hidaka et al. [50] developed a composition strategy for a graph-based model transformation approach, inspired by the Unstructured Data Query Language (UnQL), the compositional graph querying language. UnQL is a powerful query language that is based on a pattern-matching technique to express queries in a structural recursion style. This style enables the composition of two or more queries that are designed in a structural recursion way to be expressed as a single query [17].

Hidaka et al. [50] extended the UnQL by adding three editing structures that support the direct specification of graph transformation, namely, deleting, extending, and replacing a subgraph. At the end, they came up with a graph transformation language UNQL+ that supports the composition of graph-based model transformation. This homogenous composition approach is implemented as a framework that supports the development of model transformations in-thelarge. In this instance, a large model transformation can be systematically designed by gluing simpler model transformation units together via efficient intermediate models, as illustrated in the Class-to-RBDMS transformation case study [50].

When transforming attributes to columns of a particular table, there is a need to gather all information from directly and indirectly associated classes that have relationships with that table. Therefore, designing an intermediate model to associate directly all indirect associated classes to a particular table as a simple transformation (query) is an example of efficient intermediate models. From that query result, it is possible to generate directly primary/foreign keys for the specific table.

Wagelaar [123] proposes a technique for rule-based model transformations which he calls Module Superimposition. It is implemented using the ATL transformation language [33]. Module Superimposition is an internal composition style that allows for the constructing of smaller maintainable and reusable transformation modules, based on transformation rules, to perform together a mega transformation.

The basic idea of this technique aims at splitting up transformation modules into modules of manageable size, which each consists of a number of rules that are superimposed on top of each other and are executed as one rule. It is equivalent to applying a union with override operation to transformation rules that have the same names in a pair of transformation modules, in which the original rule is replaced (overridden) by the new one. The rule then will be executed once. Therefore, the approach is about a rule overriding technique that operates at the ATL helper methods and helper attributes levels [33]. It can also work in the QVT Relations [95] language by considering the transformation rules (relations) as the atomic level of modularity [123].

Goknil et al. [46] provide a composition approach based on two levels of granularity using two transformation languages (ATL [33] and Tefkat [73]). The first level is a traditional rule-based composition supported by ATL [33], whereas the second one is an operation-based composition for handling complex structures in the source model, supported by Tefkat [73]. The external level of composition aims to improve the quality of transformation rules by reducing the number of rules used in each transformation module. It considers the basic transformation operations such as Create, Read, Update, and Delete operation as atomic parts of transformation instead of whole transformation rules.

The idea behind this approach is to avoid the problem of scattering changes on several mapping rules when a change has occurred on the source pattern. This problem affects negatively the consistency and modifiability of transformation rules. Their solution, proposed here, is to consider the source's complex structure (pattern) as a distinct construct and then map it into a single module provided by a language that supports the finer grained decomposition (e.g. Tefkat [73]); then it will be executed as a single transformation rule. In rule-based model transformation languages such as ATL, the transformation rule consists of three patterns, namely, source pattern (SP), target pattern (TP) and action.

$TR = SP, TP, Action$

The source and target patterns are elements in the input and output models, which might be simple or complex constructs, whereas the action is an atomic part of the rule, which might include one or more operations, such as the basic CRUD operations. These operations might be composed into a single transformation rule, which is not supported in ATL (action part,which cannot be composed of multiple operations). However, using another model transformation language such as Tefkat [73] enables having multiple operations to be composed as one action part in a single rule. The action performs one or more operations, which part have three patterns, namely, source pattern *(SP)*, target pattern *(TP)* and operation type $(OpTy)$. The types of composition operations are specified in the operation type part.

Action =
$$
(Op1 + Op2 + + Opn)
$$

Operation = SP, TP, OpTy

From that we achieve two compositions at different levels of granularity. The first composition is between the transformation module (specification) and rules, whereas the second level is between the action part of the rule and basic operations.

2.3.3 Limitations in the Current Model Transformation Approaches

According to [119], the rule-based model transformation approach suffers from several deficiencies that prevent a full implementation of reusable transformation components forming as a transformation chain. For instance, in the Relational approach, e.g. MTF, where declarative mapping rules express the relationships between source and target elements, the transformation can be initiated from any of these rules resulting in (producing) different outputs. Therefore, there is a need to determine an initial rule and also the direction of transformation (forward, or backward).

Furthermore, in the hybrid model transformation approach (ATL) where the transformations might be considered metamodel-independent, there is a need to select concrete metamodels when executing a transformation. This requires studying the implementations in order to identify a set of valid metamodels for each mapping rule. From these it can be said that the problem of having incomplete knowledge of the transformation's implementation exists in various types of rule-based approaches.

The direct-manipulation approach also has limitations in developing reusable transformations; particularly where there is a demand to specify the entry point of the transformation and to have complete implementation knowledge in order to use the designed transformations. As seen in SiTra [9] and ReMoDeL [101], each framework adopts a specific way to use models and to implement the mapping rules using the general-purpose programming language, Java. There is no standard technique to develop, construct, manipulate and specify source and target models. Therefore, extra specifications of a transformation are required in order to design reusable transformation components.

Unlike the declarative rule-based approach, the entry point (initial rule) of a transformation and the order of rule execution are imperatively determined in the old version of the ReMoDeL Database Generator [108] by the transformation engineer during the creation of the transformation. The framework handles this issue using an execution algorithm implemented imperatively as a series of method invocations in the top class of the translators/generators hierarchy, and controlled invocations within other methods in other sub- translators/generators classes. The translation algorithm is generic for all database normalisation scenarios, in which we follow logical normalisation steps taken by database designers, whereas it might be varied in generating constraints based on the target database vendor (e.g. generating triggers instead of field constraints).

2.3.4 Summary

This part of the survey discussed various model transformation paradigms and languages. In the direct manipulation approach, two frameworks are considered, namely SiTra [9] and JaMDA [105]. For both approaches, the use of Java for encoding the transformation rules, as well as the design of transformers (Java classes) are discussed.

Regarding the rule-based transformation style, the imperative, declarative and hybrid approaches are demonstrated with some transformation language examples. For each approach, the policy for rule definition, rule ordering and rule invocation are discussed and supported by small examples to show the ability of the language for expressing the transformation.

In relation to the above, the work developed in this thesis adopts the hybrid approach for expressing transformation rules using Java. The related works are reviewed in this part of the survey. Any rule may be a top-rule; and there is no distinction (as with some other MT approaches) between top-rules and dependent rules. The following table (2.1) summarises general model transformation features (reviewed in the literature) that are adopted in the proposed method in the thesis.

Table 2.1: Features of Model Transformations

2.4 Software System Modelling Techniques

Whereas the above considered approaches used in model transformation, below we consider general modelling approaches in software engineering. We select those approaches which bear some affinity with model-driven engineering.

2.4.1 Object-Process Methodology (OPM)

Object-Process Methodology (OPM) [31] is regarded as a holistic approach for system development. It is supported by a CASE tool (OPCAT) [90] for enabling the generation of complete systems from specifications. It integrates the object-oriented and process-oriented paradigms to introduce a single-view MDE approach. This approach produces a size-wise (fewer lines) code in contrast to other approaches based on UML models. This is demonstrated in the comparison between the UML CASE tool Rhapsody [51], invented by ILogix, and OPCAT [90] tool proposed in [96].

In a multi-view approach, the generated code from each view is limited and represents a partial code that needs be filled by further details generated from other views. This might cause inconsistencies in the code output by different generators, as well as causing incompetence of generated behaviour. OPM [31] tries to tackle this issue due to the crosscutting representation of the system specifications using multiple views. This is achieved by providing a complete model of a system in a single model (view) [96].

The OPM [31] model consists of a set of interconnected OPM [31] concepts or entities, represented graphically via a workflow-like Object-Process (graph) Diagrams (OPD), and textually via a dual-purpose and natural-like Object-Process Language (OPL) based on context-free grammar [60]. The metamodel introduces two general concepts (things with states, and links). The thing concept is to represent an object or a process, and the *link* is used to express structural or procedural connections between things to connect objects to processes. Processes are considered as transformation behaviours for translating objects [60]. Therefore, it can be argued that OPM [31] has a notation that combines static and dynamic aspects of the system in one view, which facilitates generating a complete code at the end rather than code skeletons [96].

OPM [31] has its refinement/abstraction mechanisms that enable the representation of a complete system at different levels of abstraction without losing the overall consistency of its internal components. These techniques are applied to OPM [31] entities as follows:

- Folding/unfolding mechanism for refining/abstracting the structure hierarchy of OPM [31] objects.
- In-zooming/out-zooming for hiding/exposing the inner details of OPM processes.
- State expressing/suppressing for hiding/exposing the state of OPM objects.

As the OPM methodology [31] is supported by a CASE tool (OPCAT) [90], the OPM-GCG is the generator component for generating generic code. It consists of two main parts, namely, Template for Implementation Programming (OPCAT TIP), and Implementation Generator [96].

The OPCAT TIP enables users to encode transformation rules in a language-specific template (Template and Translation DB). The translation is expressed imperatively as an ordered set of operations that contains conditions and actions (Event-Condition-Action (ECA) paradigms). The TIP then generates corresponding XML [122] files. These XML [122] documents are used as inputs to the Implementation Generator [96].

The system uses the underlying representation of the OPM model, which is an OPL-XML [122] script (automatically generated from the graphical notation OPD), with a template to create the actual system, including UI, code, and database schema. The Implementation Generator takes a language-specific OPL template and OPL-XML [122] scripts as inputs, and captures the match between corresponding elements. It executes required operations, and then generates programming languages encoded into XML [122] format. The next step is parsing these output XML [122] files of the system, which is based on user preferences to select transformation options: code (e.g. Java), mark-up (e.g. HTML), or comment [96, 31].

2.4.2 Pure Object-Oriented Method

It is a widely-recognised that a multi-view modelling approach is adopted by various development methodologies and tools, such as the Model-Driven Architecture (MDA) [89]. Generally, in the UML-based approach, each UML [70] model is used for representing a system from a different angle; for example, the approach proposed earlier by Chow et al. [20]. is regarded as a two-stage multi-view MDE approach for generating Java code based on five UML models [70], namely, Class, Component, Statechart, Sequence, and Activity diagrams. The approach captures the static structure of a system at the first stage and then adds the behaviour in the second one [20].

Rhapsody [51] by ILogix is another example that is able to generate partial behaviour of a system. It uses the UML Class and Statechart [49] with translation rules for generating the system, the UML Interaction diagrams to define objects and their interconnections, and the UML Use case and Activity diagrams for analysing and documenting the system [96]. In both examples, it can be argued that maintaining the consistency between these UML views (models) is not a trivial task.

ZOOM Approach

ZOOM separates software modelling into three components: the structural, behavioural and User Interface (UI) models. This separation of concerns allows each aspect of the system to be specified separately, and makes it possible for us to use an appropriate formal specification language to specify each of these unique aspects. The Structural model is defined formally, and visually represented by UML 2 class diagrams. It is completely separated from the other aspects.

Besides this, the Behaviour model is regarded as a communication component that links the structural model with the UI model. It is formalised using state diagrams. It is represented visually using a UML 2 Statechart diagram, and textually using a UML 2 Finite State Machine (FSM) that contains rich syntactical grammar with formal semantics. Therefore, the behaviour model will consist of a number of FSMs for different user behaviours. The changes (transitions) from one state to another are based on user inputs and commands.

ZOOM uses a separate model for representing the UI. As a result, any change of the interface specification will lead to a change in the UI models only without any crosscutting effect in other types of model. The UI model is formally defined using pre-defined XML [122] schemas ZOOM-UIDL as a textual and concrete syntax.

2.4.3 Integrated Methods

The examples mentioned above are based only on the OO method for modelling system specifications. Several attempts have been made to integrate the OO method with the Process-Oriented and Structured (Functional) Methods to enable better and more comprehensive capturing of the system specifications at the design level. Here, we consider some of them to show how this integration contributes to the MDE strategy.

FOOM Approach

Functional and Object-Oriented Methodology (FOOM) is an analysis and design methodology for developing Information Systems (IS). It combines two well-known software engineering paradigms, namely, the OO approach and the Functional approach. FOOM aims to cover the structural and behaviour representations of ISs [60].

FOOM adopts the ADISSA methodology (Architectural Design of Information Systems) that extends the *Structured Analysis Methodology*, during the analysis and design stages. ADISSA is used to extract: (a) menu-tree interface, as an external view of the system to users, and (b) system basic transactions, as an internal view to represent the user interaction with the system via different events. In addition to this, (c) database schema, and (d) system Input/Outputs are also obtained from above using the ADISSA method, which is supported by a CASE tool [60].

The analysis phase of FOOM produces a data model and a functional model. The data model is noted as an initial class diagram without methods. It represents only real-world data entities derived from the requirement analysis stage. The detailed methods will be added later in the design phase. The functional model is considered a hierarchical *Object-Oriented Data Flow* Diagram (OO-DFD) that includes initial classes instead of data stores (external entities). It is used to specify the functional requirements of a system. Like traditional DFD, the OO-DFD consists of a number of functions (decomposable and elementary), and external entities (e.g. user, time entities) [60].

The design phase produces a complete class diagram and user interfaces (UIs) based on the ADISSA methodology. The class diagram contains detailed descriptions of methods, including input and output screens. This phase is organised in sub-levels, one for extracting system transactions from the OO-DFD model. Each transaction contains a chain of functions and external entities to perform a particular process of the system. This will produce the detailed descriptions of the transaction-methods (high-level descriptions) attached to classes. The process logic of transaction at this level is expressed textually using standard structured programming such as pseudo code, or visually using message charts [60].

The next step is constructing a menu-tree interface that is derived from the OO-DFD (the second level according to ADISSA methodology). A menu item is created for each function connected to a user entity and attached in a parent menu forming a hierarchal menu tree. A generic class Menu will be added to the Class diagram, and all extracted menus become objects of it [60].

Input and output screens are designed in the third step. For each input/output command in the high-level descriptions of transaction, there is a form and report screen respectively. Generic form and report classes will be added to the class diagram and all form and report screens become become instances of these classes. Then behaviour methods are designed in the last (fourth) step by converting the high-level transaction descriptions into a detailed description of methods, such as CRUD operations [60].

MOSYS Approach

In addition to this, MOSYS is another integrated approach that is aimed at the construction of Distributed Real-Time Systems (DRTS). It combines the functional and OO methodology and uses UML [70] with an extended DFD (E-DFD) in order to enable the automatic identification of objects and classes from high-level specifications. The components of E-DFD are weighted processes and weighted attributes. The weights are used to model the time constraints and complexity of processes [13].

After identifying the entities that interact with the system, a use case model is created to describe the system functionality. Then activity diagrams or E-DFDs are used to model the use case functionality. Then the functional model (e.g. E-DFD) is mapped into a graph that consists of edges and tasks (nodes). Based on clustering criteria, different object identifications might be obtained [13].

Besides that, [38] a practical technique in integrating notations of UML [70] and the DFD is proposed that aimed at the development of embedded systems. The combined approach is introduced as a part of an MDE process in which a system is modelled using different views; model transformations then force the integration rules between these views. In this approach, the designed DFDs can be transformed into UML object or class diagrams according to the transformation algorithm.

The development process consists of a series of mapping between different models (views). It starts decomposing each use case of an embedded system into three objects: data, control and interface. Therefore, a transformation step is performed on the use case model to produce an initial object diagram (IOD) including refactoring processes on it. Then the data flow model is obtained from the initial object diagram representing different details of the system. This step involves specifying data stores and processes, and identifying the connection between them [38]. Then an object-oriented model is derived from the constructed DFD using one of the following approaches:

- Applying a direct mapping to an object diagram, in which processes and data store will be mapped into objects.
- Applying advanced mapping into a class diagram, which classifies processes and data stores within the DFD based on their type. Processes become logical methods, whereas each data store is translated into a separate class.

DFD net Approach

Another novel approach has been introduced by [113] to complement the OO method with functional decomposition for realising uses case using extended DFDs (DFD net) at the analysis stage. This means that all processes and data flows are transformed into OO classes, including attributes and operations, at the design stage.

In DFD net, a distinction between various data flows has been considered. For instance, distinctions between main-input, inter-process, and non-inter-process data flows are made using double-solid, single-solid, and single-open arrowhead respectively. The process starts by specifying a main-input data flow for a collection of primitive processes in the lowest-level DFD net and connecting them with a local buffer to form a single group. Similarly, processes that share the same data as input or output are grouped together, then classes are generated for each separate group. The development process at the design stage is automatable and supported by a tool to perform transformations [113].

2.5 Metamodelling approaches for Automatic Generation of Information Systems

This section aims to highlight different metamodelling approaches for generating information systems. It considers system aspects that are taken into account during different modelling stages (analysis and design), as well as the degree of automation for generating a complete executable code, user interface and business logic. A domain-specific language framework (MOD4J) [77] for developing administrative enterprise applications, and a framework for developing Data-Driven Applications (XPage) are not included in this survey.

Different metamodelling approaches are proposed for developing web-based information systems (IS) such as [28, 118]. These approaches are normally based on capturing the structural aspects of ISs (entities, relationships, and constraints) using metamodels or appropriate domainspecific languages. They allow the automatic generation of user interfaces (UI) and in other cases basic CRUD operations. These approaches have remarkable limitations in which they suffer from the lack of business logic support and UI development [27].

2.5.1 Enterprise Information Systems (EIS) Metamodel

With regard to these limitations, several approaches have appeared to overcome the issues mentioned above. For example, the EIS approach [27, 28] for developing information systems that integrate business domain, processes and user interface descriptions is proposed to enable automatic code generation. Three separate metamodels for each aspect are introduced, namely the, Business Domain, Human-Interaction, and Business Process Metamodel.

The Business Domain Metamodel (BDM) uses the specialisation strategy instead of using logic at another modelling level to solve data modelling problems when a given entity type participates in an association type with various roles. On the other hand, the Human-Computer Interaction (HCI) metamodel expresses the appearance and behaviour of UI concepts at a higher level of abstraction. Each behaviour triggered by the UI is treated by either a business rule or application functionality generated to connect to one or more UI rules. This information is specified in the metamodel as application rules for providing the integration of UI with the information system [27, 28].

The framework contributes to integrating concepts of HCI patterns with business data aspects, for forming an Information System metamodel. Thus, the integration between these aspects is promoted via metamodels. It also supports the usage of OCL [85] for generating stored procedures in the backend database system that implements the modelled business rules. It does not offer any end users customisation. [27].

2.5.2 Metamodel to support End user Development of Web-based Business Information Systems

Another metamodelling approach has been introduced for enabling end-users to contribute in web-based information systems development [29]. The semantics of the proposed metamodel helps end-users, using their logical thinking and domain knowledge, to start the development process with little knowledge in the technical domain.

The proposed concepts in the hierarchical metamodel are similar to those in the metamodel of the UML-based Web Engineering (UWE) approach. The main difference is the three levels of abstraction in [29] approach. They indeed grouped related common aspects at a separate level of abstraction and produced a three-level hierarchal metamodel for describing information systems, namely, Shell, Application, and Function level. Common aspects in all web applications such as user, object, access control and navigation are four models that are expressed at the shell level. Access control and other specific aspects (inherited from shell level), such as workflow (business rules) are modelled at the application level. The detailed implementation of functions is modelled at the function level [29].

End users might contribute in the development at a particular stage or level; for example, to configure a specific web application they only need the knowledge presented in the shell metamodel. Moreover, in order to specify the application level, users need to use the domain knowledge presented in the application metamodel.

2.5.3 Using Weaving Models to Automate Model-Driven Web Engineering

In [120], a model-driven approach for the development of Service-Oriented Web applications (SOD-M) the use of model weaving is proposed. It is noted as a part of the MIDAS framework for developing Web Information Systems (WIS) to use the MDE approach. Unlike the promise of MDD to enable a full automation of the whole development process, SOD-M aims to include designer-decisions during the development stages and consider it before executing each model transformation (customisation). This is done in order to reduce the complexity of model transformation resulting from the nature of the behavioural model specified at an early stage.

Modelling PIMs consists of constructing and mapping two models, namely, Extended Use Case (EUC), and Service Process models. The first model is used to capture the service functionalities of the system at a lower granularity, whereas the second model is a type of Activity Diagram used to represent processes. The mapping between these PIMs is defined from the EUC to the Service Process one. The Weaving Model is used as an annotation model to introduce design decisions during executing transformations. It is a container for the extra data absent from the behavioural metamodels, but required for transformation execution. For instance, it helps in mapping the order of executing *include* relationships in the case of having a use case with more than one *include-case* [120].

The annotated model is also used to establish and handle the links between model elements using ATLAS Model Weaver (AMW). Both it and a source model are then used as an input to the transformation to generate the target model. The annotation in the source model (e.g. EUC) is used to fill the missing data that is required to execute transformation. It also allows the customisation of model transformation [120].

The development process starts normally by transforming a high-level business model into a service composition model using several layers of transformation. As a common task in MIDAS, the rules of transformation are expressed with natural language at the initial stage, then it is formalised using graph transformation rules. ATL [33] is used to implement the formalised rules to be used in expressing the mapping between PIM and PSM within the framework [120].

2.5.4 UWE and MDUWE Metamodelling Approaches

Koch [117] proposes a model-driven engineering framework in UML-based Web Engineering (UWE), which is supported by a CASE tool integrated with Eclipse. In UWE [117], aspects of web applications are captured using different models (UML profiles[70]). These models are then integrated and transformed to produce a business logic code, web pages and configuration files. In MDA [89] terms, these models are CIMs (UWE profile use case models), PIMs, and PSMs (other UML profiles). The MagicUWE tool and any tool that supports UML modelling might be used to design web applications using UWE in order to provide a semi-automatic generation of web software. A UME approach employs various transformation languages such as ATL [33], QVT-P, and Java for implementing model transformations [117].

Kraus, Knapp, and Koch in [68], also proposed a complementary approach, called the Model-Driven UWE (MDUWE) approach, which is a based on transformations and metamodels. On the one hand, the transformation rules are implemented using the ATL transformation language [33]. A number of transformations are designed to obtain different aspects of the web systems, such as *Requriement2Content* for deriving the content model from the use cases.

On the other hand, the metamodel of UWE is structured into requirements, content, navigation, process, presentation and more packages or sub- metamodels using UML 2.0 profiles [70]. For example, the presentation metamodel defines elements that are used to specify the layout and user interface UI elements on the web page reflecting underplaying processes and navigation [68]. From that, it can be said that end-users still need to learn the UWE modelling language and some technical (designing) skill to model the different views of the web system.

2.5.5 Model-Driven Web Engineering (WebML)

WebML [124] is a domain-specific language for expressing the concepts and mechanisms of the domain of web engineering. It allows the specification of the conceptual model of web applications, such as data, service, navigation, and processes. It has its own formalism and abstraction levels of models. WebML [124] is supported by the WebRatio [4] tool for enabling automatic code generation for the J2EE platform [16].

The WebML [124] metamodel consists of several views, namely, the data model, the hypertext model, and the presentation model. The data model is used for presenting the data schema including entities, attributes and relationships. It is based on the standard ER model (Entity-Relationship). The hypertext model is regarded as a graph of linked pages that represent different information (view of the site) and also the navigation path between these pages. It establishes the overall structure of the domain as a collection of units, such as views, areas, pages, and content units. These units are connected together forming the WebML [124] hypertext model. The presentation is concerned with the visual representations and corresponding styles of the pages on screen [83].

Authors describe in [16] a transformation approach, integrated with Eclipse, for transforming WebML [124] to MDA [89] representations (MOF metamodelling layers [87]). The approach consists of three transformation stages, namely, Metamodel Generation, model Generation, and Higher Order Transformation. In Metamodel Generation stage, the mapping between the DSL metamodel and MOF M2 [87] is considered by mapping DTD and EBNF representations of the DSL metamodel into Ecore. The generation starts by parsing the concrete syntax of DTD and EBNF in order to derive the instance of the defined metamodel (DSL Injection). The next step is applying a model transformation that is implemented using ATL [33] between the DSL Metamodel and the Ecore one [16].

In the following stage, Model Generation, a transformation from an XML [122] representation of a WebML [124] project to an instance of the WebML [124] metamodel (generated at previous stage) is performed. Similar to the Metamodel Generation stage, this phase consists of an injection and transformation step. A Java injector is used in the DMS Injection step to convert the XML [122] document to an instant XML [122] metamodel. In the next step, an ATL [33] transformation is used to map an XML [122] and EBNF model to an Ecore model (MOF M1 [87]), which is an instance of the generated DSM metamodel [16].

2.5.6 Model driven, Existence dependency Relationship, Object oriented Development (MERODE)

MERODE is a model-driven engineering approach for developing enterprise systems. It adopts the MDA strategy, in which it aims at the creation of a complete platform-independent domain model (PIM) that is translated into a platform-specific one (PSM) and executable code [75, 21]. The method tackles the problem of UML semantics ambiguity by using a limited number of formally defined and semantically clear UML models (Class and State Machine), supplying them with an existence dependency graph and Object-Event Table [21]. Moreover, it uses the Process Diagrams from the BPMN as a behavioural business domain model to capture topmost concepts of the system functional requirements at the early stage of development (requirement/analysis) [21, 103].

MERODE has several modelling activities during the development process of IS, namely, data modelling, interaction modelling and life cycle modelling [103]. The data modelling step in MERODE is a process of creating structural aspects of business objects. A refined UML Class Diagram is used to construct the existence dependency graph using MERODE specific notation [21, 103]. This diagram has a key benefit to MERODE, in which it performs better consistency checking between the three MERODE diagrams rather than other modelling approaches [103].

Furthermore, the interaction modelling step in MERODE is achieved via the creation of the Object-Event Table (OET). It is a tabular representation that consists of all business objects with related business event types. The type of interaction can be either create (C) , modify (M) or end (E) . In addition, *life cycle modelling* is drawn according to each object in the domain to prevent events occurring in a wrong or random order during the life cycle of that object [103].

It is worth saying that the MERODE method is supported by an object-oriented tool (JMermaid) that consists of graphical editors for constructing and modifying specific system artifacts consistently. The tool aims to help software engineers to construct enterprise systems using a number of formal models [74]. The *JMermaid* adopts a template-based code generator mechanism that is able to successfully generate a multi-layer system that consists of a GUI, event handling (session beans) and persistence layer [104].

2.5.7 Reactive System Development Support (UML-RSDS)

The Reactive System Development Support (UML-RSDS) is a subset of UML with a precise semantics [71]. It aims at generating executable code from high-level specifications, in which it is supported by an MDE tool [71] to provide a rapid automatic creation of software systems from high-level standard UML artefacts, namely, *Class, Use case Diagram* and OCL [85] language.

The targeted multi-tier EIS applications, which are developed from UML-RSDS MDE tool, consist of five tiers, namely, Client, Presentation, Business, Integration and Resource tier. The Client tier consists of a number of HTML files as interfaces to the end-users. Moreover, the Presentation tier contains Java servlets classes and JSP files. In addition, the Business tier has some *session*, entity beans and value objects to represent the business services. Furthermore, the *Integration* and *Resource tier* consist of database interface, as well as web services, and the back-end data stores of the system respectively [71].

In UML-RSDS, two forms of use case can be constructed and used, EIS and General, using the standard mechanisms of structuring and composing use cases (via *extend, include,* inheritance and more). The EIS use case represents a simple straightforward operation (e.g. CRUD), whereas, the General ones are used to specify model transformations [71].

The UML-RSDS approach provides a formal specification approach for model transformation $[72]$. The transformation is defined in a hybrid style using *OCL constraints* that express the relationships between system models. In the approach, OCL is used in various places, such as, class invariants, class diagram constraints, operation preconditions and postconditions, use case pre-, postconditions and invariants, state invariant and transition guards [72, 71].

In order to generate more efficient code, the *constraints* are defined explicitly as *operations*, specified by pre- and postconditions, determining how a source element is translated to the target one in the target model [72]. The precondition of each rule identifies when this rule is applicable to which source element in the source model, whereas, the postcondition identifies what utilisations to elements and links should be constructed between the source and the target element [71].

2.5.8 Summary

This part of the literature review discussed some general modelling approaches in software engineering that are supported by MDE CASE tools. From the modelling perspective, we considered single-view modelling approaches and multi-view modelling approaches. It is noticed that many pure OO approaches employ, or adopt the UML [70] Class Diagram as a central artefact for capturing the structural view of systems, while, the UML [70] State Machine Diagram is used for capturing the behaviour side of the systems. On the other hand, in the integrated approaches, a modified version of Data Flow Diagram are used to capture system's behaviour and then mapped to OO models, such as the UML Class Diagram and other models.

Besides this, examples of DSL-based approaches are also considered in this part of the literature. We tried to identify what are essential and minimal views to be used at early phases of development. Furthermore, we tried to justify how complex the artefact (s) are, in both types of approaches, with respect to their suitability to business end-users. The following table (2.2) summarises drawbacks of the above approaches that are covered and improved in the proposed method in the thesis.

Table 2.2: Drawbacks of other MDE approaches

2.6 Overview of End-User Development (EUD) With Respect To MDE

The End-User Development (EUD) [106] is a development technique that aims to empower endusers, who have limited technical knowledge, to become involved in the process of designing and/or customising their systems to increase their productivity and satisfaction [107]. It also aims to translate accurately and comprehensively the informal description of domain problems to reduce the gap between what the exact user desires and what functionalities the implemented system has [30].

2.6.1 Component Based EBusiness Application Development and Deployment Shell (CBEADS)

Producing end-user tools for constructing web applications, such as *DEMIN* and *Mashups*, CBEADS [45] is a widely-known example of applying EUD [106] in the real-world to tackle the problem of the lack of web developing skills for non-programmers [97]. This kind of tool considers the business-users perspectives or mental model [97], encodes developers knowledge as rules [92] and enables users to easily tailor software to meet their individual needs [107].

Ginige and De Silva, in [45], have introduced a meta-modelling approach for enabling endusers to be involved in the continuous development process, side by side with developers, end users with little technical knowledge. It also aims to enable them to employ effectively the metamodels in such as customisable environment [45]. The concepts in the meta-model are similar to those in the UWE approach (UWE, 2011). However, all common aspects are grouped at a separate level of abstraction for describing information systems, namely, Shell, Application and Function level [30]. These models are embedded into a component-based Shell introduced as developers templates for end-users to instantiate a meta-model instance and populate it at one or more levels. The *CBEADS* and other related (SMART) tools are used to generate business objects as well as functions, users interfaces and SQL queries [45].

2.6.2 Model-Driven Development for End-User development (MDD4EU)

This approach provides a mechanism to allow end-users to collaborate with developers at the early development stage, rather than transform end-users description of their needs into developers. It is supported by a modelling language (DSL) called Pervasive Modeling Language (PervML). PervML is used on the construction of pervasive services in the context of smart home systems [91].

MDD4EU is mastered by developers, in which they determine what knowledge (properties) end-users are allowed to supply to models during the development process. After that, developers also determine an appropriate language that is familiar to end-users for enabling them to supply and edit models using the language. At the end of the story, a model that holds user preferences is created with a degree of quality [91].

In order to achieve this, the MDD4EU approach is introduced to end-users via a toolkit that is specially designed in such a way that it suits end-users. The toolkit is designed to be domain-independent, in which it has an interface that might be reusable in other domains when adopting the MDD4EU [91].

2.6.3 Summary

This part of the literature discussed some model-driven engineering approaches that allow enduser participation in the development process. Both approaches provide mechanisms for allowing end-user to participate and work with developers. In contrast, the thesis aims at going beyond this and allow business end-user to lead and master the development process of information systems using less technical knowledge. The movement from this business knowledge to more technical one is performed and managed by transformation rules.

2.7 Outlook on the Chapter

This chapter discussed the related literature, including a brief comparison between alternative methodologies such as, AOSD, MDA, xUML and Software Factory. From this, it is possible to realised where the proposed approach fits, in contrast to other related model-based approaches, which is a multi-view MDE framework for developing enterprise information systems. The approach employs a user-friendly and semantically clear UML-like modelling language to describe system specifications at a higher level of abstraction than existing UML approaches.

Furthermore, the second part of the literature investigated model transformations, including the languages, paradigms, composition techniques for model transformations and code generation strategies. This formulated the adopted model transformation mechanism and language in the proposed MDE method in the thesis, which is a hybrid model transformation style, encoded using *Java* programming language.

Additionally, a number of model-driven engineering approaches are covered in the rest of this chapter. The discussion focused, mainly, on the complexity of metamodels, modelling languages and notation, used for defining and expressing domain concepts, and the ability to generate complete executable code from abstracted models. Core concepts of the proposed modelling language are defined using a metamodel rather than an UML profile. Moreover, endusers participation, in some approaches, is highlighted and compared in order to introduce a development process led by business end-users.

3 Framework Overview and Analysis

"Make everything as simple as possible, but not simpler" Albert Einstein

3.1 Context

The chapter presents a general overview of the proposed solution showing how it supports Business-user Model-Driven Engineering. It introduces, in brief, the notion of a simplified user-friendly modelling language, which we call the Micro Modelling Language (μML) ; and, it discusses briefly the purposes and engineering activities in each development phase of the proposed Business-User Information-Led Development (BUILD) framework.

In each stage, the chapter motivates the idea of adopting one or more μ ML model(s), and shows how this suits the conceptual (natural) thinking of business users. It also demonstrates how the μ ML models collaborate for capturing critical aspects of the system, which lead to significant design decisions. Each model is exemplified by a small portion of a real-world business process.

The adopted structure of model transformations throughout the development lifecycle is also reviewed, including modularisation, reusability and composition. The chapter also formalises the basic foundation of the language metamodel that will be used frequently in the following chapters.

3.2 BUILD: Business-User Information-Led Development

The proposed Business-User Information-Led Development (BUILD) framework is a model driven engineering approach to enterprise information systems development. It aims to derive detailed implementations of information system layers from high-level business-user specifications using a lightweight modelling language called the Micro Modelling Language (μML) . The inspiration for this methodology is the Reusable Modelling Design Language (ReMoDeL) approach[101], which is based on a multi-layered, forward-transformation strategy that turns

multiple abstract representations of the requirements into concrete design models and generated code.

As discussed in earlier chapters, end-users understand their business well in a whole, but it is difficult for them to articulate, systematically, the software system that they require using their limited technical skills. BUILD encourages end-users to employ their business knowledge to construct a number of user-friendly μ ML models, and let the transformation rules evolve an executable prototype system. This strategy avoids the mistake of asking users to construct complicated models that are full of technical design detail and formal constraints.

In contrast, end-users express business processes, entities and the impact of each process on one or more business entities, using three simple requirements models with semantically clear notation (Task, Impact and Information Model). The following figure (Figure 3.1) presents a wider picture of BUILD using a traditional flowchart. The development process within the framework is divided into four main phases (illustrated by horizontal dashed lines), namely, Requirements Sketching, Analysis, Design and Code Generation (Section 3.2.2). The development progress, from one stage to another, is handled and proceeds by model transformation rules (steps) that are responsible for the evolution of the model artefacts representing the system, from initial requirements through to generated executable code.

In Figure 3.1, the μ ML models used in each phase of the BUILD method are depicted as Document nodes in the flowchart notation. These models are sorted as ASTs in memory and XML files to express various system views to be used during the model transformation steps. Additionally, all forward translations from a single aspect (model) to a new one, merging two aspects into a new one and performing in-place modification of a particular aspect, are depicted in Figure 3.1 using the Arrow flowchart notation. External human participation, which is needed to construct models or supply additional data, is also depicted in the figure using manual input notation.

3.2.1 What is Micro Modelling Language (μML) ?

The Micro Modelling Language (μML) is a user-friendly modelling language that is used to express basic structural and behavioural aspects of enterprise information systems. It has graphical UML-like representations that correspond to underlying XML parse trees. It aims to simplify the modelling activity using simpler notation with cleaner semantics than existing UML-based, or Model Driven Architecture (MDA) approaches. It accelerates the development process by using effectivily business-users' knowledge of their desired system.

It can be argued that the simplicity and clarity are achieved in the proposed μ ML when each notation has the same meaning in each model, but in appropriate way that suits the context of each model. For instance, μ ML unifies the visual notation of the generalisation and compostion relationship between elements in each model it appears in. This issue is discussed in more detail in chapters 4, 5 and 6, where the notation and its semantics is described.

In BUILD, μ ML models are used to represent various information system views. There are three major groups of models, with each group being associated with a specific development stage in order to capture the essential characteristics of the system. The models in the top group are employed by model transformation to produce further detail and new aspects (models) of the following group to be used in the next stage of development. μ ML concepts, notations and their formal foundation are discussed in detail in chapters 4, 5 and 6, according to the BUILD

Figure 3.1: BUILD. A wider picture

stage that the model belongs to.

3.2.1.1 The Conceptual μ ML Metamodel Hierarchy

Conceptually, the μ ML main concepts are organised in a hierarchical structure of model nodes (elements), which form a core metamodel for the language. The following Figure (3.2) illustrates the main generalisation and composition relationships between μ ML elements. It consists of higher level concepts of the metamodel that express the core elements of the language. These concepts are: Element, Node, Edge and Model.

The μ ML metamodel includes a number of intermediate nodes that group some terminal nodes based on their common features. For instance, the kinds of Arcs subdivide into the Flow, Transition and Participation arcs. Additionally, the kinds of Relationships subdivide into: Association, Generalisation, Composition and Dependency relationships. The metaclass Model consists of a collection of the kinds Element.

In μ ML model(s), all *Arcs* and *Relationships* are defined as instances of one of their descendant (metaclasses). Each of the leaf (terminal) nodes, appears in the metamodel (Figure 3.2), corresponds to a concept in one or some μ ML model(s), which has a specific meaning and

Figure 3.2: The Metamodel for μ ML

a graphical notation. Furthermore, there is an element in the underlying representation (AST) of the models to represent each terminal concept.

3.2.2 Development Phases in BUILD

The overall structure of BUILD can be divided into four main phases: Requirements Sketching, Analysis, Design and Code Generation. Each phase in BUILD focuses on a variety of abstract system views to present the system on several levels of detail and encapsulates various model transformation activities that occur during the process of software development.

3.2.2.1 Requirement Sketching Phase

In this context, requirements engineering can be introduced as common practice in the software development lifecycle. It refers to the elicitation of stakeholders' desires and all other activities involved in discovering the requirements for a SE process. This stage involves the initiation of communication and collaboration between end-users and the developer team in order to map system requirements and objectives onto a collection of formally designed software characteristics and functions.

Engaging different stakeholders in the initial phase of the development process is considered critical in all traditional software development approaches, such as Waterfall Model, V-Model, Incremental Model and Spiral Model [94]. The functional requirements have to be explained clearly and transferred precisely from the business users to the developer team. This happens via holding a number of formal meetings to discuss requirements and the ability to accommodate them within the system. In the end, all details are stated textually in the contract.

By contrast, the current trend in new approaches to software development, e.g. Model Driven Software Development (MDSD) and Domain Specific Language (DSL), is to assign the task of expressing the functional requirements and system structure to skilful designers. These designers must have enough knowledge to deal with complex modelling language and models rich in detail. They are responsible for modelling the system to meet the functional requirements provided by the users. From there, the initial step of passing the end-users requirements to the designers must also be taken.

In BUILD, the business-user engagement is different from that found in traditional approaches. During the Requirement Sketching Phase, the business analyst, who is aware of the exact requirements of the organisation, establishes formal agreements by describing the functional and structural requirements of an information system, in models, using clean and simple modelling language. The adopted μ ML notations are lightweight and tailored to capture, visually, the intellectual logical thinking of end-users about their business. The artifacts produced are evolved directly through a chain of model transformations.

Practically, business-users describe the real world business data in terms of entities and relationships, making the entities organised, structured and easy to understand. Moreover, they describe a skeleton structure of the system that contains business tasks, goals and real world interactions with stakeholders. Additionally, users also express how the task affects the business entities in terms of the impact on the system object during the execution.

A number of system aspects must be prepared in order to progress to the second stage (*Analysis*). The proposed μ ML models, Task, Impact and Information Model, can adequately cover the end-users specifications that act as a starting point for the first collection of transformation steps to satisfy the requirements of the analysis phase.

Task Model

The Task Model is a **manually-constructed** structural model that describes the structure of all significant enterprise activities. It typically captures the wider context of the business, including interactions with external stakeholders. By using the conceptual (natural) thinking of business-users to express the functional requirements, in terms of the business process and stakeholder interactions (Figure 3.3), the Task Model can be constructed naturally in a straightforward way, stating the business tasks and the stakeholders involved in them. The notation and concepts of the Task Model are presented and discussed in detail in Chapter 4.

Figure 3.3: Task Model. Library Circulation System example

Figure 3.3 above illustrates a Task Model example of a small portion of Library Circulation System. It indicates that the Circulation consists of two business tasks: Issue Loan and Discharge Loan, modelled using a part-of relationship. Furthermore, a *Borrower* and *Service* Reader participate by supplying the system with some information and/or receiving back some responses from it.

Impact Model

The *Impact Model* is a kind of behavioural model, **manually-constructed** by business users, that describes the impact of business tasks on the internal objects portion of the system data. Exploiting the end-users awareness and understanding of the business to sketch the interconnection between tasks and objects enables the capturing of the internal behaviour of these tasks. This is extracted by model transformations to estimate the control flow between tasks (the order of execution). The following figure demonstrates an example of the Impact Model of a Library Circulation System. The notation is discussed and formalised in Chapter 4.

Figure 3.4: Impact Model. Library Circulation System example

Figure 3.4 above demonstrates an Impact Model example of a small portion of Library Circulation System. It shows the internal behaviour of the Circulation task components: Issue Loan and *Discharge Loan*, collaborating with the logical business objects: *Copy*, *Loan* and *Borrower*. The figure indicates that the *Issue Loan* task reads *Copy* and *Borrower* objects to create a *Loan*, whereas the *Discharge Loan* reads only a *Copy* to destroy the related *Loan*.

Information Model

The Information Model is a **manually-constructed** structural (abstracted) representation that describes real-world business aspects and the relationships that might exist between them. This model is obtained through two different methods: a partial model is derived (automatically) from a pre-defined Impact Model or a complete one is provided (manually) by end-users. The first method is considered a simpler approach for Information Model construction because the framework takes responsibility for manufacturing the required relationships between entities, and leaves the task of specifying their end-role multiplicities to the end-users.

The structural relationships between objects in the Information Model can automatically be extracted from the Impact Model in the first approach. A number of mapping rules are applied to achieve this shift between these system views. This leads to the construction of Information Model Aggregation and Association relationships between the captured business entities. Then the properties of the entities (attributes) and the end-role multiplicities are added as part of a separate design activity.

On the other hand, in the second method, the end-users are in charge of maintaining consistency during the construction of each model to avoid any clash of inter-relationships between objects. They must model manually all business entities and their relationships. The objects that appear in the Information Model might be physical objects and documents, logical information about stakeholders or other physical items used in the business. The inter-relationships between these objects are presented in an un-normalised form. It is worthwhile emphasising that the current implementation of BUILD supports this construction approach.

Figure 3.5: Information Model. Library System example

The above figure (Figure 3.5)shows an example of the *Information Model* of a *Library* Circulation System. It illustrates a conceptual whole/parts relationship (Composition) between three objects: Loan, Copy and Borrower. The Loan object represents the whole side of the relationship, whereas Copy and Borrower are its parts. The notation is discussed and formalised in Chapter 4.

3.2.2.2 Analysis Phase

Several intermediate models that contain extended detail about the required system are constructed in the Analysis Phase. This stage is inspired by the early Structured Systems Analysis and Design Methodology [43]. Our method attempts to specify the requirements in terms of initial business tasks that need to be decomposed into a number of atomic ones, based on a functional decomposition technique and the use of a data flow diagram.

As the outcome of this stage, a number of artifacts are produced in an attempt to reveal more detailed aspects of the system. The internal behaviour of each business task and the state of the system are clearly identified to the designers in terms of data flow (detailed DataFlow Model) and state-transition (*State Model*). In addition to this, the business entities captured previously during the Requirement Sketching Phase are refined and re-described in the notion of the Data (Dependency) Model. All artefacts at this stage are **automatically-generated** by rules, except the initial DFD model that needs end-users engagement via annotating data on flows.

It is worthwhile noting that during the shift from the Requirement Phase, the number of complex transformations varies for each model, for instance, the creation of the detailed DataFlow requires two forward translating steps; the initial $DataFlow$ (DF) is created first and then the detailed DF is derived from it. On the other hand, the compilation of the *State Model* construction requires one forward translating step from the detailed DF model and one in-place modification for generating states of failure business scenarios. In contrast, the Data Dependency Model is constructed directly using a single translating step from the *Information* and/or *Impact* $Model(s)$.

Indeed, the Model Transformation Framework takes Task and Impact Models as source models and then applies translation rules to fold their concepts and generate the initial DataFlow Model. For each task that appears in both source models, new components in the target model arise after combining the external interactions provided/received by actors and the internal ones

with objects that lie within a logical boundary representing this task. The resulting boundaries introduce a number of independent event-driven business process transactions of the system.

To exemplify this, new Input/Output subtasks are generated after extracting the system interactions with the actors. Furthermore, CRUD subtasks are produced after translating the internal interactions with the system objects (back-end entities). Both kinds of subtasks are linked and grouped within a logical boundary of the original business task. It is important to mention that the boundaries in the initial *DataFlow* have no knowledge about the data on flows and there are no interconnections between subtasks. These missing parts are considered in the next transformation step.

The existence of data on flows is fundamental to the approach in order to estimate the possible order of task execution. Thus, the manual engagement of business end-users is required to specify the kind of data flow on the initial $DataFlow$ artifact. As a consequence of visualising the DataFlow Model using μ ML notations, users can add data to related flows on the model in a straightforward way. These extra details are used simultaneously with the type of CRUD operation assigned to tasks during the task decomposition step in order to construct the detailed DataFlow Model.

Accordingly, boundaries in the detailed model express the data flow between linked subtasks. This version of *DataFlow* is used as a source in a second translating step in order to generate a State Model, after prioritising the atomic tasks in the previous translating step, whereas the Data (Dependency) Model is derived directly from the predefined Information Model. It is worth emphasising that the *Analysis Phase* ends after decomposing all the business tasks that require a number of atomic actions into tasks with a single action that are presented with a logical order of execution.

Data (Dependency) Model

The Data Model is an intermediate view, **automatically-generated** from a pre-defined Information Model, that describes the dependency of the logical data in the system, and supports the development to a point where a logical database schema may be generated. The model consists of logical objects linked by a number of dependency relationships, representing the direction of data dependency.

Figure 3.6: Generated Data Model. Library Circulation System example

Figure 3.6 illustrates a part of a Data (Dependency) Model of a Library Circulation System. It demonstrates dependency relationships between three objects: Loan, Copy and Borrower. The Loan object depends on both Copy and Borrower.

DataFlow Model

The DataFlow Model is an **automatically-generated** intermediate model that describes how information data is transformed in tasks, flows via system components and is manipulated into data sources at a particular level of abstraction. The model consists of a number of entities

and participants (agents) that are connected to some tasks via a variety of flows. It provides a detailed description of the business tasks and their aggregated subtasks that supply and perform operations on data in terms of data flows.

It is worth mentioning that the creation of a complete $DataFlow Model$ requires two independent mapping steps. The first one performs automatic merging and gathers concepts from both Task and Impact Models to derive a number of business tasks that are connected by some (typed) flows with system stakeholders and internal objects. The second step starts by a manual engagement from end-users who annotate the generated data flow model by data on flows. Then, the automatic task decomposition step is applied to the annotated intermediate DFD to produce the complete DFD. The detailed description of these construction steps is presented later in chapter 7.

Figure 3.7: Generated DFD Model. Library Circulation System example

Figure 3.7 above shows a portion of a DataFlow Model of a Library Circulation System. It illustrates the contents of the $Sign\ In$ business task. It can clearly be seen that the task consists of only one task, Validate ID, that is connected to three flows: Input, Read and Output. A stakeholder, User, interacts with the system by supplying login details to be checked via the task against the data stored in the login entity. The *Validate ID* then responds back to the user with a suitable message.

Screen State Model

The Screen State Model is an **automatically-generated** intermediate representation, derived from a pre-generated detailed DataFlow Model, to demonstrate, in an abstract way, the behaviour of the system in terms of screens and their navigations. It is mainly based on the notion of a State Machine[116] in which each screen appears as a state and associated navigations appear as transitions.

DataFlow Boundaries and Tasks are mapped into boundaries that contain States (each represents a systems state after executing an action) and Transitions, which trigger their actions, respectively. Any flow connected to an object (either source or target) is translated into a Ready state, and any $Input/Output$ flow connected to an actor is translated into a *Waiting* state. The Waiting state indicates that the system is waiting for some external activity related to stakeholders, such as input or output information. On the other hand, the Ready state indicates that the system is ready to execute a back-end operation (e.g. a database action), such as an update, insert or delete of a record.

Figure 3.8: Generated State Model. Library System (Issue Loan) example

The above figure (Figure 3.8) demonstrates the content of the *Issue Loan* business process. It consists of two Waiting states for inputting copy and borrower detail and one Ready state for (creating) inserting a new Loan record.

It is worth mentioning that the creation of a complete State Model requires an in-place model modification step, which it considered to be a separate transformation step. A number of Error states are generated to handle failure cases of the business task, such as invalid/null inputs or database connection failure. The following figure (Figure. 3.9) illustrates a complete State Model of the Issue Loan task after applying the in-place modification step.

Figure 3.9: Generated State Model. Library System (Issue Loan) example

3.2.2.3 Design Phase

The *Design Phase* is considered to be a stage that is led by model transformations to construct a number of low-level design models. This phase aims to reach an adequate level of detail to enable a full code generation of an implemented solution. This involves the creation of the technical Database and Query (DBQ), Graphical User Interface (GUI) and Code Model. All constructed models are **automatically-generated**, platform-independent and are ready for the Automatic Code Generation Phase to produce a platform-specific code.

The Database and Query Model is derived by translating business entities that appear in the pre-constructed Data (Dependency) Model and by consulting the type of CRUD effects for each entity in the detailed $DataFlow$ of the previous stage. As an outcome, a complete definition of relational database tables with clear information about their keys is presented in the model. Moreover, a number of predefined stored procedures and triggers associated with the tables also appear, including their implementations. The resulting artifact is generic in that it can be used as an entry point for generating executable relational database schemas, such as $MySQL$ and Oracle.

Additionally, the *Screen State Model* is used as a source in the translating step to construct a *GUI Model*. The produced artifact consists of multiple connected event-driven screens that perform the business, triggered by the system end-users. Furthermore, states of the possible failure scenarios are also translated into screens linked to the related screen in the successful scenarios or to the termination of the process.

The resulting GUI Model is considered to be platform-independent and complete enough to be employed in the next Code Generation Phase. For each generated screen, minimal component features, e.g. size, name and text, are inferred in the model as a consequence of crosschecking with object specifications and types in the *Data Dependency Model*. Moreover, tracing the type of CRUD operation and its carried data indicates the most appropriate type of generated GUI control. For example, in the case of Insert action, the possible type of generated GUI control on the screen is TextField.

Apart from this, there is a room for a separate folding transformation step that is applied to a number of analysis models: Task, DataFlow, and State Model, to construct the (System) Code Model. This step is independent and not part of the main IS development process (optional), as the current version of BUILD aims at generating 2-tier applications. This model is optionally used to represent the *business logic* layer in the 3-tier applications, because the business logic is by default translated into stored procedure logic in the database schema directly. The generated *Code Model* is also crosschecked with entities and attributes in the *Data Model* in order to maintain consistency in the object-relational mapping process when required.

Database and Query (DBQ) Model

The Database and Query Model is an **automatically-generated** lowest level representation that describes the normalised design of the logical schema in a relational database, as opposed to a conceptual information model. A model is considered platform-independent that includes the generic data schema definitions, query expressions, and logical representation of tables and fields. The main objective is that a comprehensive level of abstraction for representing all generic specifications, required for any database generation, is achieved and expressed at this level of detail. This low-level model is used as a source model in the code generation approach. Figure (3.10) illustrates an example of the generated DBQ model, taken from the Library System example.

Figure 3.10: The Generated DBQ Model. Library System example

Graphical User Interface (GUI) Model

The Graphical User Interface Specification Model is an **automatically-generated** lowest level textual representation AST (XML format) that describes the structure and the components of information system screens, in terms of windows with some controls and events. It is an abstracted platform-independent model in which there is no implementation detail presented in the model. This model does not have a graphical representation in the current version of the proposed μ ML ver.1.0.

Code Model

The Code Model is an **automatically-generated** lowest level AST (XML) representation that describes the abstract specification for Object-Oriented (OO) code of the system in terms of business entities and tasks (classes), their methods, attributes, and expressions. It aims at representing a comprehensive business logic layer that determines how to manipulate the data in a separate layer, as part of an (optional) three-tier architecture. In the Current version of BUILD, the model can be constructed independently using a separate folding step, as it is not a mandatory part of the main development process of IS. Similar to the GUI Model, this model does not have a graphical representation in the current version of the proposed μ ML ver.1.0.

3.2.3 Model Transformation Strategy in BUILD

The transformation strategy in BUILD consists of a number of independent translation steps, where each step contains a collection of *Translator* agents. It takes the given knowledge from one level and evolves it to a new more detailed level. As an outcome, several intermediate and lower level design artifacts are constructed based on end-user specifications. Generally, the strategy attempts to reach an adequate level of detail that enables full code generation, which is a completely independent phase.

The transformation approach relies on a *hybrid* (*imperative - declarative*) style. Each rule is designed separately, in which it does not know how other rules work or how they are implemented. There are dependency relationships between rules. Each rule fires selectively, if its input matches a conditional guard, and may employ further rules to execute part of the transformation. The design of the rule-tree typically follows the compositional structure of the model. While rules are always invoked in some given order on collections of artefacts, this order is not significant, in the sense that rule-dependency is the only constraint, and multiple firings of the same rule on the same inputs are idempotent. The rules of transformation are directly encoded in Java classes to define the map between input elements and output ones. Each set of classes represents a Javatranslator that is responsible for orchestrating the translation rules in order to derive target models from source ones.

At each development stage, these transformations are capable of initiating an automatic shift from the current phase to the next one, introducing more detailed artifacts without the direct engagement of business users. However, on a specific layer of the *Analysis Phase*, the designer supplies manual data flow detail to be considered in the next step of the transformation. It is worth mentioning that the BUILD framework has two model-to-model transformation layers: Requirement-to-Analysis and Analysis-to-Design, whereas the Code Generation Phase is a completely independent layer and is totally automatic.

The *Translator* agents resemble the *Visitor* and *Composite design patterns*. All models can be regarded as a kind of content containing other similarly-structured content. This feature makes it possible to apply transformations as a kind of *Composite design pattern*, in which each rule delegates to other rules to continue the transformation, recursively. In addition, each translating agent has knowledge (rules) embedded within it to be applied to a source model.

Additionally, the translators behave like the *Visitor pattern* in the way they traverse a model. Each translator is responsible for a scope of the source model (parse tree). It focuses only on a particular part and applies translation rules to generate a part of the target model. The agent names are based on the source and target mappings. The following list addresses the translation steps within the Requirement-to-Analysis transformation step:

- Information to Data Dependency Model: This one-to-one translator derives the dependency relationships, between business entities, from binary association, generalisation, and aggregation relationships that appear in the sketched Information Model.
- Task and Impact to DataFlow Model: This two-to-one translator merges concepts from given Task and Impact Models and produces output elements in the (initial) DataFlow Model.
- *Initial DataFlow to Detailed DataFlow Model:* Based on the business-user supplying extra information about the nature of the data on each flow, this transformation takes the initial $DataFlow$ with the additional information, provided by users, and construct the detailed DataFlow Model that contains the interconnections between atomic tasks.
- Task, DataFlow and Data Dependency to State Model: A merge of DataFlow and Data Dependency Model and a one-to-one translation from Task Model are considered together to produce to produce a Screen State Model.

As an outcome of the transformations described above, the artefacts of the Analysis phase are completely constructed and ready to act as source models for the next transformation step (Analysis-to-Design). This step of transformation adopts the identical structure to the previous one, in which it consists of a number of Java-translators for producing richer detail artefacts as follows:

- State to GUI Model: This translator is responsible for constructing GUI detail from the generated State Model, including actions and widgets controls based on the screen type.
- Data Dependency to Database and Query Model: The translator employs the dependency information, expressed in the source model to manufacture foreign keys, for the final form the translated business entities (normalised database tables). For each table, all fields are formally typed and specified in the way that it is interoperable by the relevant code generator to produce complete table definition syntax in SQL.
- DataFlow to Database and Query Model: The translator considers tasks attached to objects in the source model in order to create equivalent stored procedures.

As a result of the above transformations, all required design models are constructed and ready for the *Code Generation* phase. This stage is presented with further detail in the next section.

3.2.4 Code Generation in BUILD

The Code Generation Phase, the final stage of development, performs model-to-code generation in a particular running environment. It consists of a collection of generators that take the responsibility for generating executable code from detailed models. The overall architecture of the Code Generation Phase in BUILD also exemplifies the Visitor and Composite design patterns, in that it consists of a number of independent Java-generators. These generators are designed to support code generation for different specific platforms.

Unlike the overall structure of the *Java-translators* in the previous model-to-model transformation stage, we assign a Java-generator to produce the implementation for each information system tier: *Presentation, Business processes* and *Data*. These generators are:

- Database Schema Generator: The Database Schema Generator takes concepts from the design DBQ Model to generate executable dump files that contain SQL scripts. In this thesis, MySQL Java-generator is presented as a proof of concept, Chapter 7.
- GUI Generator: This generator creates Windows application screens with their widgets to visualise the designed system. A Java-generator for Java Swing Application with JDBC is presented as a proof of concept, Chapter 7.
- OOP Code (System) Generator (Optional): This generator is responsible for constructing a separate business logic layer that consists of a number of business entity classes, if required, Chapter 7.

3.3 Foundation of μ ML

It is worth emphasising that, in formalising μ ML, we are not looking for a degree of automated analysis and verification, all we are seeking is to add a formal flavour to the semantics of μ ML concepts and notation. This section presents an approach to formalise μ ML using standard First-Order Predicate Logic (FOPL) with extensions for subtyping $(\langle : \rangle)$, equality $(=)$, membership $(∈)$, exclusive disjunction $(oplus)$ and uniqueness quantification $(!)$.

According to previous work in this field [14, 58, 82, 99, 125], it can be noted that using FOPL is a widely adopted approach for formalising the constructs of UML diagrams and other DSLs. It mainly used for describing the language semantics and the interrelationships between its elements via mapping the concepts to a number of FOPL predicates and symbols.

The spirit behind the proposed strategy for formally specifying the semantics of μ ML concepts is similar to approaches introduced in [99, 125] particularly. In [99], for instance, a formal approach is presented to define the descriptive semantics of modelling languages by mapping metamodels concepts to a number of FOPL sentences, representing the core concepts of the language via predicates and functions. In the same context, [125] proposes mappings from UML models and metamodels to FOPL statements for describing their semantics. As an outcome, models creation policies and mapping rules can be expressed formally using FOPL.

In Comparison with OCL

OCL is based on FOPL and set theory, but it uses a syntax similar to programming languages and closely related to the syntax of UML. OCL is intended to be a light-weight formal specification language for annotation purposes e.g. invariants, guards, pre- and post-conditions for methods or UML diagrams, rather than to be used as the basis for extensive automated behaviour analysis [78].
3.3.1 Preliminary Principles

This thesis represents the basic concepts of μ ML through predicate logic. The predicates are derived from the concepts appear in the μ ML metamodel, such as, *Model*, *Node*, and *Arc*. Thus, the description of each concept in μ ML models is provided with respect to its equivalent concept that appears in the metamodel. For instance, given m is a *Node*, a unary predicate $Node(x)$ denotes x is an instance of Node.

Similarly at the model level, concepts like Task and Object are both considered Nodes at the metamodel level. By applying the notion of subtyping, we can express this in our logic as: given $Task(x) \rightarrow Node(x), Object(x) \rightarrow Node(x)$ denotes that if x is a Task then it is also a Node, this can be expressed in our logic as $(Task \lt;: Node)$, or if x is an Object then it is also a Node, this can be written in logic as $(Object \leq Node)$.

Two syntactic sugar declarations are used for these predicates. The unary predicate for assigning the type, e.g $Model(x)$, can be shorten and re-expressed as $x : Model$. In addition, the symbol (,) is also used to replace the logical And (\wedge) operator some occasions. For example, the statement $n1, n2$: Node is equivalent to $(n1 : Node) \wedge (n2 : Node)$.

On some occasions, we need to reason about the logical value of properties. For example, one of the main features of μ ML language is that *each Element in the Model has a unique Iden*tifer. This means the distinction of Elements is based on the value of their Identifiers. From that, if we want to say that *Node (a) and Node (b), in such a Model (m), are distinct*, we need to prove logically that their *Identifiers* are not the same. As a consequence, the need for *equality* $(=,\neq)$ emerges to add this expressive feature to our logic. In order to add *equality* to our logic, three axioms must be declared:

- Reflexivity: $\forall x \bullet (x = x)$ (a variable x is equal to itself).
- Function substitution: $\forall x, y \bullet (x = y) \rightarrow (f(..., x,...) = f(..., y,...))$ (if you substitute y for x in function f , the results are equal).
- Formula substitution: $\forall x, y \bullet (x = y) \rightarrow (p(..., x, ...) \rightarrow p(..., y, ...)$ (if you substitute y for x in the predicate p , the truth-value is the same.

From that, other properties of equality can be expressed as follows:

- Commutativity: $x = y \rightarrow y = x$.
- Transitivity: $(x = y) \land (y = z) \rightarrow x = z$.

In our context, based on the axioms mentioned above, given $id1$ and $id2$ are Identifiers, $id1 = id2$ denotes that id1 and id2 are equal. Similarly the statement id1 $\neq id2$ denotes that that id1 and id2 are different (not equal). Moreover, in order to say that y has the name of x, a simple logic function $nameOf(x, y)$ is declared.

Besides this, a binary predicate is used to express the notion of membership in our approach. For instance, given $Model(x)$ and $Node(y)$, $Member(x, y)$ denotes that element y is a member of the model x. This predicate can be re-written using (ϵ_m) symbol as $y \epsilon_m x$. Therefore, the statement says that the node y is a member of the model x can be expressed as: $x : Model, y :$ Node • $(y \in_m x)$

In the same context, a binary predicate is used to express the notion of property. For instance, given $Node(x)$ and $Property(y)$, $PropertyOf(x, y)$ denotes that property y belongs to the node x. This predicate can be re-written using (ϵ_p) symbol as $y \epsilon_p$ x. Therefore, the statement says that y is a property of the node x can be expressed as: $x : Node, y : Property \bullet$ $(y \in_p x)$.

Furthermore, another binary predicate is used to express the notion of widget. For instance, given $Node(x)$ and $Window(y)$, $WindowtIn(x, y)$ denotes that widget y is a control in node x. The node x must be a GUI window. This predicate can be re-written using (ϵ_c) symbol as $y \in_c x$. Therefore, the statement says that y is a GUI control in the node x can be expressed as: $x : Node, y : Widget \bullet (y \in_c x).$

To abbreviate this, the symbol \in without any subscript character is used for indicating Member, WidgetIn and PropertyOf predicates. The interpretation of \in is obtained based on the types of elements appearing in the logic statement.

Additionally, in our formalisation appraoch, there is a need for the use of an exclusive disjunction operator (\oplus) operation to emphisize that only one of the operand predicates is true but not both. For instance, any Element in the Model can be either a Node or an Edge. This cannot be expressed adequately using the direct logical (\vee) , in which it produces true output when one or both operand predicates are true. The statement $Node(x) \oplus Edge(x)$ can be bootstrapped using logical And (∧), Or (∨), and Negation (¬) operations as: (Node(x) \wedge $\neg Edge(x)$ \lor $(\neg Node(x) \land Edge(x)).$

Moreover, the uniqueness quantifier (!) is used, here, to indicate that exactly **one and** only one unique property exists in a certain *Element*. For instance, the phrase a *Node has a* unique Identifier cannot be expressed without (!) to emphisize that the Identifier is unique and only one. The uniqueness quantifier $\exists ! \text{ id} : Identifier \bullet Id(n, id)$ appears in the statement: $\forall n : Node \bullet (\exists! id : Identifier \bullet Id(m, id))$ can be bootstrapped using basic existential (\exists) and universal (\forall) quantifiers with the logical And (\wedge) , Or (\vee) , and Negation (\neg) operations as: $\exists id1 \bullet (Id(n, id1) \wedge \neg \exists id2 \bullet (Id(n, id2) \wedge id2 \neq id1)).$

Regarding the transformation rule definition, it is required to consider, for each rule, what is the next-level-down translation issue in order to specify exactly what goes where. Therfore, it is assumed that a general polymorphic translation function tr is defined: \forall Source, Target • tr: Source \rightarrow Target. Each of subsequent translation rules, presented later in Chapter 8, is a particular type-instantiation of this function, with a different overloaded rule. This means that the function $tr(x)$ can be used in a rule-body to refer to the translation of the next bit of sub-structure x. Section 8.3.2.1 in Chapter 8 is an example of translating $DEntity \rightarrow Table$.

3.4 Specifications of the μ ML Metamodel

This section introduces a set of common policies that are applicable to all μ ML model elements. The presented laws are used to specify the concepts of μ ML metamodel using FOPL. In order to construct valid instance models, the following laws must not be violated. The following table (3.1) summaries all predicates used in describing μ ML Metamodel.

Law 1. Each model m has one identifier id.

$$
\forall m: Model
$$

$$
\longrightarrow
$$

$$
\exists! id: Identifier \bullet Id(m, id)
$$
 (3.1)

Law 2. Each model m has some nodes.

$$
\forall m: Model
$$
\n
$$
\longrightarrow
$$
\n
$$
\exists n: Node \bullet (n \in m)
$$
\n(3.2)

Law 3. Each node n in a model m has one unique identifier id .

$$
\forall m: Model, n: Node
$$
\n• $(n \in_m m)$ \n
$$
\longrightarrow
$$
\n
$$
\exists! \ id: Identity \bullet Id(n, m, id)
$$
\n(3.3)

Law 4. The node can be either a task, goal, object, entity, actor, state, or window.

$$
\forall n: Node
$$
\n
$$
\rightarrow
$$
\n
$$
Task(n) \oplus Object(n) \oplus Entity(n) \oplus Actor(n)
$$
\n
$$
\oplus Goal(n) \oplus State(n) \oplus Window(n)
$$
\n(3.4)

Law 5. Each model m has some edges.

$$
\forall m: Model
$$

$$
\rightarrow
$$

$$
\exists e: Edge \bullet (e \in m)
$$
 (3.5)

Law 6. Any edge can be either an arc or a relationship.

$$
\forall a: Edge
$$

\n
$$
\longrightarrow
$$

\n
$$
Arc(a) \oplus Relationship(a)
$$
\n(3.6)

Law 7. Any arc can be either a flow or a transition or a stakeholder interaction (participation).

$$
\forall a: Arc
$$

\n
$$
\longrightarrow
$$

\n
$$
Flow(a) \oplus Transition(a) \oplus Partition(a)
$$
\n(3.7)

Law 8. A relationship can be a generalisation, composition, dependency or an association.

 $\forall r : Relationship$ \longrightarrow $Generalisation(r) \oplus Composition(r)$ \oplus Dependency(r) \oplus Association(r) (3.8)

Law 9. Each Node n has one unique identifier id .

$$
\forall n: Node
$$

$$
\rightarrow \exists! id: Identifier \bullet Id(n, id)
$$
 (3.9)

Law 10. In the model, each edge connects two nodes

$$
\forall m: Model, a: Edge, (n1, n2): Node
$$

• $(a, n1, n2 \in m), \text{Connects}(a, n1, n2)$ (3.10)

Law 11. In the Model, each Edge connects a Source Node to a Target Node

$$
\forall m : Model, a : Edge, (n1, n2 : Node)
$$

\n• $(a, n1, n2 \in m)$
\n $\land \text{Connects}(a, n1, n2)$
\n(3.11)

 $Source(a, n1) \wedge Target(a, n2)$

Law 12. In the Model, each Dependency connects a Source Node to a Target Node

$$
\forall m : Model, d : Dependercy, (n1, n2 : Node)
$$

\n• $(d, n1, n2 \in m)$
\n $\land Connects(d, n1, n2)$
\n \rightarrow
\n $Source(d, n1) \land Target(d, n2)$
\n(3.12)

Law 13. In the *Model*, each relationship links two nodes n_1 and n_2 , when it connects n_1 to $n2$ and connects $n2$ to $n1$.

$$
\forall m: Model, r: Relationship, (n1, n2: Node)
$$

\n
$$
\bullet ((r, n1, n2 \in m)
$$

\n
$$
\land Links(r, n1, n2))
$$

\n
$$
\rightarrow
$$

\n(3.13)

 $(Connects(r, n1, n2) \lor Connects(r, n2, n1))$

Law 14. In the *Model*, each association links two nodes.

$$
\forall m: Model, a: Association, (n1, n2: Node)
$$

\n•
$$
((a, n1, n2 \in m) \land Links(a, n1, n2))
$$
\n
$$
(3.14)
$$

According to the μ ML metamodel (Figure 3.2), the concept *Relationship* serves to declare directed and undirected relationships between two nodes or a node with itself. *End-roles* of this kind of Edge hold some additional features and accept further boolean predicates than those roles appear in Arc, such as multiplicities. Based on the type of relationship, the end-role might be, for instance, a whole/part, as it appears in Composition relationships, and a parent/child as it occurs in Generalisation relationships.

It worth saying that the predicate $Links(x, y, z)$ represents the undirected relationship (connection) between the node y and the node z, whereas, the $Connects(a, b, c)$ expresses the directed relationship between the node a and the node b , in which the relationship a connects the source b to the target c .

3.5 Outlook on the Chapter

The chapter introduced an overview of the proposed solution for the research problem, idenfitied earlier in Chapter 1. More specifically, the overall structure of the BUILD method is discussed, including the related system views, example of the graphical notation of models and critical transformation steps at each stage of developement. Development activities that are mastered by end-users or are led automatically by rules are also highlighted for each phase.

Furthermore, the overall structure and components of model transformation and code generation phase, within BUILD, are presented. This includes the internal architecture of transformation steps, the style of designing mapping rules *(hybrid)*, and programming language that is used to implement them (Java).

Moreover, the modelling language (μML) , which is used in BUILD, for expressing required system specifications is presented briefly in this chapter. The core elements of the language were defined using metamodelling strategy in order to express all critical concepts of the information system domain. Semantics of the metamodel elements are formalised using First-Order Predicate Logic (FOPL) with a number of selected extensions, namely, subtyping, equality, membership, exclusive disjunction and uniqueness quantification.

4 μ ML Concepts and Notations

In the Requirement Sketching Phase

"The most difficult part of requirements gathering is not the act of recording what the user wants, it is the exploratory development activity of helping users figure out what they want"

Steve McConnell

4.1 Context

This chapter describes, in detail, the Micro Modelling Language (μML) models appearing in the Requirement Sketching Phase, namely, Task, Impact and Information Model. For each model, the chapter presents how the model satisfies its purpose within the framework. Additionally it discusses each model concept and graphical notation for that concept, also providing a formally defined semantics using a set of FOPL policies.

4.2 Overview of the Task Model

A Task Model is the highest level, manually-constructed, diagrammatic representation of a business in terms of stakeholders, the tasks they perform and the goals to be achieved. It is a structural model that captures the wider context and the initial requirements of the business in a straightforward way by linking every task with a purpose or goal that is significant and worth recording in the business domain.

The participation of stakeholders, human or external system, in tasks is depicted in a clear way to be modelled naturally by the business end-user. Stakeholders might be specified / generalised to appear at different levels of abstraction. Precise types of participation and task occurrence, namely, multiplicity and optionality, are adopted to reflect the designer's natural intuitions about the ways in which they interact with the system and how the business is performed. These basic specifications are developed during the transformation stages into further design decisions used for implementation, such as iterations and conditional branches.

Apart from this, business tasks in the Task Model are fully compositional, and they may be aggregated or decomposed as desired to reveal their $part/whole$ structure. Additionally, they may be presented at different degrees of abstraction to reveal their *general/specialised* structure. There is no restriction on the ideal granularity of a task; a stakeholder might perceive tasks arbitrarily at different levels of detail.

The possible ways in which parts collaborate in order to perform the main task and meet the intended goal are considered from the end-users perspective. One part, for instance, might be executed independently to perform all aspects of the original task. Consequently, the parts are noted as options or choices. In some cases, internal flows between the parts might conceptually exist, each part performing an aspect of the main task. Both features can be captured by end-users using a simple notation.

Each group of business tasks that falls within the scope of a business practice might be modelled as a boundary that represents a separate part of the system. This feature can be adopted in order to document and organise the incremental delivery phases if planned. Technically, boundaries are used to describe the details of a composed task, including subtasks and their interactions with stakeholders.

The Task Model, in the proposed metamodelling hierarchy, is considered a kind of Model. Thus, concepts of the Task Model can be defined using FOPL and relate to the corresponding μ ML metamodel concepts.

Law 15. The TaskModel is a kind of the Model. Thus, every TaskModel is a Model.

$$
TaskModel \le: Model
$$
\n
$$
\forall m \cdot TaskModel(m)
$$
\n
$$
\longrightarrow
$$
\n
$$
Model(m)
$$

4.2.1 Notation and Semantics of the Task Model

In this section, each concept that appears in the Task Model is introduced with some details about its usage and how it is visualised graphically in the model. It is worth emphasising that a number of UML notations in the Use Case and Class Diagram are adopted and reused with some differences in their semantics. The differences allow us to apply the same notations in more modelling contexts, with a well-defined interpretation. This means our concepts might have a slightly different meaning from standard UML[70]. The following table (4.1) summaries all predicates used in describing μ ML Task Model.

In order to motivate the list of policies that specifies the Task Model, some definitions must be identified first. According to the μ ML metamodel hierarchy, the core elements of the Task Model are subtypes of either *Node* or *Arc*:

(Actor $\langle : \text{Node} \rangle \wedge (Task \langle : \text{Node} \rangle \wedge (Goal \langle : \text{Node} \rangle \wedge (Participation \langle : \text{Arc} \rangle$) \wedge (Input \langle : Participation) \wedge (Output \langle : Participation)

Predicate	Meaning	Syntactic Sugar
TaskModel(x)	x is a Task Model	x:TaskModel
Task(x)	x is a Business Task	x:Task
Goal(x)	x is a Business Goal	x:Goal
Actor(x)	x is an Actor	x: Actor
Human(x)	x is (Human) user	
ExSystem(x)	x is an external system	
Partition(x)	x is a Participation	x:Participation
Input(x)	x is an Input participation	x:Input
Output(x)	x is a Output participation	x:Output
Composition(x)	x is a Composition	x:Composition
Total(x)	x is a Total Composition	
Generalisation(x)	x is an Generalisation	x: Generalisation
Disjoint(x)	x is a Disjoint Generalisation	
Role(x)	x is an end-role	x:Role
General(x, y)	x is general in <i>Generalisation</i> y	
Specific(x, y)	x is specific in <i>Generalisation</i> y	
Whole(x, y)	x is whole in <i>Composition</i> y	
Part(x, y)	x is part in <i>Composition</i> y	

Table 4.1: Predicates for μ ML Task Model

From that, any Node in the model can be either a Task, Goal or an Actor. This can be expressed in logic as:

$$
\forall n: Node, \forall m: TaskModel
$$

\n• $(n \in_m m)$
\n \longrightarrow
\n $Action(n) \oplus Task(n) \oplus Goal(n)$

In the same context, any Arc appears in the Task Model is a Participation:

$$
\forall a: Arc, \forall m: TaskModel
$$

\n• $(a \in_m m)$
\n \longrightarrow
\n*Participation(a)*

4.2.1.1 Business Goal

The notion of a business goal is used to express some high-level aim, which is not directly measurable. It is an essential concept in the model from business management prespective. Some business tasks might be restructured during transformation steps. Relating some business tasks to a goal, at the early stage, enables the achievement of that goal to be measured, in the later design.

Because it is more abstract than a business task, we use a dashed outline ellipse to represents abstract aims in the model, as it illustrated in Figure 4.1. The semantics of Goals can be formalised using the following FOPL policies:

Figure 4.1: Task model. Goals

According to the μ ML metamodel, *Business Goal* is defined as a subtype of the *Node* element (Goal \leq : Node). Thus, it is identified using a unique identifier. This can be written formally in logic as:

$$
\forall m: TaskModel, \ \forall g:Goal \ \bullet \ (a \in_m m),
$$

$$
\exists! \ id: Identifier \ \bullet \ Id(g, m, id)
$$

4.2.1.2 Business Task

To depict tasks, we reuse the ellipse notation for a UML use case, but do not restrict this to mean a single, small-scale interaction, or use case. This allows for larger, aggregated tasks to be described, as well as single use cases. Tasks always achieve measurable objectives, within the business domain. Figure 4.2 illustrates the graphical notation of tasks in Task model.

Figure 4.2: Task model. Tasks

Similar to Business Goal, Business Task is defined as a subtype of the Node element (Task \leq : Node). Thus, it is identified using a unique identifier that formalised in FOPL as:

> $\forall m : TaskModel, \ \forall t : Task \bullet (t \in_m m),$ $\exists ! \ id : Identifier \bullet Id(t, m, id)$

4.2.1.3 Participant

The system stakeholders are represented as UML actors that appear in the model as Human or an external system participant $\forall a : Actor \longrightarrow Human(a) \oplus ExSystem(a)$. Human actors are depicted as stick figures, Figure 4.3 (left), whereas external systems are depicted as box nodes, reused from the UML deployment diagram, Figure 4.3 (right).

Figure 4.3: Task model. Actors

As it described in the μ ML metamodel, Actors in the Task Model are Nodes, (Actor \langle : Node), each participant (actor) has only one unique identifier:

$$
\forall m : TaskModel, \ \forall a : Actor \ \bullet \ (a \ \in_m m),
$$

$$
\exists! \ id : Identifier \ \bullet \ Id(a, m, id)
$$

4.2.1.4 Participation

Participation arrows capture system interactions with stakeholders, which can be represented as Input or Output participation. Both types are considered subtypes of Participation in the μ ML metamodel $(Input, Output) \leq: Participation$. Therefore, any *Participation* arrow appears in the Task Model either a Input or an Output arrow:

$$
\forall p : Participation, \forall m : TaskModel
$$

\n• $(p \in_m m)$
\n \longrightarrow
\n $Input(p) \oplus Output(p)$

The input arrow indicates that an actor participates in a task by supplying some external inputs (e.g information). This is determined by the direction of the arrow from the actor toward the task. Similarly, the output arrow indicates that an actor participates in a task by receiving some information from the system. Figure 4.4 illustrates the graphical notation of participation in the Task model.

Figure 4.4: Task model. Participation

The *Input* participation can be defined in the following theorem that follows law (axiom) 10, substituting variously Task or Actor for Node, and Input for Edge. This can be written as:

> $\forall p : Input, \forall t : Task, \forall a : Actor$ \bullet (Connects(p, a, t) $\land Source(p, a) \land Target(p, t))$

Following the similar way of defining *Input* participation, It can be said that any *Output* connects a task, as Source, to an actor, as Target. This can be written as:

> $\forall p : Output, \forall t : Task, \forall a : Actor$ \bullet (Connects(p, t, a) \land Source (p, t) \land Target (p, a))

4.2.1.5 Generalisation

Tasks may be described in general, or specific terms, in order to express the "is-a" relationship. In consequence, a task is regarded as an ancestor if it has a specialised task that inherits from it. In the same context, actors are treated likewise. We reuse the UML white triangle arrowhead notation for generalisation (inheritance), pointing to the general task (Figure 4.5) . The following figure illustrates that the Task A is a task generalisation of the Task B and C.

Figure 4.5: Task model. Generalisation

According to the μ ML metamodel, generalisation is considered a kind of *Relationship* (Generalisation \leq : Relationship) in which it connects the more specific task to the general one. The following logic expression represents the notion of Generalistion in the TaskModel:

$$
\forall t1, t2 : Task, \forall g : Generalisation, \forall m : TaskModel
$$

\n
$$
\bullet ((t1, t2, g \in_m m) \land Connects(g, t1, t2)) \longrightarrow
$$

\n
$$
General(s, t2) \land Special(s, t1)
$$

 $General(q, t2) \wedge Specific(q, t1)$

4.2.1.6 Composition

Aggregate tasks may be composed of other tasks to represent the "is part of" relationship. The notion of composition is a primary concept in compositional systems. We therefore prefer this term to the alternative term aggregation. We reuse the UML white diamond arrowhead and disjoint notations for whole/parts structures, pointing to the composed task (as a consequence, the include-dependency is redundant).

According to the μ ML metamodel, composition is considered a kind of Relationship (Composition \lt : Relationship) thatconnects the part tasks to the whole one. This relationship appears in the Task Model in two styles: total or partial composition $\forall c : Composition, m : TaskModel$. $((c \in_m m) \land (Total(c) \oplus \neg Total(c))$. The following logic expression represents the notion of Composition in the TaskModel:

$$
\forall t1, t2 : Task, \forall c : Composition, \forall m : TaskModel
$$

\n• ((t1, t2, c \in_m m) \land Connects(c, t1, t2))
\n
$$
\longrightarrow
$$

\n
$$
Whole(c, t2) \land Part(c, t1)
$$

From the business-user perspective, all "part" tasks involved in a white diamond (with *disjoint*) arrowhead composition represent options, or choices. In each execution, one part may perform independently the complete job of the ancestor task. This feature can be expressed using a unary predicate as: $c: Composition \wedge Disjoint(c)$. Figure 4.6 (a) below illustrates the visual representation of task optionality feature in the Task Model.

(a) Task model. White diamond (b) Task model. White diamond Composition (with disjoint) Composition

Figure 4.6: Task Model. Compositions

Besides, part tasks involved in a white diamond arrowhead composition indicate that there are such data flows and sequence interconnections between them, that they must be considered as all collaborating together to perform the complete job of their ancestor. Technically, there are internal variables that hold and pass data from a task to another to complete the ancestor mission. This feature can be expressed using the negation of *Disjoint* predicate as: c : Composition $\wedge \neg Disjoint(c)$. Figure 4.6 (b) demonstrates this.

As previously discussed, it is worth saying that the Task Model takes a different position from the use case model, where composition would create something larger than a use case. The current UML Use Case Diagram, does not formally compose use cases (the composition would no longer be a use case). The *Composition* relationship used here obviates the need for extra "include dependency" relationships. UML[70] uses composition to denote a specific flavour of aggregation, which unfairly limits the use of the term *composition*, which is used more generally in English to describe composed relationships.

4.2.1.7 Boundary

A rectangular system boundary may be drawn around groups of tasks, to indicate that these tasks fall within the scope of the system to be developed, while other tasks fall outside. Therefore, the Task Model might contain a number of logical Boundary nodes: $\forall m : TaskModel, \exists b :$ Boundary • $(b \in_m m)$.

4.2.2 The Significance of Task Model

The importance of Task Model emerges through its use of subtle concepts that allow the model transformation process to take different design decisions. These features are highlighted below and discussed later in detail in Chapter 4.

- Differentiating between the two types of stakeholders at early stage contributes to critical design decisions of user interfaces or system inter-operations, which will later be required.
- Differentiating between the two types of participations, $Input$ and $Output$, at early stage allows the automatic prediction of the Graphical User Interface (GUI) components and controls required in systems' screens.
- Differentiating between the two types of compositions at this stage automatically determines the control flow between the part tasks. Each subtask, produced at later stages, might be independently executed to perform the main task, or a particular flow might exist between them based on their atomic CRUD operations.

4.3 Overview of the Impact Model

The Impact Model is the highest level, **manually-constructed**, diagrammatic representation of a business task's interaction with the logical data within a system. It does not describe the time or order of processing but rather what information is produced and consumed by each task. It is assumed that each task has an effect on objects of reading, writing or updating (both reading and writing). These influences appear in the model via several kinds of arcs, each with a distinct meaning, that represent each *CRUD* (create, read, update and delete) operation.

The model might be compositional, expanding the level of detail at which a task is described. Like the Task Model, this feature is captured using system boundaries that contain the parts of the decomposed task with their interactions with the data. It is essential that the tasks in both the Task and Impact Model remain at the same level of granularity in order to generate a consistent layer in our layered model transformation approach.

The business data appears in the *Impact Model* as objects, corresponding to physical objects and documents, or logical data used by the business. It can be said that any object holds useful information about stakeholders or physical business items. Each object in the model has a lifespan that is determined by tracing the effects of CRUD on it.

The notions of object multiplicity and alternatives are also presented in the Impact Model to reflect end-user logical thoughts about how an object might participate in a task. This is determined from the perspective of the tasks using simple and easily-adopted notations. For each execution of the task, we identify how many other objects are affected. In addition, alternative kinds of data may be acceptable as inputs to or outputs from certain tasks. This is captured using a simple notation that does not complicate the model.

Law 16. The *ImpactModel* is a kind of *Model*.

$$
\begin{aligned} ImpactModel <: \ Model \\ &\forall \ m \ \bullet \ ImpactModel(m) \\ &\longrightarrow \\ Model(m) \end{aligned}
$$

4.3.1 Notation and Semantics of the Impact Model

In this section, each concept that occurs in the *Impact Model* is introduced with some details about its usage and how it is visualised graphically in the model. Similar to our strategy in adopting UML for the Task Model notation, a number of UML notations taken from the Class and Use Case Diagram are reused with slightly different meanings than the usual UML semantics. The semantics of the proposed concepts are expressed in the following sub-sections. The following table (4.2) summaries all predicates used in describing μ ML Impact Model.

Table 4.2: Predicates for μ ML Impact Model

In order to motivate the list of policies that specifies the *Impact Model*, two basic laws must be identified first as follows:

Law 17. Any node in the *ImpactModel* can be either a task or an object or a conjunction.

 $(ImpTask \leq: Node) \wedge (ImpObject \leq: Node)$ $\forall n : Node, \forall m : ImpactModel \bullet (n \in_m m)$ \longrightarrow

 $ImpTask(n) \oplus ImpObject(n) \oplus ImpConjunction(n)$

Law 18. Any flow in the *ImpactModel* can be either create, read, update, delete, write, input, or output flow.

 $(Flow \leq: Arc) \wedge (InputFlow \leq: Flow) \wedge$ $(OutputFlow \leq: Flow) \wedge (CreateFlow \leq: Flow) \wedge$ $WriteFlow \le: Flow) \wedge (ReadFlow \le: Flow) \wedge$ $(U\text{pdateFlow} <: Flow) \wedge (DeleteFlow <: Flow)$ $\forall f : Flow, \forall m : ImpactModel \bullet (f \in_m m)$ −→ $CreateFlow(f) \oplus ReadFlow(f) \oplus UpdateFlow(f)$ \oplus DeleteFlow(f) \oplus WriteFlow(f) \oplus InputFlow(f) \oplus OutputFlow(f)

Law 19. Any node in the *ImpactModel* is involved in at least one flow as a source or a target.

$$
\forall n: Node, \forall m: ImpModel
$$
\n• $(n \in_m m)$ \n
$$
\Rightarrow
$$
\n
$$
\exists f \bullet Flow(f) \bullet (f \in_m m) \land (Source(f, n) \oplus Target(f, n))
$$

4.3.1.1 Impact Task

Tasks are depicted as in the task model, using the ellipse node. In this model, the emphasis is on the impact that tasks have upon objects at a given level of granularity, so no task structure is shown. Likewise, no actors or external systems are depicted. Figure 4.7 below demonstrates the A and B are two Impact tasks.

Figure 4.7: Impact model. Tasks

In regard to the μ ML metamodel, Impact Task is considered a subtype of the metaclass Node element ($ImpTask \leq: Node$). Therefore, it is identified using a unique identifier. This can be written formally in logic as:

> $\forall m: ImpactModel, \forall t: ImpTask \bullet (t \in_m m),$ $\exists !\ id : Identifier \bullet Id(t, m, id)$

Given (ImpTask \langle : Node) from the defined metamodel, it can be deduced that any ImpTask in the ImpactModel can be either a source or a target of such a flow. This is a theorem that follows from the earlier law 19, given previously in Chapter 3, by substituting of Node.

−→

 $\forall t : ImpTask, \forall m : ImpModel \bullet (t \in_m m)$

 $\exists f \bullet Flow(f) \bullet (f \in_m m) \land (Source(f,t) \oplus Target(f,t))$

4.3.1.2 Impact Object

The UML rectangular nodes, taken from the UML class diagram, are used to represent information objects, corresponding to physical business objects (entities) and documents, or logical data used by the business in the real-world (Figure 4.8). Any object that bears useful information may be modelled. This may include objects storing logical information about human actors, or other physical things.

Figure 4.8: Impact model. Objects

Similar to the way of declaring *Impact Task* according to the μ ML metamodel, Impact Object is also considered a subtype of the *Node* element $(ImpObject \leq: Node)$. Therefore, it is identified using a unique identifier. This can be written formally in logic as:

> $\forall m : ImpactModel, \forall obj : ImpObject \bullet (obj \in_m m),$ $\exists ! \ id : Identifier \bullet Id(obj, m, id)$

Similar to the ImpTask, given $(ImpObject \leq; Node)$ from the defined metamodel, it can be deduced that any ImpObject in the ImpactModel can be either a source or a target of such a flow. This is a theorem that follows from the earlier law 19, by substitution of *Node*.

$$
\forall \text{ obj}: ImpObject, \forall m: ImpModel \bullet (\text{obj } \in_m m)
$$
\n
$$
\exists f \bullet Flow(f) \bullet (f \in_m m)
$$
\n
$$
\land (Source(f, obj) \oplus Target(f, obj))
$$

4.3.1.3 Impacts (flows)

The impacts that occur on objects are shown visually by diverse-headed arrows. The style of arrowhead indicates the direction of information flow, or the impact, between a task and an object. Figure $4.9¹$ below illustrates the various types of impacts in the *Impact Model*.

Figure 4.9: Impact model. Impacts

Given (Flow \langle : Arc \langle : Edge), (ImpModel \langle : Model) from the constructed metamodel, it can be deduced that each Flow connects two distinct nodes. This is a theorem that follows from Law 10, by substituting for Model and Edge.

> $\forall (n1, n2): Node, \forall m: ImpactModel, \forall f: Flow$ • $((f \in_m m) \land Connects(f, n1, n2))$

In order to define formally the various types of impacts in the model, a basic law is presented:

Law 20. Any flow, in the *Impact Model*, that connects an object to a task is a read flow, whereas any flow connects a task to an object can be any flow out of a *create*, delete or *write* flow. This can be expressed in FOPL as:

$$
\forall m: ImpactModel, \forall f: Flow, \forall e: ImpObject,\n\forall t: ImpTask \bullet (e, t, f \in_m m) \land Connects(f, e, t)\n\rightarrow\n\qquad\n\begin{array}{r}\n\text{ReadFlow}(f) \\
\land Connects(f, t, e) \\
\end{array}
$$
\n
$$
\rightarrow\n\begin{array}{r}\n\text{ReadFlow}(f) \\
\end{array}\n\Rightarrow
$$
\n
$$
\text{CreateFlow}(f) \oplus WriteFlow(f) \oplus DeleteFlow(f)
$$

Visually, the open-headed arrows are utilised to represent the impact on objects, which may be read, or written, Figure 4.9(B). Figure 4.10a below demonstrates that the *Object Obj1* is read by the the Task A, whereas, Figure 4.10b shows that the $Obj1$ is written by a Task A.

Figure 4.10: Impact Model. Read/Write Object

¹The interpretation of Read/Write flows depends on whether the source or target is an object.

In addition to this, a double-ended arrow indicates both reading and writing from the same object, which presents an updating effect, Figure $4.9(C)$. The following law expresses the decomposition of an Update into separate Read and Write flows.

Law 21. Any flow, in the Impact Model, that connects bi-directionally an object to a task is an update flow. This can be expressed in FOPL as:

$$
\forall m: ImpactModel, \forall f: Flow, \forall e: ImpObject, \forall t: ImpTask
$$

\n•
$$
e, t, f \in_m m) \land Connects(f, e, t) \land Connects(f, t, e)
$$

\n
$$
\longrightarrow
$$

\n
$$
UpdateFlow(f)
$$

The basic interpretation of this decomposition step means that each object, involved in this type of impact, must be read first, before it can be updated or written. This interpretation comes into play later on, when timing becomes significant, since the current model is simply concerned with flow structure. The following figure (4.11) shows that the *Obj1* is updated by a Task B.

Figure 4.11: Impact model. Update

Furthermore, two additional types of arrows are used to express the creation and deletion of objects, as illustrated in Figure $4.9(a)$ and $4.9(d)$ respectively. These Impacts are presented via the following figures. Figure $4.12(a)$ demonstrates that the object Obj1 is created by the the Task C, whereas, Figure 4.12(b) illustrates that the *Obj1* is destroyed (deleted) by a Task B. It is significant to note that objects must already exist (created) first to participate in any other type of dataflow.

Figure 4.12: Impact Model. Create/Delete Object

4.3.1.4 Multiplicity and Optionality of Impacts

The proposed μ ML language adopts the UML style for expressing cardinality to represent the number of elements involved in such a relationship or flow. It generally expresses the statement: at least m but no more than n objects, in the UML multiplicity adornment as $(m..n)$. The notion of upper bound and lower bound are also exists to specify the range of elements. Each bound might be an exact positive integer number or unlimited number of elements, denoted by asterisk symbol $(*)$. The following table (4.3) exemplifies multiplicity:

As a default interpretation of the model, any Impact represents a one-to-one correspondence. This means that one instance of each datum, modelled as an Impact Object, read or written for each execution of the task. For simplicity, there is no need to attach $(1..1)$ multiplicity for each flow unless the impact correspondence is not one-to-one.

4.3.1.5 Disjoint Impact Combinator (Alternative)

The disjoint flow combinator joins two or more optional flows, where these alternate exclusively. This exclusive alternation concept is shown in the *Impact Model* by placing a clear circle symbol between a number of flows. Figure 4.13 below illustrates this.

Figure 4.13: Impact model. Disjoint Impact Combinator

In the foundation of the *Impact Model*, a unary predicate $(ImpConjunction)$ is defined to determine Nodes that are disjoint flow combinators in the model. Following the similar way of declaring Impact Task and Object according to the μ ML metamodel, Disjoint flow combinator is also considered a subtype of the *Node* element $(ImpConjunction <:Node)$. Therefore, it is identified using a unique identifier. This is a theorem that follows from the earlier law 9, by substituting for Node.

> $\forall m : ImpactModel, \forall c : ImpConjunction \bullet (c \in_m m),$ $\exists !\ id : Identifier \bullet Id(c, m, id)$

Because disjoint flow combinators link some flows in ImpactModel, it can be either a source or a target similar to the $ImpTask$, any $ImpObject$. This is a theorem that follows from the earlier law 19, by substituting for Node.

 $\forall c : ImpConjunction, \forall m : ImpModel \bullet (c \in_m m)$ \longrightarrow $\exists f \bullet Flow(f) \bullet (f \in_m m)$ \wedge (Source(f, c) ⊕ Target(f, c))

4.3.1.6 Impact Boundary

Similar to the previously presented Task Model, a rectangular system boundary may be drawn around groups of tasks, to indicate that these tasks fall within the scope of the system to be developed, while other tasks fall outside. From that, the *Impact Model* might contain a number of logical ImpBoundary nodes: $\forall m : ImpactModel, \exists b : ImpBoundary \bullet (b \in_m m)$.

4.3.2 The Significance of Impact Model

The Impact Model aims to generate a partial ordering of tasks by realising data dependency captured via data participation in tasks. This contributes to predicting and constraining the possible orders of task execution in later design stages. The execution of a task may depend on the prior existence of certain data, and it may in turn enable other tasks to be performed due to the data it produces.

Therefore, the order of execution can in part be determined by tracing the CRUD effects of tasks on a single object or multiple objects that fall in the same business scope (system boundary). By considering each task in Figure 3.4, shown previously in Chapter 3, the creation of the Loan object, for instance, depends on the completion of reading the Copy and Borrower objects. Likewise, the deletion of Loan requires reading Copy first.

Furthermore, partial object life histories can also be determined by tracing the CRUD effects of tasks on single objects. Each creation and deletion event occurs only once, but other events may occur multiple times in any unspecified order. As a consequence, data dependency, the fundamental property of the *Data Model*, may be determined directly via analysis of the object lifetimes in the impact model. When more than one obejct is involved in the same task, with different types of *impacts*, this means that there exist dependency relationships between these objects, in which shorter-lived objects always depend on longer-lived ones.

Tracing the CRUD effects of a single task execution on individual objects (viz. a set of individuals) also provides useful information for the Data Model since knowing the relative extent (life expectancy) of each object also informs the notion of data dependency. From this, we can assume that it is possible to predict target concepts to generate a partial Data Model to express the interdependency of the business entities.

4.4 Overview of the Information Model

The Information Model is the highest level, **manually-constructed**, diagrammatic representation of the structure of business objects. The main purpose of the Information Model is to capture sufficient information on business data to generate a complete Data Model. We are following the standard approach taken in software engineering in which the model is presented in terms of conceptual entities (objects), attributes and inter-relationships. As a basic inspiration, we view conceptual modelling as a separate activity from database design, performed by business analysts or users in a very abstract way.

Entities in the model represent physical objects and documents or logical data used by the business that may, or may not, be in normal form. Any object that bears useful information may be modelled, including those storing logical information about stakeholders or other physical items. It is common for these kinds of objects to require further logical decomposition, especially if they contain repeating groups of data.

An entity consists of properties, modelled as attributes, which describe that entity (object). It is important for a business-user who is a domain expert to test the properties of each object in order to ensure that they are atomic and that their values depend fully on the object instance owning them. For each object instance, the attribute may take on a distinct value.

Apart from this, in the real-world, business objects relate and interact to present the domain of the business. The Information Model represents this via undirected conceptual associations, which link objects together, indicating that the related objects are acquainted in some fashion. Additionally, relationships that require grouping sets of objects that are treated together or share common properties, such as composition and generalization, respectively, can be expressed as other types of association to imply further semantics.

The various relationships provide clean and precise semantics about the notation of the model elements for the business-user. They show details, such as multiplicity constraints, which hide technically significant information to determine data dependency, using our promoting and qualifying policies until relationships are all of the many-to-one kind, with the many always depending on the one.

The Information Model, in the proposed metamodelling hierarchy, is considered a kind of Model. As a consequence, concepts of the Information Model can be defined using FOPL and relate to the corresponding μ ML metamodel concepts.

Law 22. The *InformationModel* is a kind of *Model*.

(InformationModel <: Model) $\forall m \bullet InformationModel(m)$ \longrightarrow $Model(m)$

4.4.1 Notation and Semantics of the Information Model

In this section, each concept that appears in the *Information Model* is introduced with some details about its usage and how it is visualised graphically in the model. The following table (Table 4) summaries all predicates used in describing μ ML Information Model.

In order to motivate the list of policies that specifies the Information Model, two basic laws must be identified first as follows:

Law 23. Any node in the *InformationModel* is an object.

$$
(InfEntity <: Node)
$$
\n
$$
\forall n: Node, \forall m: InformationModel
$$
\n
$$
(n \in_m m)
$$
\n
$$
\longrightarrow
$$
\n
$$
InfEntity(n)
$$

Table 4.4: Predicates for μ ML Information Model

Law 24. Any edge in the *Information Model* is a relationship.

$$
(Relationship <: Edge)
$$
\n
$$
∀ a : Edge, ∀ m : InformationModel
$$
\n
$$
• (a ∈ m m)
$$
\n
$$
Relationship(a)
$$
\n
$$
Relationship(a)
$$

Law 25. Any relationship in the *Information Model* can only be one out of: a generalisation or composition or association.

(Generalisation <: Relationship ∧ Composition <: Relationship ∧ Association <: Relationship) $\forall r : \mathit{Relationship}, \forall m : \mathit{InformationModel}$ \bullet $(r \in_m m)$ \longrightarrow $Generalisation(r) \oplus Association(r) \oplus Composition(r)$

4.4.1.1 Object

Similar to the object's notation used in the *Impact model*, the UML rectangular nodes, taken from the UML class diagram are adopted to present information objects in the model (Figure 4.14). The information objects correspond to physical business entites and documents, or logical data used by the business. Any object that bears useful information may be modelled. This may include objects storing logical details about human actors, or other physical things.

Figure 4.14: Information model. Objects (Entities)

It is worth saying, physical objects are often more complex than their logical counterparts, because they may contain repeating groups of data that should eventually be split. According to the μ ML metamodel, *Information Object* is also considered a subtype of the *Node* element $(InfEntity <:Node)$. Therefore, it is identified using a unique identifier. This can be defined in FOPL as:

> $\forall m : InfModel, \forall obj : InfEntity \bullet (obj \in_m m),$ $\exists ! \text{ id}: \text{Identifier} \bullet \text{ Id}(\text{obj}, m, \text{ id})$

Each object in the model might involved in one or more than one relationship with other objects. This can be expressed formally as:

 $\forall m : InformationModel, \ \forall r : Relationship,$ \exists obj1, obj2 : InfEntity \bullet ((obj1, obj2, $r \in_m m$) \longrightarrow $(Connects(r, obj1, obj2) \lor Connects(r, obj2, obj1)))$

4.4.1.2 Association

Associations express conceptual "has a" relationships between objects, between objects of any class, expressed formally as Association <: Relationship. The UML straight lines are used to represent the links, that are usually bounded by events that involve the objects concerned. These relationships are undirected, meaning that the dependency of one object upon another is not yet known or understood. Figure 4.15 below, A is an object associated with another object B. Associations may optionally be named.

Figure 4.15: Impact model. Associations

The object might involved in one or more than one association with other objects, or it might refer to itself (self-relationship). This theorem follows from Law 13 by substituting for *Model*, Node and Relationship:

 $\forall m : InformationModel, \forall r : Association,$ \exists obj1, obj2 : InfEntity • $((obj1, obj2, r \in_m m))$

$$
\over\longrightarrow
$$

 $(Connects(r, obj1, obj2) \lor Connects(r, obj2, obj1)))$

4.4.1.3 Composition

Composition is a conceptual whole/parts relationship, or "is part of", in the domain, expressed formally as: Composition

 \leq : Relationship. It groups sets of objects that are treated together, indicating data dependency of the whole on the parts, which is assumed by default. The parts may exist independently of the whole. Composition is shown using the white UML diamond arrowhead notation, with the arrow pointing to the whole concept (Figure 4.16). The notion of composition is a primary concept in compositional systems. We therefore prefer this term to the alternative term aggregation. UML uses composition to denote a specific flavour that we call total composition (section 4.4.1.4 below).

Figure 4.16: Information model. Composition

The following policy defines the concepts of Composition in the Information Model and the notion of its Parts and Whole.

Law 26. Any compostion, in the Information Model connects Part objects to the Whole one. This can be expressed in FOPL as:

$$
\forall m: InformationModel, \forall c: Composition, \forall obj1, obj2: InfEntity
$$

\n• ((c, obj1, obj2 $\in_m m)$ \land $Connects(c, obj2, obj1))$
\n
$$
\longrightarrow
$$

\n
$$
Whole(obj1) \land Part(obj2)
$$

4.4.1.4 Total Composition

A total composition is one in which the parts belong entirely to the whole, in the sense that they cannot exist independently without the whole. Total compositions are significant, because the direction of data dependency is reversed: the parts depend on the whole. For instance, deletion of the whole must result in a cascading deletion of the parts. This is shown by placing a filled circle1 over the composition, to indicate total ownership of the parts. The following figure (4.17) indicates that A is composed of B and C , which cannot exist without the A . It can be declared using Law 28 above, as Total is irrelevant to deducing the whole and part.

Figure 4.17: Information model. Total Composition

4.4.1.5 Generalisation

Generalisation is a conceptual general/specific, or "is a", relationship in the domain. It groups sets of objects that share some common properties, indicating data dependency of the specific on the general. Generalisation is expressed using the UML triangle arrowhead notation, with the arrow pointing towards the more general concept. The following figure (4.18) illustrates that B and C are specialised cases of A.

Figure 4.18: Information model. Generalisation

Law 27. In the *InformationModel*, each *Generalisation* Connects *Child* objects to a *Parent* one. This can be formally expressed as:

$$
\forall m: InformationModel, \ \forall g: Generalisation, \ \forall \ obj1, obj2: InfEntity
$$

\n• ((f, obj1, obj2 \in_m m) \land Connects(g, obj2, obj1))
\n
$$
\longrightarrow
$$

 $Parent(obj1) \wedge Child(obj2))$

4.4.1.6 Disjoint Generalisation

A *disjoint* generalisation is one in which the specific concepts are mutually exclusive, in the sense that no object from the domain could be an instance of both concepts simultaneously. This is shown by placing a clear circle over the generalisation, to indicate that the specialisations are disjoint. Figure (4.19) indicates that A may be either B, or C, but not both in the given domain.

Using the same law for defining the (overlapping) generalisation (law 29), the disjoint generalisation also connects Child objects to a Parent one.

```
\forall m : InformationModel, \forall g : Generalisation, \forall obj1, obj2 : InfEntity• ((f, obj1, obj2 \in_m m) \land Disjoint(g) \land Connects(g, obj2, obj1))
```
 $Parent(obj1) \wedge Child(obj2)$

 \longrightarrow

Figure 4.19: Information model. Generalisation

4.4.1.7 Attribute

Attributes represents the atomic properties of entities (objects), in which each entity has one or more attribute (\forall obj : InfEntity, $\exists a : \text{Attribute} \bullet (a \in_b \text{obj})$). Similar to the UML class diagram, they are visually located in the usual drop-down partition of the Information Model obejcts the following Figure (4.20) illustrates the representation of properties, including the name and the data type, for two information entities: A and B.

Figure 4.20: Information model. Attributes

4.4.1.8 Multiplicity

In the *Information Model*, objects are related to each other with a given multiplicity. For each instance of object at one end of the association, there exist zero, one or more instances of the second object at the opposite end. The usual UML multiplicity adornments, examplified previously in table 4.3, are adopted to mark various kinds of multiplicity, namely, one-to-many, zero-to-many, many-to-many, and optional.

Figure 4.21: Information model. Multiplicity

Figure 4.21 (A) indicates that there are *zero-to-many* $A(s)$ for each entity B. However, binary associations in Figure 4.21 (B) represents the decomposition of a ternary or higher-arity association between three entities: C, D, and E. Each N-arity association must be broken down into a set of binary associations in order to be analysed in a later stage.

4.4.2 The Significance of The Information Model

- Associations are analysed later to reveal data dependencies, using rules based on multiplicity. This analysis is later used to normalise tables ready for database implementation.
- Properties (Attributes) are captured for each entity. The attributes later become the columns of database tables. Some attributes may be marked as uniquely identifying the owning entity.
- In circumstances where it is impossible to find a unique identifier for an entity, this may be manufactured automatically.
- The Information Model may still capture atomic data at a high level of abstraction, for example, using types such as Currency, Text, Time, which may need translation in later stages.

4.5 Outlook on the Chapter

This chapter discussed, in-depth, the concepts, adopted graphical notations, and formal semantics for each Micro Modelling Language (μML) model appearing in the *Requirement Sketching* phase. It covered all critical elements, their usage, and their all possible interpretations within the various contexts of the μ ML Task, Impact and Information Model.

Additionally, the chapter discussed how the selected system views contribute to the main goals of the thesis. More specifically, it showed how the proposed user-friendly modelling language, which has simpler and cleaner semantics than current UML-based approaches, captured the most critical aspects of the system using business user knowledge. Moreover, the chapter discussed how the core elements of this μ ML notation succeed in raising the level of abstraction, compared to UML, at which business-users may express their model specifications.

The semantics and notation for each model and element are compared to similar or closely related UML concepts and notation, showing the simplicity and clarity of the proposed language. A number of unary and n-ary predicates were introduced for defining concepts and their interrelationships within each model in respect of the previously defined metamodel.

5 μ ML Concepts and Notations

In the Analysis Phase

"The Analyst wants to get at the simplest form of the system which has the features they are interested in" Analysis Wisdom - ultradark

5.1 Context

This chapter describes in detail the Micro Modelling Language models used in the Analysis Phase, namely, Data (Dependency), DataFlow (initial and detailed) and (Screen) State Model. For each model, the chapter presents how the model satisfies its purpose within the framework. The chapter also discusses each model concept and its corresponding graphical notation, defining the concept formally using FOPL.

5.2 Overview of the Data Dependency Model

The Data Dependency Model is an **automatically-generated** intermediate level diagrammatic representation that is intended to describe the logical data of the system and support the development to a point where a logical database schema may be generated. The model consists of a number of logical objects linked by a number of dependency relationships, representing the direction of data dependency. This also determines which additional attributes will be used for references between objects and handled in the next stage of the layered transformation to generate the Database Schema Model. In other words, this determines where the forthcoming foreign keys will be located. The following figure (Figure 5.1) demonstrates the visual representation of the Data Model elements.

The entities represent logical, rather than physical, data used by the business in the realworld from which all repeating groups have been isolated. Any entity (object) bearing useful logical information, including stakeholders or other business items, is modelled. It contains

Figure 5.1: Data model. (A) Entity, (B) Dependency

attributes corresponding to the named columns of tables in a relational database that will be generated in the next stage of transformation.

Attributes are the atomic properties of objects. They are wholly functionally dependent on their owning object for their value; for each object instance, the attribute may possibly take on a distinct value. Only simple properties that are not further decomposable are modelled as attributes. Decomposable properties must be modelled as distinct objects instead, unless the target database can treat them in the same way as a basic type. At a higher level, the business-user specifies one or more attributes that eventually identify the object uniquely. This is taken for granted in the generated $Data Model$. Otherwise, if the identifier cannot be found. an artificial identifier is created.

In the current version of BUILD, a complete *Data Dependency Model* is constructed automatically from a multiplicity analysis of conceptual associations in the pre-designed Information Model. All associations are either promoted or translated directly into directed dependencies in which the many-sided object depends on the one-sided one. In worth mentioning that the Data Dependency Model may be, alternatively, constructed partially from an impact analysis of events in the *Impact Model*, an object being created from a pre-existing one. Either or both of these prior business user predefined models may be used as a source for the data model and may be crosschecked for consistency.

The Data Dependency Model, in the proposed metamodelling hierarchy, is considered a kind of Model. Thus, conepts of the Data Dependency Model can be defined using FOPL and relate to the corresponding μ ML metamodel concepts.

Law 28. The DataModel is a kind of Model.

$$
DataModel \le: Model
$$

\n
$$
\forall m \bullet DataModel(m)
$$

\n
$$
\longrightarrow
$$

\n
$$
Model(m)
$$

5.2.1 Notation and Semantics of the Data Model

In this section, each concept that appears in the Data Model is discussed in detail. This includes its usage and how it is visualised graphically in the model. It is worth emphasising that a number of UML notations in the Class Diagram are adopted and reused with slight different meaning from its original semantics. For instance, underlining attributes in the UML Class Diagram is used to express static attributes. The meaning of underlining attributes is changed here to express primary keys The following table (5.1) summaries all predicates used in describing μ ML Data Dependency Model.

Predicate	Meaning	Syntactic Sugar
DataModel(x)	x is a Data Model	x:DataModel
DEntity(x)	x is a Data Model Entity	x:DEntity
DAttribute(x)	x is a Data Model Attribute	x:DAttribute
Dependency(x)	x is a Dependency	x:Dependency
DependsOn(x, y)	x depends on y	
DRole(x)	x is a Data Model Role	x:DBole

Table 5.1: Predicates for μ ML Data Dependency Model

In order to motivate some policies that specifies *Data (Dependency) Model*, basic declarations of core elements must be identified first. According to the μ ML metamodel hierarchy where Data Model entity is a subtype of Node and Dependency is a subtype of Arc, expressed as: (DEntity \langle : Node \land Dependency \langle : Arc), then any Node and Arc in the Data Model represents an entity and a dependency respectively. This can be expressed formally as:

$$
\forall n: Node, \forall a: Arc, m: DataModel
$$

\n• $(n, a \in_m m)$
\n \longrightarrow
\n $DEntity(n) \land Dependercy(a)$

5.2.1.1 Entity and Attributes

Similar to Information Model objects and their properties, entities in the Data Model are drawn as rectangular nodes and their attributes are treated likewise in the previously described model, in which each entity has some attributes $\forall e : DEntity, \exists a : DAttribute \bullet (a \in_n e)$. The representation of Entity and Attribute are adopted from the well-known UML class diagram. This can be shown in Figure 5.1 (A) above, Figures 5.3 and 5.4 below.

5.2.1.2 Dependency

As illustrated in Figure 5.1 (B) above, dependencies in the Data Model are drawn as a directed edge, with an open arrowhead pointing toward the target from the dependent source on which it depends. Figure (5.2) demonstrates that the object A is dependent upon the object B, which can be formalised in FOPL as:

$$
(DEntity \le: Node \land Dependency \le: Arc)
$$

\n
$$
\forall m : DataModel, d : Dependency, (e1, e2 : DEntity)
$$

\n•
$$
((d, e1, e2 \in_m m) \land Connects(d, e1, e2))
$$

 $DependsOn(e1, e2)$

Figure 5.2: Generated Data model. Dependencies

5.2.1.3 Identifier

In the Data Model, each object (entity) must have at least a single attribute that identifies it uniquely, or several attributes, whose values taken together uniquely identify the object.

```
(Identifier <: DAttribute)\forall m : DataModel, \ \forall e : DEntity,∃! id : Identifier • ((Id(obj, m, id)) ∧ id ∈<sub>p</sub> e)
```
Under all circumstances, if no identifier can be found in such an entity, an artificial identifier is created. In the model, identifiers are shown by underlining the attributes concerned.

Figure 5.3: Generated Data model. Primary Key

This contrasts with the usual meaning of underlining in UML notation, which denotes a static (shared or class) attribute. From a data modelling perspective, a static attribute is not wholly functionally dependent on its owning entity, so logically should belong to a separate single entity, representing the constants of the class, on which the current entity depends by association.

5.2.1.4 References

For each dependency between one object and another, certain additional attributes must be chosen for the source object to uniquely refer to the target object. The references of an object correspond exactly to the identifiers of the objects to which they refer by association. The references are simple, if the identifier of the target is simple, and compound if the identifier is compound. From that, it can be said that the structure of each reference is identical to the structure of the *identity* of the obejct it refers to.

Figure 5.4: Generated Data model. Foreign Key

References are implicit in the dependency structure, but may optionally be visually shown in a second drop-down box (Figure 5.4). This replaces the use of the second drop-down box in the UML class diagram, which represents the methods of a class.

5.3 Overview of the DataFlow Models

The DataFlow Model is an **automatically-generated** intermediate level diagrammatic representation that that shows how data flows between the internal tasks and data stores, and external entities at a particular level of abstraction. The model consists of a number of entities and participants (agents) that are connected to some tasks via various flows. It provides a detailed description of the business tasks and aggregated subtasks that s perform operations on data in terms of data flow.

Unlike the *Impact Model*, which demonstrates how a single object is impacted by tasks, the DataFlow Entity represents a collection of objects that are directly matched to a formalised entity appearing in the *Data Model*. This allows the representation of crosschecked instances and attributes that are carried on the flows linked to an entity.

In addition, data on flows is expressed using formal expressions in terms of attributes, instances of entities (objects), local variables and values. This textual notation also represents filtering, projecting and assigning values when read from or written to an entity. It is worth saying that the design of the textual statements considers the business-users capability to express the data flows and related operations without any difficulty. As a result, the produced DataFlow Model holds readable data with more accurate interpretations of the flows.

Much like the previously introduced *Impact Model*, it does not describe the time-order of processing but rather what data passes between tasks, users and entities. A partial graph of flows between the system components of each part of the system is automatically derived by merging concepts from the matched boundaries that appear in the pre-constructed Task and Impact models. This produces a number of artifacts that are divided into boundaries equivalent to those existing in the requirement models, forming an initial DataFlow Model. It is worth mentioning that flows between tasks are not allowed, at this level of abstraction, in which there is insufficient information about task ordering that prevents predicting these additional flows.

The initial *DataFlow Model* is considered to be a top level partial model with no precise information about the data on flows. Business users, who are aware of the kinds of data transferred between the system components, are responsible for specifying the data and expressing basic operations meant by flows, using restricted statements recognized by the model. Using this external information enables the complete evolution of the detailed *DataFlow Model*. Figure 3.7, Chapter 3, demonstrates the contents of DataFlow Model, including, a task, flow, entity and boundary concept.

The detailed DataFlow Model is considered to be a bottom level comprehensive model that consists of a number of boundaries of decomposed tasks (atomic tasks) based upon their atomic actions. Each subtask, which is also a task, has only one associated flow and one atomic action. This new collection of boundaries forms the final DataFlow Model that contains comprehensive information regarding data flows and related operations. Figure 9.21, Chapter 9, illustrates a detailed DFD model including its concepts similar to the initial DFD model.

The DataFlow Model, in the proposed metamodelling hierarchy, is considered a kind of Model. Thus, concepts of the *DataFlow Model* can be defined using FOPL and relate to the corresponding μ ML metamodel concepts.

Law 29. The DataFlowModel is a kind of Model.

 $DataFlowModel$ <: Model $\forall m \bullet DataFlowModel(m)$ −→ $Model(m)$

5.3.1 Notation and Semantics of the DataFlow Model

In this section, each concept that appears in the DataFlow Model is introduced with some details about its usage and how it is visualised graphically in the model. It is worth mentioning that the notation introduced in the previously defined Task, Impact and Information Model are resused in the DataFlow to indicate slightly different meanings. The following table (5.2) summaries all predicates used in describing μ ML DataFlow Model.

Table 5.2: Predicates for μ ML DataFlow Model

In order to motivate the list of policies that describes DataFlow Model, two basic laws must be declared first as follows:

Law 30. Using the notion of subtyping in our FOPL approach, any Node in the DataFlowModel can be either a DfTask , a DfEntity or a DfActor.

$$
(DfActor \le: Node) \land (DfTask \le: Node) \land (DfEntity \le: Node)
$$

$$
\forall n: Node, m: DataFlowModel \bullet (n \in_m m)
$$

$$
\longrightarrow \qquad \qquad \longrightarrow \qquad \qquad \qquad \longrightarrow \qquad \qquad \qquad \qquad \qquad \longrightarrow
$$

$$
DfActor(n) \oplus DfTask(n) \oplus
$$

$$
DfEntity(n)
$$

Law 31. Using the notion of subtyping in our FOPL approach, any Flow in the DataFlowModel can be either DfCreateFlow, DfReadFlow, DfUpdateFlow, DfDeleteFlow, DfWriteFlow, DfInput-Flow, or DfOutputFlow. In the following sections, each concept that appears in the DataFlow Model is introduced with some details about its usage and how it is visualised graphically in the model.

 $(Flow \leq: Arc) \wedge (DfInputFlow \leq: Flow) \wedge$ $(DfOutputFlow \leq: Flow) \wedge (DfCreateFlow \leq: Flow) \wedge$ $(DfWriteFlow \le: Flow) \wedge (DfReadFlow \le: Flow) \wedge$ $(DfUpdateFlow \leq: Flow) \wedge (DfDeleteFlow \leq: Flow)$ $\forall f : Flow, m : DataFlowModel \bullet (f \in_m m)$

−→

 $DfCreateFlow(f) \oplus DfReadFlow(f) \oplus$ $DfUpdateFlow(f) \oplus DfDeleteFlow(f) \oplus$ $DfWriteFlow(f) \oplus DfInputFlow(f) \oplus$ $DfOutputFlow(f)$

5.3.1.1 Task

Similar to the notation adopted for tasks in the Task and Impact model, DataFlow Tasks are depicted using the ellipse node. In this model, the emphasis is on the actual data produced/consumed by business tasks from/to system entities at a given level of granularity. The external data usage and production with actors or external systems are also depicted. It is worth mentioning

that during the development process, the decomposed DataFlow Tasks are also presented in the detailed DFD using the same notation (ellipse node). Figure 5.5 below depicts the two Data Flow tasks, A and B.

Figure 5.5: Generated DataFlow model. Tasks

According to the previously defined μ ML metamodel hierarchy $(DfTask \leq: Node)$, each dataflow task has only one unique identifier. This can be expressed using FOPL predicate as: $\forall m : DataFlow, \forall t : DfTask \bullet (t \in_m m), \exists ! i : Identifier \bullet Id(t, m, i)$

5.3.1.2 Entity

Similar to the notation for objects in the *Impact*, *Information* and *Data model*, the UML rectangular nodes ar used to depict $DataFlow Entity$. The Entity is totally derived from the Impact Model and might be cross-checked with the Data/Information Models. In this model, it refers to a collection of objects, or a type, which is a distinct meaning compared to its meaning in other models. Figure 5.6 below depicts two Data Flow entities A and B.

Figure 5.6: Generated DataFlow model. Entites

A DataFlow Entity is considered subtype of Node in the μ ML metamodel hierarchy (DfEntity \lt : *Node*), therefore, it has only one unique identifier $\forall m : DataFlow, \forall e : DfEntity \bullet (e \in_m$ m), $\exists!$ i: Identifier • Id (e, m, i) . They replace the notion of data sources in the traditional DFD, in which each entity consists of a number of instances $\forall e : DfEntity, \exists x :$ Instance \bullet (Instance $\mathcal{O}(e, x)$). It worth mentioning that entities in the DataFlow Model have no attributes shown in the diagram, but theythey are presumed to correspond to identicallynamed objects (with attributes) that appear in other μ ML models.

5.3.1.3 Actor

Actors in the DataFlow Model are totally derived from Participant concept in the Task Model. It is used to represent the same concept here. Therefore, to maintain the consistency between models, an Actor is presented in the DFD using the same notation as the Task Model Participant, that is, a human stick figure, or a 3D box to represent a human, or system actor, respectively $\forall a : DfActor \rightarrow Human(a) \oplus ExSystem(a).$ A stick figure node is used to represent a human actor, as it seen in Figure 5.7 (left), whereas external systems are depicted as box nodes, taken from the UML deployment diagram, Figure 5.7 (right).

Figure 5.7: Generated DataFlow model. Actors

Each DfActor, in DataFlow Model, has only one unique identifier $\forall m : DataFlow, \forall a :$ $DfActor \bullet (a \in_m m), \exists ! i : Identifier \bullet Id(a, m, i).$

5.3.1.4 Data Flows

Data flows describe the movement of information (data drift) within a system, between its Tasks, Actors and Entities. The direction of such a flow is indicated by an arrow-head at one end (unidirectional) to represent (Read, Write, Create, Destroy (Delete), Input and Output), or at both ends (bi-directional) to represent (Update) at the initial DataFlow Diagram.

Figure 5.8: DataFlow model. Flows

Figure 5.8 (a,b) above illustrates the notation of the unidirectional and the bi-directional flow respectively, whereas (c) shows the way to express bi-directional update flows at the detailed DataFlow Diagram.

Figure 5.9: Generated DataFlow model. Create and Delete Flow

Figure 5.9 above shows the notation of two types unidirectional flows that are used in the DFD model to express Create (left) and Delete (right) flow.

Law 32. Any flow, in the *DataFlow Model*, that connects an actor to a task is an input flow, whereas any flow that connects a task to an actor is an output flow. It can be expressed as:

∀ m : DataF lowModel, f : F low, a : DfActor, t : DfT ask • (a, t, f ∈^m m) ∧ Connects(f, a, t) −→ DfInputF low(f) ∧ Connects(f, t, a) −→ DfOutputF low(f)

Law 33. Any flow, in the *DataFlow Model*, that connects a task to an entity can be either a create, a write, or a delete flow. On the other hand, any flow that connects an entity to a task is a read flow. This can be expressed in FOPL as:

$$
\forall m: DataFlowModel, f: Flow, e: DfEntity,\nt: DfTask \bullet ((e, t, f \in_m m) \landConnects(f, e, t)\n\rightarrowDfCreateFlow(f) \oplus DfDeleteFlow(f) \oplus DfWriteFlow(f)\n\land Connects(f, t, e)\n\rightarrowDfReadFlow(f)
$$

Law 34. In the *DataFlowModel*, each *Flow* links two distinct Nodes.

$$
\forall m : DataFlowModel, f : Flow \bullet (f \in_m m)
$$

$$
\rightarrow \exists (n1, n2 : Node) \bullet Connects(f, n1, n2)
$$

$$
\land (n1 \neq n2)
$$

Law 35. Any Flow in the DataFlowModel transmits some data.

$$
\forall m : DataFlowModel, f : Flow \bullet (f \in_m m)
$$

$$
\rightarrow \exists d1 : Datum \bullet (Transmits(f, d1))
$$

Law 36. Every Node in the DataFlowModel is involved in at least one Flow as a Source or a Target.

−→

 $\forall n : Node, m : DataFlowModel \bullet (n \in_m m)$

 $\exists f \bullet Flow(f) \bullet (f \in_m m) \land (Source(f, n) \oplus Target(f, n))$

Law 37. Every *DfTask, DfEntity* and *DfActor* in the *DataFlowModel* can be either a *Source* or a Target.

 $\forall t : DfTask, m : DataFlowModel \bullet (t \in_m m)$ −→ $\exists f \bullet Flow(f) \land (Source(f,t) \oplus Target(f,t))$

 $\forall e : DfEntity, m : DataFlowModel \bullet (e \in_m m)$ −→

 $\exists f \bullet Flow(f) \land (Source(f, e) \oplus Target(f, e))$

 $\forall ac : DfActor, m : DataFlowModel \bullet (ac \in_m m)$ −→

 $\exists f \bullet Flow(f) \land (Source(f, ac) \oplus Target(f, ac))$

Law 38. Each create flow creates an instance of a target entity

 $\forall m : DataFlowModel, cf : CreateFlow, t : DfTask,$ $e: DfEntity$ • ((t,e, cf ∈_m m) ∧ Connects(f,t,e)) \longrightarrow $\exists x : \text{Instance} \bullet (\text{InstanceOf}(e, x) \land \text{CreateBy}(f, x))$

Law 39. Each create flow assigns some carried data into some attributes of the target entity's instance

 $\forall m : DataFlowModel, cf : CreateFlow, t : DfTask,$ $e: DfEntity, d1: Datum \bullet ((cf, t, e \in_m m))$ \land Connects(cf, t, e) \land Transmits(cf, d1) −→ $\exists x:Instance, att1: DfAttribute$ • ((att $1 \in_{p} x$) ∧ InstanceOf(e, x) ∧ $CreateBy(cf, x) \wedge Assigns(cf, d1, att1)$

Law 40. Each delete flow deletes some instances (one or more) of a target entity

 $\forall m : DataFlowModel, df : DeleteFlow, t : DfTask,$ $e: DfEntity$ \bullet ((t, e, df ∈m m) ∧ Connects(df, t, e)) −→ $\exists x : \text{Instance} \bullet (\text{InstanceOf}(e, x) \land \text{DeletedBy}(df, x))$

Law 41. Each flow, in the *DataFlow Model*, that connects an entity to a task is a read flow

−→

 $\forall m : DataFlowModel, f : Flow, e : DfEntity, t : DfTask$ • $((e, t, f \in_m m) \land Source(f, e) \land Target(f, t))$

$$
DfReadFlow(f) \ \wedge \ \mathit{Connects}(f, e, t)
$$

Law 42. Each read flow reads one or more instance(s) of a target entity

 $\forall m : DataFlowModel, rf : ReadFlow,$ $t: DfTask, e: DfEntity \bullet ((t, e, rf \in_m m))$ \wedge Connects(rf, t, e)) −→ $\exists x, y : Instance \bullet (InstanceOf(e, x))$ $∧ InstanceOf(e, x) ∧ ReadBy(rf, x)$ \land ReadBy(rf, x) \land (x \neq y))

−→

Law 43. Each flow, in the *DataFlow Model*, that connects bi-directionally an entity to a task is an update flow

 $\forall m : DataFlowModel, f : Flow, e : DfEntity,$ $t: DfTask$ • ((e,t,f ∈_m m) ∧ (Source(f,e) $\land Target(f,t)) \land (Source(f,t) \land Target(f,e)))$

 $DfUpdateFlow(f) \wedge Connects(f, e, t) \wedge Connects(f, t, e)$

Law 44. Each update flow that connects bi-directionally an entity e to a task t, in the DataFlow *Model*, consists of a read flow that connects e to t and a *write* flow that connects t to e .

```
\forall m : DataFlowModel, uf : UpdateFlow,e: DfEntity, t: DfTask \bullet (e, t, uf \in_m m)\longrightarrow\exists rf : ReadFlow, wf : WriteFlow
                                                • ((rf, wf \in_m m) \land Breakdown(uf, rf, wf))∧ Connects(rf, e, t) ∧ Connects(wf, t, e)
```
Law 45. Each write flow updates an instance of a target entity

 $\forall m : DataFlowModel, wf : WriteFlow,$ $t: DfTask, e: DfEntity \bullet ((t, e, wf \in_m m))$ \wedge Connects(wf, t, e)) \longrightarrow $\exists x : \text{Instance} \bullet (\text{InstanceOf}(e, x) \land \text{UpdateBy}(wf, x))$ Law 46. Each write flow modifies some attibute values of the target entity's instance by some carried data.

 $\forall m : DataFlowModel, rf : ReadFlow, wf : WriteFlow,$ $t: DfTask, e: DfEntity, d1, d2: Datum$ • $((rf, wf,t,e \in_m m) \land Connects(rf,e,t)$ \land Connects(wf, t, e) \land Transmits(wf, d1) \land Transmits(wf, d2) \land (d1 \neq d2) −→ $\exists x:Instance, att1, att2: DfAttribute$ • $((att1, att2 \in p \ x) \land InstanceOf(e, x))$ \land ReadBy(rf, x) \land Assigns(wf, d1, att1) \land Assigns(wf, d2, att2) \land (att1 \neq att2))

5.3.1.5 Boundary

The final generated *DataFlow* artefact is the detailed *DataFlow Model*. It consists of a number of logical boundaries: $\forall m : DataFlowModel, \exists b : DFBoundary \bullet (b \in_m m)$. The DataFlow DfBoundary represents a business task that is decomposed into its atomic actions (tasks) within this boundary. The rectangular system boundary notation is drawn around groups of tasks and entities, to indicate that these element fall within the scope of the original business task.

5.3.1.6 Textual Expressions For Describing Data of Flows

Business users are able to annotate the generated initial DFD model using a structured textual expressions on flows. The following table (5.3) describes this:

Concept	Textual expression	Comment	
Single variable	\boldsymbol{x}	primitive naming convention οf Java datatypes	
Multiple variables	x, y, z	separated by commas	
Single object	Y	Java naming convention of objects	
Value (number)	665	numbers only without quotation	
Value (text)	"value"	any valid value within quotation	
Condition expression	$\left[\right]$	a logical operator between brackets	
Multiple conditions	[AND AND]	logical operators between brackets separated by "AND, OR " or comma	
Logical operator	$>,\leq, \leq, AND, OR$		
Attributes	@x, @y	" ^o " symbol for indicating attribute	
Projection	$\left[\ldots\right] \mathbb{Q}x$	the attribute $\sqrt[n]{\alpha}$ is projected	
Assignment	$@y = 5$	assignment operator without brackets	

Table 5.3: Textual Expressions for Describing Data on Flows

For each flow connected to an entity, the type of the return value on that flow is the type of entity, unless a projection expression exists. For instance, suppose we have an entity called Student and we want to retrieve the name of the one who has $id = 10012$. This is written as:

 $Student[@id = 10012]$ $@name$

The following expression represents how the status of a particular student is updated after completing their degree.

 $Student[@id = 10012]@status = "graded"$

5.4 Overview of the (Screen) State Model

The Screen State Model is the lowest level, **automatically-generated**, analysis representation that demonstrates, in very abstract detail, the behaviours of the system in terms of states and transitions. It is mainly based on the notion of the UML State Machine Diagram[116], which contains characteristics from both Mealy and Moore finite state machines[10]. Each state is used to indicate a current system status (displayed screen), and when an event occurs, the system will transit to the next state due to a transition triggered (labelled) by an event and/or conditions.

The model is regarded as navigational, the specifications of the GUI widgets in each screen being out of its scope. The purpose of the model is to describe, in an abstract way, the legal and illegal scenarios involved in business processes and the effects of actions within each process. As a consequence, the complete State Model expresses the most successful scenario of a process as well as all possible unsuccessful ones in order that any errors that occur during their execution can be reported and dealt with. Those scenarios occurring in a particular task are grouped together and modelled inside a logical boundary. As a result, a collection of boundaries that represent the internal behaviours of business tasks jointly form the content of the State Model.

Two kinds of errors are captured in the model, one being regarded as user-input validation and the other as related to failures caused by data source operations. The user-input error might be raised when feeding the system with values that have an unexpected data type or a null value. This issue is handled by reporting the error to the user and going back to the current state (screen). On the other hand, errors related to the back-end database system, such as connection failure, denial of access, or query/update failure, raise exceptions that notify the user about the type of error and give the option of terminating or resetting the task again.

Achieving a complete State Model leads to a comprehensive picture of the behaviours of the information system. This model can be understood by business users as being organised within a set of boundaries. Figure 9.11, 9.12 and other figures, Chapter 9, demonstrate some examples of the generated State Model including its core concepts. Each boundary represents the business task with the internal behaviours of its subtasks.

Additionally, transition labels, containing an action's name and precondition (if applicable) only, with the descriptive states name, showing whether the system is waiting for the users input or is ready to perform a database action, produce a more a readable model from an analysis perspective. The State Model, in the metamodelling, is considered a Model. Their concepts are defined using FOPL and relate to the corresponding μ ML metamodel concepts.

Law 47. The *StateModel* is a kind of *Model*.

$$
StateModel \le: Model
$$
\n
$$
\forall m \bullet StateModel(m) \longrightarrow
$$
\n
$$
Model(m)
$$

5.4.1 Graphical Notations of the State Model

In this section, each concept that appears in the *State Model* is introduced with some details about its usage and how it is visualised graphically in the model.

Figure 5.10: Generated State model. Core Elements

This model can be similar to state machine diagrams, employed in other approachs, in which each boundary of the system starts start off by an *initial* state (pseudostate) and end up by an end state. Each boundary has at most one initial and one end pseudostate. These states show only the start or end of the main business task, which do not have any internal behaviour.

Predicate	Meaning	Syntactic Suger
StateModel(x)	x is a State Model	x:StateModel
State(x)	x is a Screen State	x:State
Waiting(x)	x is a Waiting State	
Ready(x)	x is a Ready State	
Transition(x)	x is a Transition	x: Transition
Variable(x)	x is a local Variable	x:Variable
Start(x)	x is an initial pseudostate	x:Start
End(x)	x is a final pseudostate	x:End
StBoundary(x)	x is a logical boundary (region)	x:StBoundary

Table 5.4: Predicates for μ ML State Model

In order to motivate the list of policies that specifies *State Model*, two basic laws must be identified first as follows:

Law 48. Any Node in the StateModel can be either a State, Start, or End.

(State <: Node) ∧ (Start <: Node) ∧ (End <: Node) $\forall n : Node, m : StateModel \bullet (n \in_m m)$ −→

 $State(n) \oplus Start(n) \oplus End(n)$

Law 49. Any Arc in the StateModel is a Transition.

 $(Transition \leq: Arc)$ $\forall a : Arc, m : StateModel \bullet (a \in_m m)$ \longrightarrow $$

Law 50. Any boundary in the *StateModel* contains one *Start* and one *End* pseudostate, and some States and Transitions.

 $\forall m : StateModel, \forall b : StBoundary \bullet (b \in_m m)$ −→

 $\exists t : Transition, \exists s : State, \exists! init : Start,$ $\exists ! e : End \bullet (t, s, init, e \in_m b)$

5.4.1.1 State

A *State* in the *State Model* is considered a *simple state*, in respect to the UML terminology, in which it does not have substates. A State is depicted in the model, using the circle node, to represent a *current status* of the system at a time. Figure $5.10(b)$ illustrates the graphical notation of *State*. It is worth mentioning that there is no difference between the visual representation of waiting, error and ready state. Every *State* that appears in the model can be either waiting, error or ready. This can be formalised as:

> $\forall m : StateModel, \forall s : State \bullet (s \in_m m)$ −→ $Waiting(s) \oplus Ready(s) \oplus Error(s)$

State is defined as a subtype of the *Node* element (State \lt : *Node*). Thus, it is identified using a unique identifier that formalised in FOPL as:

> $\forall m : StateModel, \ \forall s : State \ \bullet \ (s \in_m m),$ $\exists ! \ id : Identifier \bullet Id(s, m, id)$

5.4.1.2 Transition

A Transition is a directed arrow drawn from the current state to the next state of a system. The UML notation of transition is adopted, an open-headed arrow, to represent the Transitions in the *State Model*. Figure $5.10(c)$ demonstrates the graphical notation of Transition. The Transition can be defined in the following theorem that follows axiom (law 10) by substituting Node with State, as well as substituting Edge with Transition. This can be written as:

 $\forall m : StateModel, \ \forall t : Transition \ \forall s1, s2 : State$ • $((s, t \in_m m) \land Connects(t, s1, s2))$ −→

 $Source(s1) \wedge Target(s2)$

5.4.1.3 Start and End Pseudostate

Two types of pseudostate are used in the *State Model*, namely *Start* and *End*. The *Start* one is placed within a logical boundary to identify the starting state of the internal behaviour of a business task. The UML notation of the initial pseudostate is used, Figure 5.10(a) above.

On the other hand, The End pseudostate is placed within a logical boundary to identify the end of the execution of a state machine. Similar to the Start pseudostate, the UML notation of the final pseudostate is used, Figure 5.10(d) above.

5.4.1.4 State Boundary

Similar to the previously presented *boundary* concepts in other μ ML models, such as Task, Impact and DataFlow models, the StBoundary element in the State Model is used to bind a complete simple state machine diagram that contains *States, Transitions* and *Pseudostates*. This captures a complete behaviour of a part of the system that falls within such a subsystem scope. StBoundary can be defined using FOPL as: $\forall m : StateModel, \exists b : StBoundary \bullet (b \in m)$.

5.5 Outlook on the Chapter

This chapter presented, in detail, the concepts, adopted graphical notations, and formal semantics for each Micro Modelling Language (μML) model appearing in the Analysis phase of BUILD. It covered all significant concepts, their utilisation, and their all possible interpretations within the various contexts of the μ ML Data Dependency, DataFlow and State Model.

In addition to this, the chapter represented how the generated intermediate system views contribute to the main goals of the thesis. More specifically, it shows how the generated models together hold a comprehensive detailed view of the system using small and semantically clean notation. Moreover, the chapter discussed how notations of the designed core elements, in each model, were reused to appear in μ ML analysis models.

The semantics and notation for each model and element are compared to similar or closely related UML concepts and notation, showing the simplicity and clarity of the proposed language. A number of unary and n-ary predicates were introduced for defining concepts and their interrelationships within each model in respect of the previously defined metamodel.

\bigcap μ ML Concepts and Notations

In the Design Phase

"Adding manpower to a late software project makes it later!" Brooks Law

6.1 Context

This chapter describes those Micro Modelling Language (μML) models that appear within the Design Phase of BUILD, namely, Database and Query (DBQ), Graphical User Interface (GUI) and Code Model. For each model, the chapter introduces, in detail, concepts, underlying ASTs and graphical notations with formally defined semantics using a set of FOPL policies and discusses how the model satisfies its purpose within the framework.

6.2 Overview of the Database and Query Model

A Database and Query Model (DBQ), shortened to a Database Model, is the lowest level, automatically-generated, model that describes, in a generic way, the common concepts and behaviours that exist in various relational database systems, such as $MySQL$, Oracle and Microsoft SQL. The model consists of one or more data schemas that hold all data and query definitions. It uses familiar database terminologies to describe the core concepts of relational database systems.

Concepts in DBQ fall into two categories, namely, Data Definition and Data Manipulation and Query concepts. The Data Definition concepts are used to describe the structure of tables, constraints that are applied to fields and special fields (keys) that represent relationships between tables. A schema element is used to define the logical database schema that consists of a number of persistent tables that contain some atomic columns, defined using a DBQ Table and Column concept, respectively.

The Table element reflects the definition of a database table, that is, the table type definition with its associated column specifications, as found in a real database system. It has a number of columns and special columns (a primary key and possible foreign keys) represented by Column, $PrimaryKey$ and $ForeignKey$ concepts, respectively. DBQ Columns are considered to be atomic elements in the model that cannot be further decomposed. Together they define the type of each instance (new row) of the table they belong to.

It is worth mentioning that relationships between business entities are also expressed at this level either by foreign keys or separate tables (Linkers). The decision to convert relationships into either a table or a special column is made according to the application of a transformation rule. This is discussed separately in Chapter 7 and 8.

In addition to the *Data Definition* concepts, *Data Manipulation and Query* concepts, which are used for constructing appropriate queries about a domain, are also described in a declarative way. For example, *Trigger* and *Procedure* elements are used to define, respectively, triggers and pre-defined queries that exist in a particular schema. The trigger is used to capture *range constraints* applied to some columns of a particular table. As a result, the generated SQL script includes a BEFORE INSERT trigger attached to the table to enforce constraints that prevent or allow new data to be inserted.

The format for query expressions is regarded sufficiently general in order to be translated into relational model (SQL) for various database vendors. The declarative style of describing database queries, including the *Projection* (π) , Selection (σ) and Join (∞) operations, is inspired by Functional Algebra, as it discussed later in section 6.2.2. The following example illustrates the expression of a query using the DBQ query syntax, Relational Algebra syntax, and natural language (English).

The query: "Retrieve the name of a module that has a code equals to COM6006" can be expressed declaratively using Relational Algebra as: $\pi_{name}(\sigma_{code='COM6006'}(Module));$. The following listing (6.1) illustrates the equivalent DBQ query expression. *Query* as a higher-order function is defined. It takes *Project* as a first argument and a *Relation* as a second argument, to which the Project is applied.

Listing 6.1: The structure of the DBQ Query Expression

```
1 <dbs:Query id="14" name="Reteive_Module_Name">
2 <dbs:Project id="15">
3 <dbs:Column id="15" name="name" size="0" prefix="Module">
4 <mod:Type id="16" name="Integer"/><br>5 </dbs:Column>
         5 </dbs:Column>
6 </dbs:Project>
7 <dbs:Relation id="17" name="Module">
8 <dbs:Filter id="18">
9 <dbs:Relation id="19" name="Module"/>
10 <dbs:Operator id="20" type="boolean" symbol="equals">
11 <dbs:Column id="21" name="code" size="20">
12 <mod:Type id="22" name="INTEGER"/>
13 </dbs:Column>
14 <dbs:Literal id="23" type="VARCHAR" value="COM6006"/>
15 </dbs:Operator>
16 </dbs:Filter>
17 </dbs:Relation>
18 </dbs:Query>
```
The query language starts from the assumption that the results of all operations on tables, e.g. Filter (listing 6.1 above) and *Join*, are treated as *Relations*. It can be said that the Relation element is an abstracted term of Table. This means that any Relation in DBQ refers to either an actual database table or else filtered or joined tables.

As a consequence of using the notion of Relation in the model, direct navigation through the data is clearly expressed in a unified way for all operations. The concrete path to data within Relations is used without the necessity to be computed during transformation.

The Database Model, in the proposed metamodelling hierarchy, is considered a kind of *Model.* Thus, conepts of the DBQ model can be defined using FOPL and relate to the corresponding μ ML metamodel concepts.

Law 51. The *DatabaseModel* is a kind of *Model*.

$DatabaseModel$	$<$	$Model$
$\forall m \cdot DatabaseModel(m)$	\longrightarrow	
$Model(m)$		

6.2.1 Notation and Semantics of the Database and Query Model

In this section, each concept that appears in the *Database Model* is discussed in detail. This includes its usage and how it is visualised graphically in the model. It is worth emphasising that a number of UML notations from the Class Diagram are adopted and reused with a slightly different meaning compared to the usual UML semantics. The following table (6.1) summaries all predicates used in describing μ ML Database and Query Model.

CompositePK(x)	x is a DBQ Composite	x:CompositePK
	Primary Key	
CompositeFK(x)	x is a DBQ Composite	x:CompositeFK
	Foreign Key	
For eigenKey(x)	x is a DBQ Foreign Key	x: For eigenKey
PartOfKey(x, y)	x is a part (column) of a	
	Composite Key y	
Create(x)	x is a DBQ Insert operation x:Create	
Query(x)	x is a DBQ Select operation	x:Query
Update(x)	x is a DBQ Update operation	x:Update
Delete(x)	x is a DBQ Delete operation	x:Delete
Defines(x, y)	a column x defines the	
	characteristics of each row uniquely	
	in table y	
References(x, y, z)	a column x refers table y	
	to table z	
References(x, y)	the relation x refers to the table y	
Join(x)	x is a join operation	x:Join
Joins(x, y, z)	x is a join operation between	
	relation y and relation z	
Filter(x)	x is a filter function	x:Filter
Project(x)	x is a project function	x:Project

Table 6.1: Predicates for μ ML Database and Query Model

In order to motivate some policies that specify the DBQ Model, basic declarations of core elements must be identified first. According to the μ ML metamodel hierarchy where DBQ Model table is a subtype of *Node* as: $(Table \leq: Node)$, then any *Node* in the *Database Model* or schema represents a logical table: $\forall n : Node, m : Schemea \bullet (n \in_m m) \rightarrow Table(n).$

In addition to this, the *Procedure* element in the DBQ Model is declared in the μ ML metamodel as a Function element, in which the DBQ Procedure is a subtype of Function, (Procedure \leq : Function \leq : Expression). Therefore, any Function in the model represents a stored procedure: $\forall p : Function, s : Schema \bullet (p \in_m s) \rightarrow Procedure(n).$

6.2.1.1 Schema

A Schema is the metadata for a database, consisting of data and query definitions. According to the μ ML metamodel, the DBQ *Schema* element is also considered a subtype of the *Node* element (Schema \leq : Node). Therefore, it is identified using a unique identifier. This can be defined in FOPL as:

> $\forall s : Schema, \forall d : DatabaseModel \bullet (s \in_m d),$ $\exists ! \ id : Identifier \bullet Id(s, id)$

A *Schema* is constructed by considering the internal structure of business entities $(Table)$, e.g. named and typed columns, and the relationships between entities that are represented via key columns ($PrimaryKey$ and $ForeignKey$), and finally the various kinds of pre-defined queries (Procedure and Trigger) that will be executed over the data.

Law 52. In the *DatabaseModel*, each node in the schema is a table.

$$
\forall s: Schema, d: DatabaseModel, \forall n:Node
$$

$$
\bullet ((n \in m s) \land (s \in m d))
$$

$$
\rightarrow
$$

 $Table(n)$

Law 53. In the *DatabaseModel*, each function in the schema can be either a trigger or a stored procedure.

$$
\forall s: Schema, d: DatabaseModel, \forall p: Function
$$

\n• ((p $\in_m s$) \land (s $\in_m d$))
\n \longrightarrow
\n<

6.2.1.2 Relation

Based on the concept of a Relation in Relational Databases and also in the Relational Algebra, in which everything is considered a relation, the DBQ Relation element serves the same purpose. Whereas Table represents the type of a database table, Relation represents the actual table (the collection of rows), and, recursively, the result of performing relational algebra operations on relations. It is used for binding some internal database operations (functions), such as *filter* and join and presenting their returned result as a Relation with a given name.

$$
\forall s: Schema, \forall r: Relation \bullet (r \in_m s),
$$

$$
\exists! ((f: Filter \bullet (Child(r, f)) \oplus (j:Join \bullet (Child(r, j)))
$$

In addition to this, the DBQ Relation is also used to refer to the collection of rows in a table, defined in another DBQ Table element represented within the same database schema. This interpretation is intended when when the Relation node has no descendants.

> $\forall s : Schema, \forall r : Relation \bullet ((r \in_m s) \land (Child(r, \emptyset))),$ $\exists ! t : Table \bullet ((t \in_m s) \land (References(r, t)))$

6.2.1.3 Table

As the intention for DBQ is to express low-level logical schemas that can be directly implemented in a relational database system, business entities that appeared previously at the Requirement Sketching and Analysis phases are declared at this level as Table types. Each Table reflects the usual notion of a database table and its properties as found in real-world database systems. Figure 6.1 illustrates the structure of the DBQ table.

Figure 6.1: Generated Database and Query model. Structure of a Table

According to the μ ML metamodel, DBQ Table is also considered a subtype of the Node element (Table \leq : Node). Therefore, it is identified using a unique identifier. This can be defined in FOPL as:

$$
\forall s: Schema, \forall t: Table \bullet (t \in_m s),
$$

$$
\exists! id: Identityier \bullet Id(t, id)
$$

DBQ Tables always contains one or more columns that describe the shape of the data stored in the table: $\forall t : Table, \forall c : Column \bullet (c \in p t)$. Additionally, tables have some key columns: $\forall t : Table, \forall f k : ForeignKey \bullet (fk \in p t).$ The description of *Column, PrimaryKey*, CompositePk and ForeginKey are presented in the following sections.

Law 54. In the *DatabaseModel*, each table has a primary key.

$$
\forall s: Schema, \forall t: Table, c: Column
$$

\n• ((c $\in_{p} t$) \land (t $\in_{m} s$) \land *Defines(c,t)*)
\n \longrightarrow
\n*PrimaryKey(c)*

Law 55. In the *DatabaseModel*, a table has only one composite key.

$$
\forall s: Schema, \forall t: Table, (c1, c2): Column
$$

\n
$$
\bullet (((c1, c2) \in_p t) \land (t \in_m s)
$$

\n
$$
\land (Defines(c1, t) \land Defines(c2, t)) \land (c1 \neq c2)
$$

\n
$$
\Rightarrow
$$

\n
$$
\exists! ck: CompositeKey \bullet (ck = c1 \cup c2)
$$

6.2.1.4 Column

The Column element is used to define the actual columns of a table at DBQ. Additional features exist, which include the ability to declare a recognised datatype for each column's value, such as VARCHAR, INTEGER and DATE:

 $\forall c : Column, \forall t : Table \bullet (c \in_p t), Type(c, VARCHAR) \oplus Type(c, INTEGRER) \oplus Type(c, DATE)$

Moreover, at this level, a column can be defined whether it is special or not. Some columns are also foreign keys that represent relationships or primary keys that represent a structure of a table:

 $∀ c : Column, ∀ t : Table \bullet (c \in p t)$ \longrightarrow $(PrimaryKey(c) \vee ForeignKey(c)) \vee Column(c)$

6.2.1.5 Primary Key

A Primary Key is regarded an essential feature of database tables. Each DBQ Table has one that might be either manufactured (automatically generated) or specified manually by a user. This can be defined using FOPL as: $\forall pk : PrimaryKey, \forall t : Table \bullet ((pk \in p t) \land (Auto(pk) \oplus$ $\neg \text{Auto}(pk))$. The following figure (6.2) shows the visual representation of *Primary Key* within a table, which is the bold underlined attribute.

Figure 6.2: Generated Database and Query model. Table's Structure

Law 56. In the *Schema* (DBQ) Model, each primary key in such a table defines each row of that table.

$$
s: Schemea, \forall pk: Column, \forall t: Table
$$
\n
$$
\bullet ((pk \in_{p} t) \land (t \in_{m} s) \land Defines(pk, t)) \rightarrow
$$
\n
$$
\rightarrow
$$
\n
$$
PrimaryKey(pk)
$$

In some cases, the table might contain two or more columns that identify the uniqueness of its rows. Therefore, these columns are captured and treated as a composite primary key, (Compos*iteKey*), of that table.

Law 57. In the *DatabaseModel*, a composite key defines a table if that table has more than one primary key.

$$
s: Schema, \forall c1, c2: Column, \forall t: Table
$$
\n
$$
\bullet ((c1, c2 \in p \ t) \land (t \in m \ s)
$$
\n
$$
\land (Defines(c1, t) \land Defines(c2, t))
$$
\n
$$
\land (c1 \neq c2))
$$
\n
$$
\rightarrow \exists! ck: CompositeKey \bullet (PartOfKey(c1, ck)
$$
\n
$$
\land PartOfKey(c2, ck) \land Defines(ck, t))
$$

6.2.1.6 Foreign Key

Relationships among business objects that are previously expressed, using graphical notation, in Information and Data Dependency model appear here as a number of Foreign Key columns attached to Tables. It refers to the other table in the relationship. Figure 6.2 above illustrates the visual representation of foreign keys, which is the underlined attribute.

Law 58. In the *Schema* (DBQ) Model, each foreign key in such a table refers to another table.

$$
s: DatabaseModel, \forall f k: ForeignKey, \forall t1, t2: Table
$$

\n
$$
\bullet ((t1, t2 \in_m s) \land (fk \in_p t1))
$$

\n
$$
\rightarrow
$$

\n
$$
References(fk, t1, t2)
$$

6.2.2 The Query Language and Functional Algebra

The query language in DBQ is inspired by two types of Algebra, namely, Relational Algebra and Functional Algebra. Relational Algebra is a procedural query language that applies particular operators to one or more relations. Additionally, Functional Algebra implies a higher-order functional style, in which functions accept functions as arguments. The following subsections discuss every database operation (function) taking into account how each one relates to *Algebra* concepts.

6.2.2.1 Filter

A Filter function is one of the main operations of the algebra that maps a predicate over a set and returns a subset of values which pass the predicate test. This is similar to a filter function in Functional Algebra, which accepts a predicate and a list as arguments, returning a filtered list as the result. Filter can be defined using the following signature:

$$
filter(pred: T \rightarrow \; Bool, Set: Set[T]): Set[T].
$$

Regarding to the DBQ query language, a Filter element, which is equivalent to filter in a Functional Algebra, is considered parent of two child elements, namely, Relation and Operator:

> $\forall m : DatabaseModel, \ \forall f : Filter \bullet (f \in_m m),$ $\exists r : Relation, op : Operator \bullet (Child(r, f) \wedge Child(op, f))$

The Operator is noted as a boolean function (Operator \leq : Function) that has comparison operators, such as greater than $(>)$, less than $($ and equals $(=)$, and two arguments to be examined. The arguments are expressed as *child* nodes of *Operator* element in DBQ, which might be Column, Variable or Literal. Listing 4.1 examplifies the utilisation of Filter, in which a substantial predicate is provided to return a subset of Module records (filtered).

6.2.2.2 Project

The second core function in the proposed DBQ query language is the Project function that projects out a set of column values, determined by a function, from a set of records. Similar

to the Filter, the Project function corresponds to a map function in Functional Algebra, which accepts a function over every element in the input list. Project can be defined using the following signature:

 $project(func: T \rightarrow U, Set: Set[T]): Set[U].$

In the DBQ query language, a *Project* element is utilised to express the *project* opration that is equivalent to project in a Functional Algebra. The DBQ Project node consists of either some Columns or Relations nodes as descendants. This makes Project is not a complete query expression in the DBQ language, as the relation argument is missing, but Project is considered a function to be applied to a relation in a higher-order way, by a DBQ Query operator. However, the *prefix* name, which is an *attribute* of the *Column* element, is used to label the expected type of the second (implicit) argument.

> $\forall m : DatabaseModel, \ \forall p : Project \ \bullet \ (p \in_m m),$ $\exists c : Column, r : Relation \bullet (Child(c, f) \oplus Child(r, f))$

On one hand, when Project has a number of Columns as descendants, it means these columns are projected from the set declared in their *prefix* attribute. On the other hand, when the *child* of Project is a Relation node, it means that whole columns (tuples) are projected, which is equivalent to $(*)$ in SQL. Listing 6.1 examplifies the utilisation of *Project* element, in which it has only one *child* element (*Column*) to be projected from a set of *Modules* as declared in its prefix attribute.

6.2.2.3 Query

The DBQ Project and Filter functions are often used together to to return data of interest after a search. Therefore, another high-order function might be introduced, *query* that takes a *project* function and a set and returns a selected set. The following signature defines the *Query* function:

$$
query(func: T \rightarrow U, Set: Set[T]): Set[U]
$$

The search operation is represented in DBQ via the Query element that consists of a Project and Relation element. In the case that the Relation is filtered, it will contain a Filter element as *child.* Listing 6.1 examplifies the use of *Query* element, where as the following FOPL specifies the Query element:

> $\forall m : DatabaseModel, \ \forall q : Query \bullet (p \in_m m),$ $\exists p : Project, r : Relation \bullet (Child(p, q) \wedge Child(r, q))$

6.2.2.4 Create

Create is a usual set- theoretic operation for adding new instances (rows) into a dataset (insert). It is a destructive operation that modifies the base tables (dataset) by instantiating a new distinct instance (object) of that table and inserting a new record into it. It is obvious that there is a predicate to prevent inserting duplicated records. The uniqueness of records is based on the values of all fields and equality is judged on the basis of the declared primary keys.

From this, the *Create* function can be defined as a function that takes two arguments: an original dataset (*base*) of some type (T) , and a new tuple (*extra*) of the same type T that is going to be inserted into the *base*. The function returns the number of inserted tuples (nat) , which is a Natural number (nat $\in \mathbb{N}$), and has the side-effect of a modified dataset (base'). The following signature defines the Create function:

 $create(base : Set[T], extra : Tuple[T]) : nat, base'$

In the DBQ Create element, a number of assignment operator statements, Operator elements with symbol="assign", are used, as descendant nodes of Create, for each item in the new tuple that is going to be adding to the dataset. The following listing (6.2) demonstrates the structure of the DBQ Create element. It is clear to see that Create has a list of Operator elements that are assignment instructions (symbol = "assign") to insert every item of the new tuple, supplied as Variables, into the equivalent column in the dataset.

Listing 6.2: The structure of the DBQ Create Expression

```
1 <dbs:Create id="28" name="INSERT_Module">
2 <dbs:Operator id="29" type="boolean" symbol="assign">
3 <dbs:Column id="30" name="code" size="6">
4 <mod:Type id="31" name="INTEGER"/>
5 </dbs:Column>
6 <dbs:Variable id="32" type="INTEGER" name="code"/>
7 </dbs:Operator>
8 <dbs:Operator id="33" type="boolean" symbol="assign">
9 <dbs:Column id="34" name="title" size="20">
10 <mod:Type id="35" name="VARCHAR"/>
11 </dbs:Column>
12 <dbs:Variable id="36" type="VARCHAR" name="title"/>
13 </dbs:Operator>
14 <dbs:Operator id="37" type="boolean" symbol="assign">
15 <dbs:Column id="38" name="credit" size="2">
16 | mod:Type id="39" name="VARCHAR"/>
17 </dbs:Column>
18 <dbs:Variable id="40" type="VARCHAR" name="credit"/>
19 </dbs:Operator>
20 </dbs:Create>
```
6.2.2.5 Update

The Update function performs an update operation on every instance of a dataset. In algebra, we assume that it is possible to have functions with side effects. The Update function iterates over every row of a dataset and applies an arbitrary re-assignment operation to each record that probably satisfies a predicate. It is assumed that the dataset is modified as a side-effect of Update. The following signature defines the Update function:

 $update(pred: T \rightarrow U, Set: Set[T]) : nat.$

It is common in various database systems that the SQL UPDATE statement returns a Natural number, indicating the number of records that were modified. In the proposed DBQ Update, a number of assignment operator statements are used to express modification to express modifications to column data without considering an expression to present the return value of Update.

The following listing (6.3) demonstrates the structure of the DBQ Update. It is clear to see that Update has an Operator of the type assignment (symbol = "assign") to modify the value of the column *title* by the value of the variable *title*.

Listing 6.3: The structure of the DBQ Update Expression

```
1 <dbs:Update id="40">
2 <dbs:Operator id="41" type="Boolean" symbol="assign"><br>3 <dbs:Column id="42" name="title" size="20">
             3 <dbs:Column id="42" name="title" size="20">
4 <mod:Type id="43" name="VARCHAR"/><br>
</dbs:Column>
5 </dbs:Column><br>6 </dbs:Variable
             6 <dbs:Variable id="44" type="VARCHAR" name="title"/>
7 </dbs:Operator>
8 </dbs:Update>
```
6.2.2.6 Delete

The *Delete* is another common data manipulation operation for removing existing records from a dataset. Similar to Create, the Delete operation is defined as a destructive operation that modifies a base tables (dataset) by removing records and reducing its size. As Delete function removes some records from a dataset, it is clear that there is a predicate that must be applied to every record in the dataset. Only records that pass the predicate are removed. That is why Delete is often used with a independant sub-operation *Query* with a Filter to achieve this.

From this, the *Delete* operation can be defined as a function that takes two arguments, an original dataset (*base*) of the type T , and a projected tuple, of the same type T from the performed *filter* (sub-operation). The function returns the number of removed tuples(*nat*), which is a Natural number (nat $\in \mathbb{N}$), has the side-effect of a modified dataset (base'). The following signature defines the *Delete* function:

```
delete(base : Set[T], to go : Tuple[T]) : nat, base'
```
Regarding the DBQ *Delete* element, a *child* element *Relation* is used to wrap the *Filter* function (sub-operation) and represents its return result set to the Delete function. This filtered Relation will be deleted (extracted and remove) from the original Relation defined inside the Filter element. The following listing (6.4) demonstrates the structure of the DBQ Delete.

Listing 6.4: The structure of the DBQ Delete Expression

```
1 <dbs:Delete id="43" name="DELETE_Code">
2 <dbs:Relation id="44" name="Delete_Module">
3 <dbs:Filter id="45">
4 <dbs:Relation id="46" name="Module"/><br><dbs:Operator id="47" type="boolean"
              5 <dbs:Operator id="47" type="boolean"
6 symbol="equals">
7 <dbs:Column id="48" name="code" size="6">
8 <mod:Type id="49" name="INTEGER"/>
9 \langle dbs:Column>
10 <dbs:Variable id="50" type="INTEGER"
11 \vert name="code"/><br>12 </dbs:0perator>
              12 </dbs:Operator>
13 </dbs:Filter>
14 </dbs:Relation>
15 </dbs:Delete>
```
6.2.2.7 Join

The *join* is a database operation (∞) that computes the inner join of two tables, resulting in a merged table, based on a constraint between the values of two columns, one from each table. It is considered a binary *cartesian product* operation (\times) with a selection condition (θ). The *join* concept can be expressed using algebra as: $\sigma_{\theta}(p \times q)$, where p and q are two relations. Therefore, the resulting *join* outcome is another relation that consists of columns from both relations p and q. The following signature defines the Join function:

 $join(Set : Set[T], Set : Set[U]) : Set : JSet[W]$

6.2.3 The Significance of the Database and Query Model

The Database and Query Model is a key model in μ ML because it describes a comprehensive representation all generic data schema specifications, required for any relational database generation. This platform-independent model is used as a source model in the code generation approach to produce executable database schema code in various target database systems.

6.3 Overview of the Graphical User Interface (GUI) Model

A Graphical User Interface Model (GUI) is the lowest level, **automatically-generated**, structural model that describes, in an abstracted way, the detailed GUI architecture that is common in various enterprise system environments. The model contributes to the development of system GUIs by representing the concrete structure of their controls (widgets) that can later generate simple and user-friendly interfaces for their system. Familiar terminologies for describing basic concepts of GUI are used to form the declarative GUI language. The model is considered to be platform- independent as no implementation details are presented in the model.

Much like the modelling strategy adopted in μ ML, each group of application windows that falls within a scope of a business practice is modelled as a boundary to represent a separate part of the system. Thus, it can be said that the GUI Model consists of a number of independent boundaries expressed via Boundary nodes. As this is considered during the early development phases in BUILD, a business task may, by this stage, already have been broken down into atomic tasks (each task having an atomic action). It is assumed here that each atomic task requires a user interface to support human-computer interaction throughout the execution of a business process. As a consequence, an appropriate screen design appears in the model for each atomic task that appeared previously in the (Screen) State Model.

It can be said that the Boundary contains a collection of Window nodes that represent IS screens. Each window in a system boundary is prioritised to be linked to another one, forming an ordered sequence or branches (choices) of screens. The screen order is based on the priority number attached to each *Window* node, the window with the highest priority score appearing before those with lower ones to reflect the internal behaviour (logic) of a business process. It is worth mentioning that the priority score is allocated for each window, except error ones, by a previous model transformation step for generating the State Model, in which the priority for each *state* is passed to the corresponding *window*, as discussed later in Chapter 7.

Generally, information system screens are classified into two major types: entry form and output report screen. The entry form is a screen that is waiting for an external input/event from a system actor. This appears in the GUI Model as Waiting Windows with a main button that triggers an *input* action when a *click* event occurs. The *output report* is a ready screen for executing an atomic business task, such as retrieving or deleting information from the system. It appears in the model as *Ready Windows* with a main button that initiates either create, read, update, delete or output *(display)* actions. It simply notifies the user that an action is going to be initiated or displays some output message to them.

In addition to this, a failure of a business task is captured by presenting an appropriate Error Reporting Window that links to each window in a successful scenario. An Error Reporting Window can be defined as a special kind of Reporting Window that is used to display (report) a meaningful error message to the user. An Exception event that is triggered when any error occurs is attached to every window. As a result, an error reporting window may appear during any step in a scenario.

In fact, the GUI model has one kind of node to represent the window concept, the Window node. The type of window, whether it is *Waiting* or *Ready*, is extracted from the generated name of each Window. A particular naming convention is adopted whereby the name of a window must include its type, such as *Input_title_Waiting* or *Input_title_Waiting_error*.

In typical enterprise applications, business-users visualise and gain access to the backend database of the system by interacting with a number of GUIs to execute some business actions in order to retrieve or manipulate data. This is known as a presentation layer in the traditional architecture of any information system. The widgets in the GUI Model hold the required information about the target database tables and fields, with their properties and constraints, along with corresponding data that is shown via these screens. This offers a consistent mapping between the presentation and database layers and leads to the correct visualisation of a window.

The *GUI Model*, in the produced metamodelling hierarchy, is considered a kind of *Model*. Thus, concepts of the GUI Model can be defined using FOPL and relate to the corresponding μ ML metamodel concepts.

Law 59. The *GUIModel* is a kind of *Model*.

$$
GUIModel \le: Model
$$
\n
$$
\forall m \bullet GUIModel(m) \longrightarrow
$$
\n
$$
Model(m)
$$

6.3.1 Notation and Semantics of the Graphical User Interface (GUI) Model

In this section, each concept that appears in the GUI Model is introduced with some details about its usage and how it is visualised graphically in the model. The following table (6.2) summaries all predicates used in describing μ ML GUI Model.

Table 6.2: Predicates for μ ML GUI Model

In order to motivate some policies that specify the GUI Model, basic declarations of core elements must be identified first. According to the μ ML metamodel hierarchy where GUI Model window is a subtype of *Node* as: (Window $\lt:$: *Node*), any *Node* in the *GUI Model* represents a GUI window:

$$
\forall n : Node, m : GUIModel \bullet (n \in_m m)
$$

$$
\longrightarrow
$$

$$
Window(n)
$$

Additionally, various GUI controls are considered Widget concept in the metamodel. According to the μ ML metamodel hierarchy, the *Widget* element is denoted a subtype of *Node* as: $(Widqet \leq: Node)$, any *Node* in the GUI window represents a GUI control:

$$
\begin{array}{l} \forall\; n : Node, w : Window\;\bullet\;\left(n\;\in_{c}\;w\right)\\ & \longrightarrow \\ \hspace{2.5cm} \hspace{2.2cm} \longleftarrow\\ \hspace{2.2cm} \hspace{2.2cm} \hspace{2.2cm} \hspace{2.2cm} \longleftarrow\\ \hspace{2.2cm} \hspace{2.2cm} \hspace{2.2cm} \longleftarrow\\ \hspace{2.2cm} \hspace{2.2cm} \hspace{2.2cm} \longleftarrow\\ \hspace{2.2cm} \hspace{2.2cm} \longleftarrow\\ \hspace{2.2cm} \longleftarrow
$$

Law 60. In the *GUIModel*, each widget in any window can be either a button, textfield, or label.

> $\forall c: Widget, \ \forall w: Window, m: GUIModel$ \bullet ((w \in_m m) \wedge (c \in_c w)) −→ $Textfield(c) \oplus Button(c) \oplus Label(c)$

6.3.1.1 GUI Window

The *Window* element is used to capture every information system screen. According to the μ ML metamodel, GUI Window is considered a subtype of the Node element $(Window <:Node).$ Therefore, it is identified using a unique identifier. This can be defined in FOPL as:

> $\forall m : GUIModel, \forall w : Window \bullet (w \in_m m),$ $\exists !\ id : Identifier \bullet Id(w, id)$

Every *Window* cosists of many GUI controls; these are all *Widget* elements. This can be expressed as:

> $\forall w : Window, m : GUIModel \bullet (w \in_m m),$ $\exists c : W \text{idget} \bullet (c \in_c w)$

Extra features are added to specify each window. For instance, the priority score of a window is declared, for each Window element, using the order attribute. Moreover, the declaration of Error Window is determind by the boolean attribute error attached the Window node.

Listing 6.5: The structure of the GUI Window node

```
1 <qui:Window id="22" name="Input code Waiting" order="4">
2 <gui:Textfield id="23" name="code" size="6"/><br>3 <gui:Button id="24" name="Exception" event="E
       3 <gui:Button id="24" name="Exception" event="Exception"
4 exit="false" hidden="true"/>
5 <gui:Button id="25" name="Input" event="Input"
6 exit="false"/>
7 \times /gui:Window>
```
The previous listing (9.8) illustrates a structure of a window for receiving *code* value from the user that consists of a Textfield and a Button to read the input value. However, it is worth mentioning that the button that fires the *Exeption* is invisible (*hidden =true*).

6.3.1.2 GUI Textfield

The Textfield element is used to capture every user interface text field control that appears in a window. It allows the actor to input textual information to be used by the system. The additional specifications of each text field are determined by some attributes attached to a Textfield node, such as the name and size attributes. Listing 9.8 above exemplifies the GUI Textfield element including its attributes. According to the μ ML metamodel, GUI Textfield is considered a subtype of the *Widget* element (*Textfield* \lt : *Widget*). Therefore, it is identified using a unique identifier. This can be defined in FOPL as:

> $\forall tf: Textfield, \forall w: Window \bullet (tf \in_c w),$ $\exists !\ id : Identifier \bullet Id(tf, id)$

6.3.1.3 GUI Label

The label is a user interface control that is used to display textual information on a window. It is represented in the GUI Model via the Label element that is used to capture every label appears in a screen. Label is also specified by some attributes attached to it, such as the name and text attribute. The text attribute the actual text that appears on a window. According to the μ ML metamodel, GUI Label is considered a subtype of the Widget element (Label $\lt:$: Widget). Therefore, it is identified using a unique identifier. This can be defined in FOPL as:

> \forall leb : Label, \forall w : Window • (leb \in_c w), $\exists !\ id : Identifier \bullet Id(left, id)$

6.3.1.4 GUI Button

The button is another kind of user interface control that enables the user a direct way to trigger an event. It is appeared in a window as a rounded rectangle with longer width than its hight, and a text in its middle. Every button in such a window is expressed in the GUI Model through the *Button* element. The GUI *Button* describes many features in regard to the button using a number of attributes, such as, *name* and *event* attribute that define the button's name and the type of event triggered by this button. Listing 9.8 above examplifies the GUI Button element including its attributes. According to the μ ML metamodel, GUI *Button* is considered a subtype of the Widget element (Button $\lt:$: Widget). Therefore, it is identified using a unique identifier. This can be defined in FOPL as:

> $\forall b : Button, \forall w : Window \bullet (b \in_c w),$ $\exists ! \ id : Identifier \bullet Id(b, id)$

6.3.1.5 GUI Boundary

Similar to the previously presented *boundary* concepts in other μ ML, such as Task, Impact, DataFlow and State, the GUIBoundary node in the GUI Model to bind one or more Window elements. This indicates that some windows fall within a particular subsystem or (scope) to be developed, while other windows fall within another scope. It can be defined using FOPL as: $\forall m : GUIModel, \exists b : GUIBoundary \bullet (b \in m).$

6.3.2 The Significance of the GUI Model

As the Micro Modelling Language (μML) seeks simplicity for designers, business-users are relieved from the task of designing the kinds of screens used, since the decision to create a waiting

or ready window is determined automatically, based on information provided in the Screen State Model. Additionally, some database tables might might be automatically assigned as data sources, when required, in some screens without any end-user manual modification. Consequently, reasonable chunks of boilerplate code for managing the database connection and exception handling may be manufactured directly at the code generation stage. Moreover, the fact that this is all done automatically provides type crosschecking to preserve consistency between the information system layers.

6.4 Overview of the Code Model

The Code Model is the lowest level, **automatically-generated**, AST representation that describes the abstract specification for the Object-Oriented (OO) code of the system in terms of business entities and tasks (classes), their methods, attributes and expressions. It aims to present a comprehensive business logic layer that determines how to manipulate the data in either a separate layer or one combined with the user interfaces, that is, either in a 3-tier or 2-tier architecture, respectively. At the current version of BUILD, it is worth mentioning that the model is NOT a part of the transformation chain for generating complete 2-tier applications. It might be used in later versions for developing 3-tier IS applications.

Following the code structure in the OOP languages, the *Code Model* consists of a number of $Clazz^1$ elements that represent the actual OO classes of a system. Each $Clazz$ node typically has some *Constructor, Attribute* and *Method* nodes as descendants to express the actual class attributes, constructors and methods at this level of detail.

Apart from this, this model is platform independent and intended to capture sufficient details to generate idiomatic constructs of executable code for a target OOP language, e.g. Java, along with adequate chunks of *boilerplate* code to establish and manage consistent protocols between the invocation of methods and database predefined queries in the system layers.

Additionally, the model is completely derived from both the detailed *Database and Query* Model and the DataFlow Model without any direct contribution from the end-user. This achieves one of the goals of BUILD, by avoiding end-users from modelling these technical specifications about their system. All modelling activities are done at a higher level of abstraction. As a consequence, one goal of our approach is achieved by relieving users from writing these technical specifications users from these technical specifications regarding their system and producing it at a higher level of abstraction. Business-users, the designers, make no contribution to the detailed architecture of their software system as our approach adopts a common architecture and is able to generate all kinds of systems.

The classes are divided into two main categories, namely, entity classes and process (task) classes. The entity classes are translated from the equivalent tables and exist in the Database and Query Model, whereas the process classes are derived from boundaries that appear in the DataFlow Model. Relationships between classes, such as composition and inheritance, are also extracted from the Database and Query Model and might be implemented in the Code Model to represent a consistent business logic layer.

¹The name *Clazz* is chosen for this node to avoid a name-clash with the built-in type "Class" in Java.

The Code Model, in the produced metamodelling hierarchy, is considered a kind of Model. Thus, concepts of the Code Model can be defined using FOPL and relate to the corresponding μ ML metamodel concepts.

Law 61. The CodeModel is a kind of Model.

$$
Code Model \ \text{&}: \ Model
$$
\n
$$
\forall m \ \bullet \ CodeModel(m)
$$
\n
$$
\longrightarrow
$$
\n
$$
Model(m)
$$

6.4.1 Notation and Semantics of the Code Model

In this section, each concept that appears in the *Code Model* is introduced with some details about its usage and how it is expressed in the model. The following table (6.3) summaries all predicates used in describing μ ML Code Model.

Table 6.3: Predicates for μ ML Code Model

In order to motivate the list of policies that specifies the Code Model, two basic laws must be identified first as follows:

Law 62. Any node in the CodeModel is a class.

$$
Code Model \le: Model
$$

$$
\forall m: CodeModel, n:Node \bullet (n \in_m m)
$$

$$
\longrightarrow
$$

$$
Clazz(n)
$$

Law 63. Any function in the CodeModel is a method in a class.

$$
Code Model \le: Model
$$

$$
\forall m : CodeModel, f : Function, c : Class \bullet (c \in_m m)
$$

$$
\longrightarrow
$$

$$
Method(f) \land (f \in_p c)
$$

6.4.1.1 Class

The *Clazz* element, in the Code Model, is used to define object-oriented class types. According to the μ ML metamodel, *Clazz* is considered a subtype of the *Node* element (*Clazz* \lt : *Node*). Therefore, it is identified using a unique identifier. This can be defined in FOPL as:

```
\forall m : CodeModel, \forall c : Class \bullet (c \in_m m), \exists ! \ id : Identifier \bullet Id(c, id)
```
A Clazz element consists of a number of attributes to describe its features and methods to describe its operations, expressed using *CAttribute* and *CMethod* nodes as descendants of *Clazz*. This can be formalised in FOPL as:

> $\forall m : CodeModel, \ \forall c : Class \bullet (c \in_m m),$ $\exists ((a : CAttribute \bullet (a \in p \ c) \land Child(c, a))$ \wedge $(f : CMethod \bullet (f \in_{p} c) \wedge Child(c, f)))$

An extra feature to specify the accessibility of a class (*public, private* and *protected*) is considered in the Code Model language using an XML attribute attached to the Clazz element $(visible="public")$.

6.4.1.2 Attribute

A CAttribute element is used to declare every fields of a class in the Code Model. It is an atomic node in the model that has no further childern. Additional features that specify the attribute are attached, as XML attributes, to the CAttribute node, such as name, visible and type. The CAttribute elements appear in the model as *descendants* of a Clazz node, indicating that these attributes are properties of that class. This can be expressed in logic as:

 $\forall m : CodeModel, \forall c : Class \bullet (c \in_m m), \exists att : CAttribute \bullet (att \in_p c)$

6.4.1.3 Method

A *CMethod* element is utilised to define the structure of every method in a class. These act as wrappers for binding internal database operations, in which they declare syntax for calling a database stored procedure (pre-defined query) and/or for managing the database connectivity, such as $JDBC$ application in Java.

CMethod has two important kinds of child nodes, namely, an CArgument element to represent its input parameter and a Call element to express a procedure invoked:

 $\forall c : Class \forall f : method \bullet (f \in_{p} c), \exists a : CArgument \land Child(f, a)$

The *Call* element is used to declare a method invocation expression, which is typically a *child* element of Method. It refers to a syntax (expression), which is a part of the body of a method, that invokes a stored procedure located in a database schema, expressed previously in DBQ. The name of the invoked pre-defined query is declared in the *Call* node via its *name* attribute. The notion of *Call* can be defined using the following policy:

Law 64. Any method, in a class, might call a stored procedure, in a database schema.

 $\forall s : DatabaseModel, \ \forall m : CodeModel, \ \forall f : Method,$ $\forall c : Class \bullet ((c \in_m m) \land (f \in_p c)),$ $\exists p : procedure \bullet ((p \in_m s) \land Invokes(f, p))$

In the same context, a *Constructor* element, which is used to declare a construction (initialising) of a class, has a number of Variable elements and some assignment Operators, as its descendants, to represent its input parameters and initialise the attributes of the class.

6.5 Outlook on the Chapter

This chapter represented, in-depth, the concepts, underlying ASTs (XML), and formal semantics for each Micro Modelling Language (μML) model appearing in the *Design* phase of BUILD. It covered all critical concepts, their usage, and all their possible interpretations within the various contexts of the μ ML Database and Query (DBQ), Graphical User Interface (GUI) and Code Model.

The Code Model was introduced to be used in later versions of BUILD for developing 3-tier ISs, whereas models may be used alone for generating complete 2-tier business applications. Business logics were expressed using predefined stored procedures as a part of database schema. Based on Relational and Functional Algebra, a query modelling language was introduced for expressing a number of database operations in the DBQ model.

Apart from this, the chapter showed how to reuse the graphical notation from the Data Dependency Model for expressing DBQ concepts. There is no graphical representation in the current version of μ ML for the GUI and Code model. The formal semantics and notation for models and their elements are compared to similar or closely related UML concepts and notation, showing the simplicity and clarity of the proposed language. A number of unary and n-ary predicates were introduced for defining concepts and their relationships in each model in respect of the previously defined metamodel.

7 The Architecture of the Model-Transformation Approach

"If something is worth doing once, it's worth building a tool to do it " Unknown

7.1 Context

This chapter is divided into two main sections. The first one provides an in depth description of the overall structure and mechanism of the model transformation approach, highlighted briefly in Chapter 3. The approach encompasses a two level transformation framework, namely, top level and concrete level.

The first sections below discuss the generic design of the top level framework. This level is an ancestor of all actual transformation rules used to generate several μ ML models via the different BUILD development stages. The top level is adopted since it is the predecessor of all actual transformation rules used to generate a variety of μ ML model(s) via the different BUILD development stages, and the architecture of the concrete (actual) model transformation framework.

In addition, the second part of the chapter clarifies, in depth, the overall structure and mechanism of the code generation approach, mentioned briefly in Chapter 3. The current version consists of two platform specific generators. Each generator has two levels of generator agents. Throughout later sections, the overall and the detailed design of the code generation framework is discussed.

7.2 Brief Overview of the Model Transformation Framework

The transformation architecture developed for BUILD is simple and flexible. Every model manipulation is classified either as a kind of translation, transformation, folding (merging) and in-place modification. These four types of model transformation rules are supported by the current version of the framework. A translation rule translates a concept from a source model into an equivalent concept in a different type of target model (exogenous). A transformation rule transforms between models expressed in the same language (endogenous). In the case of optimising a model (where only one model is involved in a rule as both its source and target), this rule is considered *in-place modification*. In contrast, when concepts from more than one model are involved in a rule as source, the rule is regarded as afolding (merging) transformation rule.

It is worth mentioning that the proposed architecture of transformations in BUILD differs from the previous ReMoDeL [101] version that was used for the MySQL Database Generator, presented in [109] and [108]. That code generation framework encodes transformation rules as methods in each generator responsible for generating a part of the target (imperative). On the other hand, the current version of BUILD encodes each rule of transformations as an independent java class, representing the declarative side of the language. Chapter 8 explained how this leads to a hybrid declarative/imperative approach, where individual transformations are both independent and idempotent, while some order of execution is eventually imposed.

The overall architecture of BUILD is divided into two main layers: abstract (generic) and specific. Two types of translation rules are used in our transformation approach, namely, (oneto-one translating) Simple and (two-to-one merging) Merging rules. Our Simple rules take an element from a model and generate a new element in another model. Translation, transformation and *in-place modification* rules are implemented in the framework as *Simple* rules. On the other hand, our *Merging* rules, in the current version of BUILD, take two elements from different models and generate a new element in a target model.

A simple forward transformation (refinement or code generation) occurs when a source model belongs to a different metamodel type than the target. This activity is frequently seen in the majority of transformation components in BUILD. It enables the shift from a higher development phase to a lower one, introducing new and richer knowledge. The following section discusses these kinds of transformations in more detail.

7.3 The Top Level Architecture of the Model Transformation Framework

The top level architecture of the model transformation framework is an Object-Oriented (OO) approach that consists of five main classes, namely Rule, MergeRule, Translation, MergeTranslation and Context, illustrated in Figure 7.1. The root classes of the framework are the Rule and MergeRule class.

Figure 7.1: Model Transformation in BUILD. The Top Level Framework

The Rule class is considered the ancestor of all rules applied to a single source that is regarded a super class of a single subclass (so far), called Translation. Based on the type of source and target, this subclass acts as a *model translator* that translates a source model to a target one with a different metamodel, a *model transformer* that translates a source model to a target one within the same metamodel, and an *in-place modifier* that modifies and evolves a source model.

Figure 7.2: Translation Rule Structure

The above Figure (7.2) exemplifies the Translation rule structure. It includes a simple rule (class) from the concrete level, called $DAttributeToColumn$ class. that is used to translate every Data Dependency attribute (*DAttribute*) into a DBQ column (*Column*).

In addition to this, *MergeRule* is considered a general class of a subclass (so far), called MergeTranslation. It takes two source elements and translates (merges) them to produce a target. The following Figure (7.3) exemplifies the *MergeTranslation* rule structure. It includes a simple rule (class) from the concrete level, called *DDiagramToSchema* class. that takes DataFlow and Data Model element to translate them into a DBQ Schema.

Both Rule and MergeRule have an appropriate accept method when the rule has particular restrictions on the valid source elements required to be examined and accepted. This method is effectively the precondition of the rule, which returns true when the source element is valid for

Figure 7.3: Translation Rule Structure

the rule. An overriden version of *accept* may appear in any subclass of Rule and MergeRule if it has particular restrictions on the valid source elements in the subclass.

The Context class is the fifth class in the top level framework. It is a record of the work done by a tree of Translation rules. A new Context is created whenever a Translation rule is created at the top level, to perform a translation. The same Context is shared by every descendant rule created by the top level Translation rule. The Context indexes all transformations performed by the tree of Transformation rules, under the name of each rule. This allows individual rules to access any existing translation and avoids repeated work.

7.4 Framework Architecture at The Concrete Level

The architecture of all translator components at the concrete level is directly influenced by the hierarchical structure of elements in the source model of each. A naming convention for each agent is specified using the following format: $\langle sourceelementname \rangle$ To $\langle targetelementname \rangle$.

For instance, considering the structure of nodes in the *Information Model*, the step of translating the Information Model into the Data Model (section 7.5.2) consists of Diagrams, containing Nodes (e.g. *DEntity*), which in turn contain *Properties* (e.g. *DAttribute* and *Identifier*). Additionally, it contains Relationship (e.g. Association, Generalisation and Composition), which in turn contains Properties (Roles).

Thus, the internal structure of each translator is based upon the concept distribution strategy. This means that each concept in a source model is handled separately by a particular agent (sub-translator), represented via a java class in the framework. Each agent is responsible for applying appropriate transformation rules to create the corresponding target concepts.

Agents present in a hierarchical series of layers delegate to other agents at the next finer level of abstraction. Sometimes these sub-translators refer back to information previously processed by the higher-level (parent) translator. Each translator, which might consist of a set of either Translation or MergeTranslation or even both rules, behaves in the same manner as the Composite and Visitor design pattern[42] as some agents, which are containers of others, have

a command *translate()* to delegate to their parts (sub-translators), and each agent traverses a part of the model. It is worth mentioning that the $translate()$ method is the top-level method. It may either look up a cached result or invoke doTranslate() to compute the result for the first time.

In addition to this, all mapping algorithms are encoded as methods of the framework. The algorithms are imperative, using an ordered collection of transformations on Task and Impact Models, represented as (as abstract syntax trees). The order of rules execution is controlled and implemented through the body of the $doTranslate()$ method in each agent.

7.5 Requirement-to-Analysis Model Transformation Approach

As previously presented in Chapter 3, the Requirement to Analysis Model Transformation step is regarded as a first forward mapping step that shifts the end-users requirement models towards the analysis phase. In order to construct the required models, which are conceptually located within the BUILD Analysis phase, a number of Java-translator components are designed to implement the (hybrid) mapping rules for each translator agent. These translators are: (Information-to-DataDependency, Task-Impact-to-DataFlow, DataFlow-to-DetailedDataFlow, DetailedDataFlow-to-State). Initially, the XML files representing the business user's input are parsed, creating source models. Then these rules are applied to the source models to generate in-memory target models, which in turn will be used for the source models in the next stage.

7.5.1 Translating Task and Impact Models into (initial) DataFlow

Translating Task and Impact Models into (initial) DataFlow is considered to be a forward merge translation step. This takes two elements as its *source* from a Task and Impact Model, producing an equivalent target element in the *DataFlow* artefact. This translation step consists of 15 translating agents that are responsible for traversing both Task and Impact Model, applying a set of transformation rules to perform the mapping between their elements and the target elements in the DataFlow Model. The following figure (7.4) demonstrates the internal structure of the translator.

From this, some agents are considered to be subclasses of the *MergeTranslation* class, at the top level of the framework. These translators are:

- DiagramToDfDiagram: The agent takes a Task Model Diagram and an Impact Model *Diagram* as source and produces an equivalent *DataFlow Model DfDiagram*. This rule consists of one subrule (*BoundaryToDfBoundary*).
- BoundaryToDfBoundary: The agent takes an (equivalent) Task Model Boundary and an Impact Model Boundary as source and produces an equivalent DataFlow Model Df-Boundary. This significant translator also exemplifies the merging process of two source elements to obtain a single target. It consists of nine subrules that are responsible for mapping other concepts from the source to corresponding ones in the target DataFlow Model.
- MergeTaskToDfTask: The agent takes a Task Model Task and an Impact Model Task as source and produces an equivalent DataFlow Model Task. This agent maps two source elements, extracted from different models, and produces a single target element.

Figure 7.4: Req-to-Analysis: Task & Impact to (initial) DataFlow Model

The rest of the agents are subclasses of the *Translation* class (*Simple* rule), at the top level framework. These translators are:

- TaskToBoundary: The agent takes a Task Model composite Task as source and translates it into an equivalent DataFlow DfBoundary. This agent is interesting because it is controlled by a precondition that must be satisfied (a task is composite) to be translated into a dataflow boundary.
- ParticipationToInputFlow: The agent takes a Task Model Participation as source and translates it into an equivalent *DataFlow InputFlow*. The type of user interaction must be an input to satisfies the precondition of this rule. ParticipationToInputFlow consists of one subrule to translate endroles of each Participation, RoleToDfRole rule.
- ParticipationToOutputFlow: The agent takes a Task Model Participation as source and translates it into an equivalent DataFlow OutputFlow. The type of user interaction must be an *output* to satisfies the precondition of this rule. ParticipationToOutFlow consists of one subrule to translate the Participation endroles, namely, RoleToDfRole rule.
- RoleToDfRole: The agent takes a Task Model endrole Role of every Participation as source and translates it into an equivalent *DataFlow DfRole* of either *InputFlow* or *Out*putFlow. Any source task from the Task Model must be simple (not composite) in order to satisfy the precondition of this rule. RoleToDfRole consists of two subrules: TaskToDf-Task and ActorToDfActor.
- TaskToDfTask: The agent takes the Task referenced by an end $Role$ either in the Task Model or Impact Model and translates it into equivalent DataFlow Model DfTask.
- ActorToDfActor: The agent takes a Task Model Actor as source and translates it into an equivalent DataFlow DfActor.
- CreateFlowToDfCreateFlow: The agent takes a *Impact Model CreateFlow* as source and translates it into an equivalent *DataFlow CreateFlow*. This rule consists of one subrule to translate the endroles, $RoleToDfRole$ rule.
- DeleteFlowToDfDeleteFlow: The agent takes a *Impact Model DeleteFlow* as source and translates it into an equivalent *DataFlow DeleteFlow*. This rule consists of one subrule to translate the endroles, $RoleToDfRole$ rule.
- UpdateFlowToDfUpdateFlow: The agent takes a Impact Model UpdateFlow as source and translates it into an equivalent *DataFlow UpdateFlow*. This rule consists of one subrule to translate the endroles, $RoleToDfRole$ rule.
- ReadFlowToDfReadFlow: The agent takes a *Impact Model ReadFlow* as source and translates it into an equivalent DataFlow ReadFlow. This rule consists of one subrule to translate the endroles, RoleToDfRole rule.
- ImpObjectToDfEntity: The agent takes a refered *Impact Model ImpObject* by either any Role in the Impact Model as source and translates it into an equivalent DataFlow DfEntity.
- ImpRoleToDfRole: The agent takes a Task Model endrole Role of every Flow as source and translates it into an equivalent DataFlow DfRole of the equivalent DataFlow Flow.

7.5.2 Translating the Information Model into the Data (Dependency) Model

Translating an Information Model into a Data (Dependency) Model is regarded as a one-toone forward translation step. It takes one element as its source from an Information Model, producing an equivalent target element in the Data Model. This mapping step consists of 10 translating agents that are responsible for traversing nodes in the Information Model, applying a set of rules to perform the transformation. The following figure (7.5) shows the internal structure of the translator.

For this step, all agents are considered subclasses of the Translation class, at the top level framework. These translators are:

• InfDiagramToDDiagram: The agent takes an *Information Model Diagram* as source and produces an equivalent *Data Model DDiagram*. It is considered a root rule of this transformation step that consists of seven subrules.

Figure 7.5: Req-to-Analysis: Information to Data Dependency Model

- InfEntityToDEntity: This agent provides a direct mapping as it takes an Information Model InfEntity as source and produces an equivalent Data Model DEntity. This agent is interesting because it can produce knew knowledge in the target model (e.g. manufactured identity), if needed.
- InfAttributeToDAttribute: The agent takes an *Information Model Attribute* as source and produces an equivalent Data Model DAttribute. Data types and additional specifications (e.g. size and null) might be generated precisely, if needed.
- InfAssocToDDependency: The agent takes an Information Model Association as source and produces an equivalent *Data Model Dependency*. This rule is restricted by a precondition, in which the source association must be many-to-one. It consists on only one subrule for translating endroles $(InfRoleToDRole)$ of that association.
- InfAssocToMergedDEntity: The agent takes an *Information Model Association* as source and produces an equivalent *Data Model DEntity*. This rule is applicable for translating many-to-many associations only (precondition). It consists of two subrules to proceed the translation of attributes and identifiers for the generated target element. These subrules are: InfAttributeToDAttribute and IdentifierToUnique.
- InfCompToDDependency: The agent takes an *Information Model Composition* as source and produces an equivalent *Data Model Dependency*. A precondition is used to control this rule to be applied to not *Total Composition*. The *InfCompToDDependency* rule consists of one subrule for translating endroles (InfRoleToDRole) of that composition.
- InfGenToDDependency: The agent takes an *Information Model Generalisation* (over*lapping*) as source and produces an equivalent *Data Model Dependency*. This rule consists of one subrule for translating endroles (InfRoleToDRole) of that generalisation.
- InfGenToMergedDEntity: The agent takes an Information Model Generalisation as source and produces an equivalent *Data Model DEntity*. This rule is restricted to *disjoint* Generalisation only, via a precondition. It consists of two subrules to complete the mapping of attributes and identifiers for the generated entities. These subrules are: InfAttributeToDAttribute and IdentifierToUnique.
- Identifier To Unique: The agent takes an Information Model Identifier as source and produces an equivalent Data Model Unique concept.
- InfRoleToDRole: The agent takes an *Information Model Role* as source and produces an equivalent Data Model DRole.

7.5.3 Translating (initial) DataFlow Model into Detailed DataFlow Model

Translating the (initial) DataFlow Model into the detailed one can be described as a one-to-one forward translation step. In particular, it is regarded as a model optimisation and evolving step that aims to decompose each business task that appears in the initial DataFlow into its collection of atomic sub-tasks after applying a number of decompositional rules to each *DfTask*. It is worth mentioning that the rule of transformation is presented in more detail further in Chapter 8.

This mapping step consists of 8 translating agents that are responsible for traversing the nodes in the initial *DataFlow Model*, applying a set of rules to perform the transformation. The following figure (7.6) shows the internal structure of the translator. For this step, all agents are considered to be subclasses of the *Translation* class, at the top level of the framework. These translators are:

- ArcToArc: The agent takes one of the (initial) DataFlow Flow elements, where its type is either *create, read, delete, input* or *output*, as source and produces an equivalent (detailed) DataFlow Model Flow.
- ArcToArcs: The agent takes an (initial) DataFlow Flow element, where its type is update, as source and produces an equivalent pair of read and write flows in the (detailed) DataFlow Model. This rule is applicable for *update* flows only.
- ArcToDfActor: The agent takes one of the (initial) DataFlow Flow elements, where its type is either *input* or *output*, as source and produces an equivalent (detailed) DataFlow Model DfActor.
- ArcToDfObject: The agent takes one of the (initial) DataFlow Flow elements, where its type is either create, read, update or delete, as source and produces an equivalent (detailed) DataFlow Model DfObject.
- ArcToDfTask: The agent takes one of the (initial) DataFlow Flow elements, where its type is either *create, read, delete, input* or *output*, as source and produces an equivalent (detailed) DataFlow Model DfTask.

Figure 7.6: Req-to-Analysis: (initial) DataFlow to Detailed DataFlow Model

- ArcToDfTasks: The agent takes an (initial) DataFlow Flow elements, where its type is update, as source and produces equivalent (detailed) DataFlow Model DfTask elements.
- DfdDiagramToDfdDiagram: The agent takes a DataFlow Model DfdDiagram as source and produces an equivalent DataFlow Model DfdDiagram.
- **DfTaskToDfBoundary:** The agent takes a DataFlow Model *DfTask* as source and produces an equivalent DataFlow Model DfBoundary.

7.5.4 Translating DataFlow Model into (Screen) State Model

Translating a (detailed) DataFlow Model into a State Model can be defined as a forward translation step that takes an element as a source from a *DataFlow*, constructing an equivalent target element in the Screen State artefact. This step consists of 6 translating agents that traverse the tree structure of the DataFlow Model, applying a set of transformation rules to map a source element to a target one in the State Model. The following figure (7.7) illustrates the internal structure of the translator.

For this step, all translating agents are considered subclasses of the Translation class, at the top level framework. These translators are:

- DfdDiagramToStDiagram: The agent takes a DataFlow Model DfdDiagram as source and produces an equivalent *State Model StDiagram*. This root rule consists of two main subrules: DfdBoundaryToStBoundary and StBoundaryToStBoundary.
- DfdBoundaryToStBoundary: The agent takes a DataFlow Model *DfBoundary* as

Figure 7.7: Req-to-Analysis: DFD to State Model

source and produces an equivalent *State Model StBoundary*. This rule consists of three subrules for translating the content the source model into equivalent *States* and *Transi*tions (ArcToStReadyState, ArcToStWaitingState and DfTaskToTransition).

- StBoundaryToStBoundary: This is an *in-place* model modification step. The agent takes a State Model StBoundary as source and produces generate a number of Error States, which capture the failure scenarios of the business process, to be added to the current boundary StBoundary.
- ArcToStReadyState: The agent takes a *DataFlow Model Flow* as source and produces an equivalent State Model ReadyState. Each acceptable arc must be one of the CRUD flows only.
- ArcToStWaitingState: The agent takes a *DataFlow Model Flow* as source and produces an equivalent State Model WaitingState. Each acceptable arc must be either *input* or output flow. Otherwise, this rule is dismissed.
- **DfTaskToTransition:** The agent takes a *DataFlow Model DfTask* as source and produces an equivalent State Model Transition.
- TaskToStWaitingState: This rule takes every *composite Task* in the Task Model to produce a waiting state.

7.6 Analysis-to-Design Model Transformation Approach

According to the overall structure of the BUILD framework, presented in chapter 3, the Analysis to Design Model Transformation step is considered to be a second forward mapping step that translates some intermediate analysis models, resulting from the previous Requirement to Analysis Model Transformation process, into lower level design artefacts that describe the system at the design phase.

Similar to the previously discussed Requirement to Analysis Model Transformation step, a number of Java-translator programs are designed to implement the (hybrid) transformation rules for each agent. These translators are: State-to-GUI and Data-DFD-to-DataBaseQuery.

They parse the *XML* representation of the μ ML analysis models, namely, DataFlow, Data Dependency, and State Model and produce a number of design models: GUI, DBQ and (optional) Code artefacts ready for the code generation phase.

7.6.1 Translating the State Model into the GUI Model

Translating State Model into a GUI Model is considered a direct forword translation step. It takes an element from a State Model, as its source, producing an equivalent target element in the GUI Specification Model. The internal structure of the translator is simple, in which the current version consists of 6 translating agents that are responsible for traversing *State Model*, applying transformation rules to perform the mapping to target elements in the GUI Model. The following figure (7.8) demonstrates the internal structure of the translator.

Figure 7.8: Analysis-to-Design: State to GUI Model

From this, all translating agents are considered subclasses of the Translation class, at the top level framework. These translators are:

- StDiagramToGuiDiagram: The agent takes a *State Model StDiagram* as source and produces an equivalent GUI Model GuiDiagram. This rule consists of one subrule for translating boundaries.
- StBoundaryToGuiBoundary: The agent takes a *State Model StBoundary* as source and produces an equivalent GUI Model GuiBoundary.
- StateToWindow: The agent takes a *State Model State* as source and produces an equivalent GUI Model Window. This rule consists of two subrules for translating widget controls for that window, namely, SvariableToLabel and SvariableToTextfield.
- SvariableToLabel: The agent takes a *State Model Variable* as source and produces an equivalent GUI Model Label.
- SvariableToTextfield: The agent takes a *State Model Variable* as source and produces an equivalent GUI Model Textfield.
- **TransitionToButton:** The agent takes a *State Model Transition* as source and produces an equivalent GUI Model Button.

7.6.2 Translating the DataFlow and Data Model into the DBQ Model

Described previously in Chapter 6, the DBQ model consists of two kinds of concepts: data definition and *query expression.* "This translation step has two source models, the DataFlow model and the Data Dependency model.

All agents that aim to produce data definition concepts in the target model are actually designed and implemented as one-to-one forward translation components. Each of these takes one element as its *source* from either a *DataFlow* or *Data Dependency Model*, producing an equivalent target element in the Database Model (DBQ).

On the other hand, agents that aim to produce query expression concepts in the target model are designed as *two-to-one* merging and *one-to-one* translation components. Each of these takes elements from one or both the *DataFlow* and *Data Model* as source and translates them into a target DBQ concept.

This mapping step consists of 12 translating agents, in total, that are responsible for traversing nodes in either DataFlow or Data Dependency or even both models, and applying a set of rules to perform the transformation. The following figure (7.9) shows the internal structure of the translator.

From this, some agents are considered to be subclasses of the *MergeTranslation* class, at the top level framework. These translators are:

- CreateFlowToStoredProcedure: The agent takes a *DataFlow Model CreateFlow* and an Data Model DDiagram as source and produces an equivalent DBQ Model Procedure. The argument of the target procedure is translated by *CulomnToArgument* subrule.
- ReadFlowToStoredProcedure: The agent takes a DataFlow Model ReadFlow and an Data Model DDiagram as source and produces an equivalent DBQ Model Procedure. This rule consists of one subrule for translating required arguments of that procedure.
- DeleteFlowToStoredProcedure: The agent takes a *DataFlow Model DeleteFlow* and an Data Model DDiagram as source and produces an equivalent DBQ Model Procedure. The argument of the target procedure is translated by *CulomnToArgument* subrule.
- WriteFlowToStoredProcedure: The agent takes a *DataFlow Model WriteFlow* and an Data Model DDiagram as source and produces an equivalent DBQ Model Procedure. This rule consists of two subrules for translating required arguments and local variables of that procedure.

The rest of agents are subclasses of the one-to-one *Translation* class, at the top level framework. These translators are:

• DDiagramToSchema: The agent takes a *Data Model DDiagram* as source and translates it into a DBQ equivalent Schema. This rule consists of seven subrules to proceed

Figure 7.9: Analysis-to-Design: Data Dependency to Database and Query Model

with this translation step.

- DEntityToTable: The agent takes a *Data Model DEntity* as source and translates it into a DBQ equivalent Table. This rule consists of two subrules for translating attributes and primary keys of that generated table.
- DEntityToPrimaryKey: The agent takes a *Data Model DEntity* as source and translates it into a DBQ equivalent *PrimaryKey*.
- DAttributeToColumn: The agent takes a *Data Model DAtribute* as source and translates it into a DBQ equivalent Column.
- DependToReferTable: The agent takes a $Data Model Department$ as source and translates it into a DBQ equivalent Table. This rule consists of one subrule for translating a foreignkey that referred to the generated table.
- DRoleToForeignKey: The agent takes a *Data Model DRole* as source and translates it into a DBQ equivalent DfEntity.
- ForeignKeyToStoredProcedure: The agent takes a *DBQ Model ForeignKey* as source and translates it into a DBQ equivalent Procedure. This rule consists of two subrules for translating required arguments and local variables of that procedure.
- OopToDbType: The agent takes an OOP data type and translates it into an equivelant target database datatype.

7.7 Alternative Model Transformation Steps

In this section, the structure of the optional translation steps are discussed. The current version of BUILD has two additional translation steps that are not a part of the information systems development process, as currently implemented. These steps are: translating the *Impact Model* into the (initial) Information Model and translating DataFlow and DBQ Models into the Code Model.

7.7.1 Translating the Impact Model into the (Initial) Information Model

Translating an Impact Model into an (initial) Information Model in is an alternative one-to-one translation step toward generating a complete data model of a system. It takes object and impact elements from an *Impact Model*, as its *source*, generating an equivalent *Information Model Entity* and Associations as target elements. This translation step consists of 6 Java agents, in total, that are responsible for traversing nodes in an Impact Model, and applying a set of rules to perform the dirct translation between elements and predict the association between them. The following figure (7.10) shows the internal structure of this translator.

Figure 7.10: Alternative Transformation: Impact to Information Model

From this, all translating components are regarded subclasses of the Translation class, at the top level framework. These translators are:

- ImpBoundaryToDiagram: The agent takes a *Impact Model ImpDiagram* as source and produces an equivalent Information Model Diagram.
- ImpBoundaryToEntity: The agent takes a boundary $(ImpBoundary)$ from the *Impact* Model as source; each enitity, within that boundary, is mapped into an equivalent enitity (*Entity*) in the *Information Model*. This rule consists of one subrule (*ImpObjectToEntity*) that continues the translation.
- ImpBoundaryToAssociation: The agent takes a Impact Model ImpBoundary as source; each flow, in that boundary, is translated into an equivalent *Information Model Associa*tion. This rule consists of a single subrule $(ImpRoleToRole)$ that is responsible for map-

ping the types of endroles of each flow into the quivalent ones in the target Information Model.

- ImpConjunctionToAssociation: The agent takes a *Impact Model ImpBoundary* that has one or more conjunction *ImpConjunction* as source; for each conjunction within that boundary, the rule generates equivalent Information Model Associations. This rule also consists of one subrule $(ImpRoleToRole)$ that is responsible for mapping endroles of each conjunction into the quivalent ones in the Information Model.
- ImpObjectToEntity: The agent takes a *Impact Model ImpObject* as source and produces an equivalent Information Model Entity.
- ImpRoleToRole: The agent takes a *Impact Model ImpRole* as source and produces an equivalent Information Model Role.

It is worth saying that this translation agent is implemented and used later in Chapter 9 for generating the Information Model in a case study. The rule of transformation is presented later with further detail in Chapter 8.

7.7.2 Translating the DataFlow and DBQ Model into the Code Model

Translating a Database and Query and DataFlow Model into an (OO) Code Model is considered to be a forward merging step. It takes two elements as its *source* from a DBQ and DataFlow Model, producing an equivalent target element in the Code model. Finalising the clean version of this translation agent, which supports all possible OO code features, is not implemented completely at the date of writing this thesis.

The agent consists of 15 translating agents that are responsible for traversing both *Database* and Query and DataFlow Models, applying a set of transformation and merging rules to perform the mapping between their elements and the target elements in the *Code Model*. For this step, some agents are considered to be subclasses of the *MergeTranslation* class, at the top level framework. These translators are:

• DDiagramToCDiagram: The agent takes a *DBQ Schema* and DFD Diagram as source and produces an equivalent *Code Model Diagram* element. This rule consists of three main subrules: TableToClass, DfBoundaryToClass and StoredProcedureToMethod for generating entity classes, process classes and methods respectively.

The rest of the agents are subclasses of the one-to-one *Translation* class, at the top level framework. These translators are:

- DfBoundaryToClass: The agent takes a *DataFlow DfBoundary* as source and translates it into a Code Model equivalent Clazz. The generated class represents a process class (business task class). This rules consists of two main subrules: DfObjectToAttribute and DfActorToAttribute for generating all required attributes (fields) of that process class.
- ColumnToCArgument: The agent takes a *DBQ Model Column* as source and translates it into a Code Model equivalent Argument element.
- ColumnToVariable: The agent takes a *DBQ Model Column* as source and translates it into a Code Model equivalent Variable.
- DbToOopType: The agent takes a *DBQ Model* datatype of every element as source and translates it into an acceptable OO datatype that is equivalent to the source one.
- **OopToDbType:** The agent takes a *Code Model* datatype of every element as source and translates it into an acceptable relational database datatype.
- DfActorToAttribute: The agent takes a $DataFlow$ DfActor as source and translates it into a Code Model equivalent Attribute. The generated attributes are a member of a class that is represent a business task (not a business entity).
- DfObjectToAttribute: The agent takes a $DataFlow\ DfObject$ as source and translates it into a Code Model equivalent Attribute. The generated attributes are a member of a class that is represent a business task (not a business entity).
- AttributeToIdentifier: The agent takes a recently generated Code Model Attribute as source and translates it into a Code Model equivalent Identifier, which might be used in such a method or a constructor.
- StoredProcedureToMethod: The agent takes a DBQ Model Procedure as source and translates it into an equivalent Code Mode's Method. This rule consists of two subrules: proArgumentToMethArgument and proArgumentToMethResult for translating its input arguments and the type of its return result.
- ProArgumentToMethArgument: The agent takes an argument of DBQ Model Procedure as source and translates it into a Code Mode's Method equivalent Argument.
- ProArgumentToMethResult: The agent takes an result (return value) of DBQ Model Procedure as source and translates it into a *Code Mode's Method* equivalent return value.
- TableToArgument: The agent takes a *DBQ Model Table* as source and translates it into an equivalent Code Model's Argument. This generated argument is used to pass object to method and/or constructor.
- TableToClass: The agent takes a DBQ Table as source and translates it into a Code Model equivalent Clazz. The generated class represents a business entity class. It is consists of two subrules: ColumnToAttribute and TableToConstructor for translating attributes and constructor for that entity class.
- TableToConstructor: The agent takes a *DBQ Table* as source and translates it into a Code Model equivalent Constructor for each generated business entity class.
- VariableToIdentifier: The agent takes a recently generated Code Model Variable as source and translates it into a *Code Model* equivalent *Identifier*, which might be used in such a method or a constructor.

7.8 The Implementation of μ ML Models

Java is adopted both for constructing the model and for the transformation technology the model and the transformation technology in BUILD. It is a widely known OO programming language and does not force any further conceptual load on developers, unlike rule-based transformation languages that require a a steep learning curve before introducing a new translation agent or a new system aspect into the framework. Models in BUILD are abstract syntax trees (ASTs), built in the same programming language (Java).

Two types of *Java Packages* are used to implement μ ML models, namely, a core package and several *concrete packages*. The current version of BUILD has one *core package* that contains several *Java* classes for representing core implementation of the μ ML metamodel elements (see Table 7.1). On the other hand, the framework contains seven *concrete packages* for representing language concepts of each μ ML model (information system view), each package has a number of classes for defining these concrete concepts (Table 7.2).

At the core metamodel level, a concept might inherit from another one. The following listing (7.1) illustrates the representation of Node.java class. It interits from the core Element class, and has constructors, attributes and a number of get/set methods. The remaining of concepts are defined using the same strategy.

Listing 7.1: Construction of Node Concept

```
1 public class Node extends Element {
\frac{2}{3}protected String name:
4
5 public Node() {
6 }
7
8 public Node (String name) {
9 this. name = name;
10 }
11
12 public String getName() {
13 return name;
14 }
15
16 public Node setName(String name) {
17 this. name = name;
18 return this;
19 }
20 }
```
As it is mentioned before, a number of *concrete packages* are used to represent the concept of each μ ML model. The following table (Table 7.2) demonstrates μ ML concepts for each *concrete* package that are defined using Java classes.

Table 7.2: Java Package of the main μ ML Concrete Elements

Elements of memory models are indexed by their names. A Java LinkedHashMap object is used to store memory elements in the same order as they appear in the model, because some transformation rules need to check the order of occurrence of some elements to make a mapping decision. The following listing (7.2) demonstrates the implementation of the Information Model Entity.java concrete class:

Listing 7.2: Construction of the Entity element in the Information Model

```
1 public class Entity extends Type {
2
3 private Map<String, Attribute> attributes;
4
5 public Entity() {
6 attributes = new LinkedHashMap<String, Attribute>();
7 }
8
9 public Entity(String name) {
10 super(name);
11 attributes = new LinkedHashMap<String, Attribute>();
12 }
13
14 public List<Attribute> getAttributes() {
15 return new ArrayList<Attribute>(attributes.values());
16 }
17
18 public Entity addAttribute(Attribute datt) {
19 attributes.put(datt.getName(), datt);
20 return this;
21 }
22
23 public Attribute getAttribute (String name) {
24 return attributes.get(name);
25 }
26 }
```
7.9 The Implementation of the Transformation Rules

In BUILD, the transformation rules are classified into two types, *top level* and *concrete* rule. The top level rules are implemented as abstract classes in the top level framework. The current version of BUILD has two abstract classes, namely, Rule and MergeRule. At the concrete level, each rule is implemented as a concrete java class that inherits either from Translation or MergeTranslation abstract classes at the top level framework. The following sections present examples of an abstract and concrete transformation rule.

7.9.1 Example of Top-Level Rule Implementation

The abstract Translation class is the ancestor of all translation rules, implemented in the concrete level, mapping from one *Source* element to one *Target* element. Every concrete forward translation rule must inherit from *Translation*, which is responsible for maintaining the shared Context of completed work. Figure 7.3 demonstrates the content of the Translation.java abstract class that expresses the generic translation rule at the top-level framework.

As mentioned earlier, the method $translate()$ first checks whether a translation already exists for the given source and, if so, it returns the corresponding target for that source. Otherwise it invokes the **abstract** method $doTranslate()$ and stores the translated target element in the context.

The shared Context is used by the *hasTranslation ()* method to check whether a translation exists for a particular source element. When such a translation exists, the $getTranslation()$ method is used to retrieve the translated target element for a given source, otherwise the method returns the *null* value. In contrast, when such a translation does not exist, the method *putTrans*lation() stores a target element as the translation for a given source one and returns it.


```
1 public abstract class Translation<Source, Target> extends Rule<Source> {
\overline{2}3 private Context context;
4
5 public Translation(Context context) {
6 this.context = context;
7 }
8
9 public boolean hasTranslation (Source source) {
10 Map<Object, Object> map = context.lookup(this);
11 return map.containsKey(source);
12 }
13
14 protected Target getTranslation(Source source) {
15 Map<Object, Object> map = context.lookup(this);
16 return (Target) map.get(source);
17 }
18
19 public Target putTranslation(Source source, Target target) {
20 Map<Object, Object> map = context.lookup(this);
21 map.put(source, target);
22 return target:
23 }
24
25 protected abstract Target doTranslate(Source source);
```

```
26
27 public Target translate(Source source) {
28 if (hasTranslation(source))
29 return getTranslation(source);
30 else
31 return putTranslation(source, doTranslate(source));
32 }
33
34 }
```
7.9.2 Example of Concrete Rule Implementation

The transformation rule for translating Data Model's DEntity into DBQ's Table is considered in this section to exemplify its implementation and illustrate how the rule invokes another one in an imperative style. Figure 7.4 represents the content of the $DEntityToTable.java$ class that expresses the translation rule for generating DBQ table from a generated Data Model.

Listing 7.4: Construction of the Entity element in the Information Model

```
1 public class DEntityToTable extends Translation<DEntity, Table> {
2
3 private DAttributeToColumn dAttributeToColumn;
4 private DEntityToPrimaryKey dEntityToPrimKey;
5
6 public DEntityToTable() {
7 this(new Context());
8 }
9
10 public DEntityToTable(Context context) {
11 super(context);
12 dAttributeToColumn = new DAttributeToColumn (context) ;
13 dEntityToPrimKey = new DEntityToPrimaryKey(context);
14 }
15
16 public boolean accept (DEntity dentity) {
17 return true;
18 }
19
20 public Table doTranslate(DEntity dentity) {
21 Table table = new Table();
22 table.setName(dentity.getName());
23 if (dEntityToPrimKey.accept(dentity)) {
24 PrimaryKey primary = dEntityToPrimKey.translate(dentity);
25 if (primary.isSynthetic())
26 table.addColumn(primary.getColumn());
27 table.addPrimaryKey(primary);
28 }
29 for (DAttribute dattr : dentity.getDAttributes()) {
30 if (dAttributeToColumn.accept(dattr))
31 table.addColumn(dAttributeToColumn.translate(dattr));
32 }
33 return table;
34 }
35 }
```
The concrete *DEntityToTable* class inherits from the *Translation* class at the top level. The rule has no preconditions to be fired, as it seen in the body of $accept()$ method. Two subrules are declared declaratively, without taking into account their order of execution, namely, DAttributeToColumn and DEntityToPrimaryKey.

The body of the *DEntityToTable*, which represents exactly how the target element is developed, is expressed imperatively as a body of the concrete $doTranslate(...)$ method. The following algorithm demonstrates how to translate *Data Model DEntities* into *DBQ Tables*:

Result: Table initialise DAttributeToColumn object; initialise DEntityToPrimaryKey object; while DEntityToTable accepts a new source element do translate the source element (DEntity); if DEntityToPrimaryKey accept the source then generate a primary key from the source; add the generated key to the generated table; end forall the DAttribute in the current DEntity do if DAttributeToColumn accept the source then translate the current DAttribute from the source; add the generated Column to the generated table; end end go back to translate the next DEntity in the source model; end

Algorithm 1: Translating Data Model DEntities into DBQ Tables (Default)

In order to illustrate the degree of modularity of the designed agents (DAttributeToColumn and DEntityToPrimaryKey), they can be reused in the opposite order to produce the same output result. From that it can be concluded that these two agents are completely independent. The following algorithm shows an alternative order to invoking the subrules than the order shown in algorithm 1.

Result: Table initialise DAttributeToColumn object; initialise DEntityToPrimaryKey object; while *DEntityToTable accepts a new source element* do translate the source element (DEntity); forall the DAttribute in the current DEntity do if DAttributeToColumn accept the source then translate the current DAttribute from the source; add the generated Column to the generated table; end end if DEntityToPrimaryKey accept the source then generate a primary key from the source; add the generated key to the generated table; end go back to translate the next DEntity in the source model; end

Algorithm 2: Translating Data Model DEntities into DBQ Tables (Alternative)

7.10 Brief Overview of the Code Generation Framework

As previously presented in Chapter 3, the *Code Generation* step is considered the final forward mapping step, within BUILD framework, that generates executable code from a number of lowlevel platform-independent design models. The input models for this phase are: the GUI and DBQ Model, whereas, the output is a platform-specific executable OOP code and a relational SQL schema script. The current version of BUILD is able to generate a *Java Swings* application with a JDBC connection to a MySQL back-end database system is also generated from BUILD, producing a runnable 2-Tier Information System.

As a proof of concept, two domain-specific generators, namely, a *Java Swing and MySQL* generator, were designed. The overall architecture of the BUILD Code Generation Framework is described as a two levels of code generator agents. The top-level framework consists of a number of abstract classes, or Java agents, for holding common features of all sub-generators that are platform-specific for a target environment, such as MySQL.

7.11 Code Generation Approach (Design-to-Code)

In order to achieve a complete code generation facility, a number of Java-generator components are designed to implement the (imperative) mapping rules for each generator agent. These generators are: State-to-GUI and DataFlow-DataModel-to-Database. They accept as source the in-memory models produced by earlier transformation steps and produce executable code of each input artefact providing an executable code for a target system.

As the approach aims to generate code for different working environments, two layers of java classes are considered in the internal structure of each generator: abstract and concrete classes. The layer of abstract classes is constructed on the top of all groups of sub-generators, for a particular IS tier, to hold and share the common behaviours of the widely-known relational database vendors and OOP languages, which might be duplicated within different kinds of generators. Figure 7.11 demonstrates the sub-generators of the Java Swing GUI code generation framework.

For instance, when constructing the data tier, generating back-end databases in MySQL, Oracle, and other database system is possible when a relevant version of database generator exists. The following figure (Figure 7.12) illustrates sub-generators for generating a MySQL database system. These specific generators are implemented separately and grouped as concrete classes in Java.

Figure 7.11: The Architecture of the Java Swings UI Generator

For each concept in the target model, a specific agent (Java generator) is introduced. The transformation rules are implemented imperatively as methods within each agent. The following listing (7.5) shows an example of a method in MySQLFieldGenerator class for generating fields of a particular MySQL table.

7.12 Outlook on the Chapter

The chapter has discussed the overall architecture and mechanism of the transformation approach in BUILD. The transformation framework consists of two levels, an abstract top level and *concrete bottom level*. In the top level, two root rules (*Java classes*) are designed to implement the various types of mapping rules within the approach (e.g. one-to-one translation, two-to-one merging rule). Furthermore, for each compositional layer of transformation, the architecture of all translation agents at the concrete level is discussed in-depth in this chapter. Independent rules for each translation step are highlighted, showing their roles in the transformation approach.

Figure 7.12: The Architecture of the (MySQL) Database Generator

8 The Rules of Model Transformation

"Simplicity is the soul of efficiency" Austin Freeman

8.1 Context

This chapter presents and discusses the transformation rule for each translation step at the concrete level of the framewok. The sections are divided based on the transformations between the different development phases in BUILD, which mainly are Requirement to Analysis phase and *Analysis to Design phase*. For each target model, the rule that is used for deriving each target element is presented. The First-Order Predicate Logic (FOPL) statements are used to express and formalise the transformation rules.

8.2 Transformation Rules at the Requirement to Analysis Phase

8.2.1 Translating Task and Impact Models into (initial) DataFlow

While each external interaction between tasks and actors is expressed in the Task Model, and the internal interaction between the same tasks and system objects is described in the Impact Model, the DataFlow Model brings together both interactions in one diagram as a result of merging equivalent task from the Task and Impact Model.

8.2.1.1 DataFlow Diagram

Equivalent Diagram elements in the Task and Impact Model are mapped into a DataFlow one.

 $mr : Diagram \times ImpDiagram \rightarrow DfDiagram$ $\forall d1 : Diagram, \forall b1 : Boundary, m1 : TaskModel$ $\forall d2: ImpDiagram, \forall b2: ImpBoundary,$ $m2: ImpactModel \bullet ((b1 \in d1) \wedge$ $(d1 \in_m m1) \wedge (b2 \in d2) \wedge (d2 \in_m m2)$ \land NameOf(d1, d2) \land NameOf(b1, b2)) −→ ∃! d3 : DfDiagram, b3 : DfBoundary, $m3: DataFlowModel$ • $((d3 \in_m m3) \land (b3 \in d3))$ \wedge (NameOf(d1,d3) \wedge NameOf(d2,d3)) \wedge $(b3 = mr(b1, b2)))$

8.2.1.2 DataFlow Task

Each DfTask corresponds to the merge of both a Task Model Task and it is equivalent ImpTask in the Impact Model. This translation is performed by the MergeTaskToDfTask agent. Therefore, the rule can be stated as: for each two equivalent tasks in both Task and Impact Model, there exist a task (DfTask) in the generated DataFlow Model. The main translation rule for this translator can be expressed in logic as:

```
tr : Task, ImpTask \rightarrow DfDTask\forall t1 : Task, m1 : TaskModel,\forall t2: ImpTask, m2: ImpactModel\bullet ((t1 \in_m m1) \land (t2 \in_m m2) \land NameOf(t1, t2))−→
                                                            \exists !\; t3 : DfTask, m3 : DataFlowModel• ((t3 \epsilon_m m3) \wedge (NameOf(t1, t3)
                                                            \wedge NameO f(t2, t3)
```
8.2.1.3 DataFlow Boundary

Composite tasks in the Task Model are treated differently, in that they are translated into a DataFlow Boundary, containing its sub-tasks. This is handled by the TaskToBoundary agent. Thus, the transformation rule is: For each composite task in the Task Model, there exists a logical boundary in model of the DataFlow Model.

> $tr : Task \rightarrow DfBoundary$ $\forall t : Task, c : Composition, m1 : TaskModel$ • $((t, c \in_m m1) \wedge (Whole(c, t)))$ \longrightarrow

 $\exists ! b : DfBoundary, m2 : DataFlowModel$

• $((b \in_m m2) \land NameOf(t, b))$

In addition to this, a direct mapping rule exists: For each equivalent boundary in the Task Model and Impact Model, there exists a logical boundary in the DataFlow Model. This can be expressed in FOPL as:

 $mr : Boundary, ImpBoundary \rightarrow DfBoundary$

 $\forall b1 : Boundary, \forall p : Partitionation,$ $\forall t1 : Task, \ \forall a1 : Actor, m1 : TaskModel$ • $((t1, p, a1 \in_m b1) \land (b1 \in_m m1)),$ $\forall b2: ImpBoundary, \forall ar: Arc, \forall obj: ImpObject,$ $\forall t2: ImpTask, m2: ImpactModel$ \bullet ((t2, obj, ar \in_m b2) \land (b2 \in_m m2)) −→

> $\exists !\ b3 : DfBoundary, t3 : DfTask, a2 : DfAction,$ $f : Flow, e : DfEntity, m3 : DataFlowModel$ • $((t3, a2, f, e \in_m b3) \wedge (b3 \in_m m3))$ \wedge (NameOf(b1, b3) ∧ NameOf(b2, b3) ∧ $(t3 = mr(t1, t2)) \wedge (a2 = tr(a1)) \wedge$ $(f = tr(p)) \wedge (f = tr(ar))))$

8.2.1.4 DataFlow Entity

Each *DfEntity* is mapped directly to an *Impact Model* object (*ImpObject*). This translation is performed by the ImpObjectToDfEntity agent. The rule of this translation is: for each object in the Impact Model, there exist a DataFlow Entity in the generated DataFlow Model. It can be expressed in FOPL as:

$$
tr: ImpObject \rightarrow DfEntity
$$

\n
$$
\forall obj: ImpObject, m1: ImpactModel \bullet (obj \in_{m} m1)
$$

\n
$$
\Rightarrow
$$

\n
$$
\exists! ent: DfEntity, m2: DataFlowModel
$$

\n
$$
\bullet ((ent \in_{m} m2) \land NameOf(obj, ent))
$$

8.2.1.5 DataFlow Actor

Each DfActor is translated simply to a Task Model actor (Actor). This translation is performed by the ActorToDfActor agent. The rule of this translation is: for each actor in the Task Model, there exist a DataFlow Entity in the generated DataFlow Model. It can be expressed in FOPL as:

> $tr : Actor \rightarrow DfActor$ $\forall a1: Actor, m1: TaskModel \bullet (a1 \in_m m1)$ \longrightarrow $\exists ! a2 : DfActor, m2 : DataFlowModel$

• $((a2 \in_m m2) \land NameOf(a1, a2))$

8.2.1.6 DataFlow Flows

Each Task Model Participation is checked first to see whether or not it is an Input or an Output. This check is carried out by *BoundaryToDfBoundary* merge translation agent. Based on the kind of *Participation*, the decision to translate it into *DataFlow DfInputFlow* or *DfOutputFlow* is made. The translation is perform by either *ParticipationToDfInputFlow* and *Participation*-ToDfOutputFlow component. From that, two basic rules can be extracted:

1- For each Input participation in the Task Model, there exist an Input flow in the DataFlow Model.

 $(DfInputFlow \leq: Arc)$ $tr : Participation \rightarrow DfInputFlow$ $\forall p : Participation, (r1, r2) : Role, m1 : TaskModel$ $\bullet ((p, r1, r2 \in_m m1) \land (Input(p)))$ \longrightarrow $\exists ! f : DfInputFlow, (r3, r4) : DfRole,$ $m2 : DataFlowModel \bullet ((r3 = tr(r1))$ $∧ (r4 = tr(r2)) ∧ (f, r3, r4 ∈_m m2))$

2- For each Output participation in the Task Model, there exist an Output flow in the DataFlow Model.

(DfOutputF low <: Arc) tr : P articipation → DfOutputF low ∀ p : P articipation,(r1, r2) : Role, m1 : T askModel • ((p, r1, r2 ∈^m m1) ∧ (Output(p))) −→ ∃! f : DfOutputF low,(r3, r4) : DfRole, m2 : DataF lowModel • ((r3 = tr(r1)) ∧ (r4 = tr(r2)) ∧ (f, r3, r4 ∈^m m2)

Besides this, the kind of Impact in the Impact Model is also checked by the BoundaryToDf-Boundary merge translation agent that traverse the Impact Model XML tree handling each kind of Impact separately. Each Impact is mapped directly to an equivalent Flow in the DataFlow Model. The translation is perform by either *ParticipationToDfInputFlow* and *ParticipationToD*fOutputFlow component. From that, four basic rules can be extracted:

1- For each create impact in the Impact Model, there exist an create flow in the DataFlow Model.

 $(DfCreate Flow \le Arc)$ $tr: CreateFlow \rightarrow DfCreateFlow$ $\forall cf1: CreateFlow,(r1,r2): ImpRole,$ $m1: ImpactModel \bullet (cf1, r1, r2 \in_m m1)$ \longrightarrow $\exists ! cf2: DfCreateFlow, (r3, r4): DfRole,$ $m2 : DataFlowModel \bullet ((r3 = tr(r1))$ $∧ (r4 = tr(r2)) ∧ (cf2, r3, r4 ∈_m m2)$

2- For each read impact in the Impact Model, there exist an read flow in the DataFlow Model.

 $(DfReadFlow <: Arc)$ $tr : ReadFlow \rightarrow DfReadFlow$ $\forall r f1 : ReadFlow,(r1,r2) : ImpRole,$ $m1: ImpactModel \bullet (rf1, r1, r2 \in_m m1)$ \longrightarrow $\exists ! \; rf\,2 : DfReadFlow, (r3, r4) : DfRole,$ $m2 : DataFlowModel \bullet ((r3 = tr(r1))$ $∧ (r4 = tr(r2)) ∧ (rf2, r3, r4 ∈_m m2)$

3- For each update impact in the Impact Model, there exist an update flow in the DataFlow Model.

 $(DfUpdateFlow <: Arc)$ $tr: UpdateFlow \rightarrow DfUpdateFlow$ $\forall u f1: UpdateFlow, (r1, r2): ImpRole,$ $m1: ImpactModel \bullet (uf1, r1, r2 \in_m m1)$ −→ $\exists ! \; uf2 : DfUpdateFlow, (r3, r4) : DfRole,$ $m2 : DataFlowModel \bullet ((r3 = tr(r1))$ $∧ (r4 = tr(r2)) ∧ (uf2, r3, r4 ∈_m m2)$

4- For each delete impact in the Impact Model, there exist an delete flow in the DataFlow Model.

 $(DfDeleteFlow \leq: Arc)$ $tr : DeleteFlow \rightarrow DfDeleteFlow$ $\forall df1 : DeleteFlow, (r1, r2) : ImpRole,$ $m1: ImpactModel \bullet (df1, r1, r2 \in_m m1)$ \longrightarrow $\exists! df2: DfDeleteFlow, (r3, r4): DfRole,$ $m2 : DataFlowModel \bullet ((r3 = tr(r1))$ $∧ (r4 = tr(r2)) ∧ (df2, r3, r4 ∈_m m2)$

It is important to highlight that these agents use the $ImpRoleToDFRole$ and $RoleToDFRole$ agents to translate the end-roles for each flow and participation, respectively, into correspondence ones in the DataFlow Model.

$$
tr: Role \rightarrow DfRole
$$

\n
$$
\forall r1: Role, m1: ImpactModel \bullet (r1 \in_m m1)
$$

\n
$$
\Rightarrow
$$

\n
$$
\exists! r2: DfRole, m2: DataFlowModel
$$

\n
$$
\bullet ((r2 \in_m m2) \land NameOf(r1, r2))
$$

8.2.2 Translating Information Model into Data (Dependency) Model

8.2.2.1 Data Diagram (Model)

A direct one-to-one mapping exists between a given Information Model and the generated Data Model, in which the Information Diagram element in the Information Model is translated directly into a Data Dependency Diagram node. This is expressed in the following rule:

 $tr : InfDiagram \rightarrow DDiagram$ $\forall d1: InfDiagram, \forall e1: InfEntity,$ $\forall a : Relationship, \forall (r1, r2) : Inf Role,$ $m1: InformationModel \bullet ((e1, a, r1, r2 \in_m d1))$ \wedge $(d1 \in_m m1))$ −→ $\exists ! d2 : DDiagram, (r3, r4) : DRole, d : Dependercy,$ $e2 : DEntity, m2 : DataModel \bullet ((e2 = tr(e1))$ \wedge (d = tr(a)) \wedge (r3 = tr(r1)) \wedge (r4 = tr(r2)) $∧ (e2, d, r3, r4 ∈_m d2) ∧ (d2 ∈_m m2)$ \wedge NameOf(d1,d2))

8.2.2.2 Data Entity

Deriving *DEntity* is done by mapping an entity in *Information Model* into a corresponding one in the Data Dependency Model. This translation is carried out by InfEntityToDEntity agent. The rule of this translation is: For each entity in the Information Model, there exists an entity in the Data Model. It can be expressed in FOPL as:

 $tr : InfEntity \rightarrow DEntity$ ∀ ent1 : InfEntity, ∀ att1 : InfAttribute, $m1: InformationModel \bullet ((att1 \in_{p} ent1))$ \wedge (ent1 \in_m m1)) \longrightarrow

∃! ent2 : DEntity, att2 : DAttribute, m2 : DataModel

• ((att2 \in_{p} ent2) \wedge (ent2 \in_{m} m2)

 \land $NameOf(ent1, ent2) \land (att2 = tr(att1)))$

In addition to this, *DEntity* may also correspond to an *Information Model* many-to-many Association. This kind of association is promoted to a full entity in the Data Model. The end-roles are converted into named attributes storing the identities of the *InfEntities* of the association as extra DAttribute model the old end Roles.

This translation is performed by $InfAssocToMergeTable$ agent. It searches in the source model tree for many-to-many *Association* nodes, and then generates a new *DEntity* element in the target model. This includes translating the *Identifiers* of each *InfEntity* into unique DAt tributes within the created entity. These sub-translations are carried out by *IdentifierToUnique* agent. From above, two main roles can be expressed as:

1- For each many-to-many association in the Information Model, there exists an entity in the Data Model.

 $tr : Association \rightarrow DEFINity$

 $\forall a : Association, (obj1, obj2) : InfEntity,$ $(id1, id2)$: Identifier, $m1$: Information Model $\bullet ((a, obj1, obj2 \in_m m1) \land ManyToMany(a)$ \land Links(a, obj1, obj2) \land Id(obj1, id1) $\wedge Id(obj2, id2))$ \longrightarrow

> $\exists ! \; e : \textit{DEntity}, (\textit{att3}, \textit{att4}) : \textit{DAttribute},$ $m2 : DataModel \bullet ((e \in_m m2))$ $\wedge (att3, att4 \in_p e) \wedge NameOf(a, ent)$ \wedge (att3 = tr(id1)) \wedge (att4 = tr(id2))

2- For each identifier of an entity involved in a many-to-many association, in the Information Model, there exists an equivalent unique attribute in the corresponding entity of that association in the Data Model.

 $tr: Identifier \rightarrow DAttribute$ $\forall a : Association, \forall e : InfEntity,$ $\forall id : Identifier, m1 : InformationModel$ • $((a, e \in_m m1) \land ManyToMany(a))$ \land $Id(e, id) \land$ $(Source(a, e) \oplus Target(a, e)))$ \longrightarrow ∃! ent : DEntity, ∃! att : DAttribute, m2 : DataModel • ((ent ϵ_m m2) \wedge (att ϵ_p ent) \land Unique(att) \land NameOf(a, ent)))

It is worth emphasising an interesting case in regard to *disjoint Generalisation*, where it is assumed that Child entities consist of the attributes of the Parent Entity, added to its original attributes. At the end the Parent entity is deleted from the model and a new version of its Child entities appear in the Data Dependency Model.

From this, a new rule might be stated: For each disjoint generalisation, which connects child entities to a parent one, in the Information Model, there exists an equivalent entity to the child entities that have the attributes of the parent. This can be written in logic as:

 $tr : Generalisation \rightarrow DEntity$ $\forall g : Generalisation, \forall e1, e2 : InfEntity,$ ∀ a1, a2 : InfAttribute, m1 : InformationModel \bullet ((g, e1, e2 \in_m m1) \land (a1 \in_p e1) $∧ (a2 \epsilon_p e2) ∧ Connects(g, e2, e1)$ \land *Disjoint(q)* \land (e1 \neq e2) \land (a1 \neq a2)) −→ ∃! ent : DEntity, ∃ att1, att2 : DAttribute, m2 : DataModel • ((ent ϵ_m m2) \wedge (att1, att2 ϵ_p ent) ∧ NameOf(e2, ent) \wedge (att1 = tr(a1)) \wedge (att2 = tr(a2)))

8.2.2.3 Data Attribute

A straightforward mapping exists between Information Model Attribute to generate equivalent Data Dependency DAttribute, in which each Attribute element is directly translated into a DAttribute. The following logic formula expresses this:

 $tr : InfAttribute \rightarrow DAttribute$ ∀ att1 : InfAttribute, m1 : InformationModel • $(\text{att1 } \in_m m1)$ \longrightarrow ∃! att2 : DAttribute, m2 : DataModel • $((att2 \in m) \land NameOf(att1, att2))$

8.2.2.4 Data Dependency

A Dependency in the Data Model might be derived by translating one of the three Information Model concepts: many-to-one association, (overlapping) generalisation, and composition. In regards to many-to-one association, the InfAssocToDDependency rule (agent) is responsible for translating each detected Association into a Dependency in which the many always depend on the one. Any many-to-many association must be promoted into a into a pair of many-toone associations first, then the multiplicity constraints must also be translated (see earlier rule above).

From the above, the first translating rule can be summarised as: For each many-to-one association between two entities in the Information Model, there exists an equivalent dependency in between the equivalent entities the Data Dependency Model where the entity on the many-side depends on the entity on the one-side. This can be expressed in FOPL as:

 $tr : Association \rightarrow Dependercy$ \forall e1, e2 : InfEntity, \forall a : Association, $m1: InformationModel \bullet ((a, e1, e2 \in_m m1))$ \land $ManyToOne(a) \land Connect(a, e1, e2) \land (e1 \neq e2)$

−→ $\exists ! d : Dependercy, e3, e4 : DEntity, m2 : DataModel$ • $((d, e3, e4 \in_m m2) \land Connects(d, e3, e2))$ \land NameOf(e1, e3) ∧ NameOf(e2, e4) ∧ $Depends(e3, e4) \wedge (e3 = tr(e1)) \wedge (e4 = tr(e2))$

In addition to this, every (overlapping) Generalisation in the Information Model is resolved by making the specific entities depend on the general one. This mapping is carried out by the InfGenToDDependency agent. The second rule of translation can be defined as: For each overlapping generalisation between entities in the Information Model, there exists an equivalent dependency relationship between the equivalent entities in the Data Dependency Model. This can be written in logic as:

−→

 $tr : Generalisation \rightarrow Dependency$ $\forall e_1, e_2 : InfEntity, \forall g : Generalisation,$ $m1: InformationModel \bullet ((g, e1, e2 \in_m m1))$ $\land \neg Disjoint(g) \land Connect(g, e1, e2) \land (e1 \neq e2)$

 $\exists ! d : Dependency, e3, e4 : DEntity, m2 : DataModel$

• $((d, e3, e4 \in m \space m2) \land Connects(d, e3, e4)$ \land (NameOf(e1, e3) ∧ NameOf(e2, e4) ∧ $Depends(e3, e4) \wedge (e3 = tr(e1)) \wedge (e4 = tr(e2)))$

Furthermore, every (not total) Composition in the Information Model is resolved by making the whole depend on the parts. The *InfComToDDependency* rule is responsible for proceeding with this mapping from *Composition* into *Dependency*. From this, the third rule for deriving Dependency can be defined as: For each composition between entities in the Information Model, there exists an equivalent dependency relationship between the equivalent entities in the Data Dependency Model. This can be written in logic as:

 \longrightarrow

 $tr: Composition \rightarrow Dependent$

 $\forall e_1, e_2 : InfEntity, \forall c : Composition,$ $m1: InformationModel \bullet ((c, e1, e2 \in_m m1))$ $\wedge \neg Total(c) \wedge Connect(c, e1, e2) \wedge (e1 \neq e2)$

 $\exists ! d : Dependercy, e3, e4 : DEntity, m2 : Data Model$

• $((d, e3, e4 \in_m m2) \land Connects(d, e3, e4))$

∧ $NameOf(e1, e3)$ ∧ $NameOf(e2, e4)$ ∧

 $Depends(e3, e4) \wedge (e3 = tr(e1)) \wedge (e4 = tr(e2))$

8.2.3 Transforming (initial) DataFlow Model into Detailed DataFlow Model

8.2.3.1 Detailed DataFlow Diagram

This one-to-one mapping aims at applying a task decomposition strategy to *DataFlow* tasks that are connected to more than one flow. From that, For each DataFlow Model, there exists a detailed DataFlow one generated from it. This can be formalised as follows:

> $tr : DfDiagram \rightarrow DfDiagram$ $\forall d1 : DfDiagram, \forall b1 : Boundary,$ $m1 : DataFlowModel \bullet ((b1 \in_m d1))$ \wedge (d1 \in_m m1) −→ $\exists ! \ d2 : DfDiagram, b2 : Boundary,$ $m2 : DataFlowModel \bullet ((b2 \in_m d2))$ $∧ (d2 \in_m m2) ∧ (b2 = tr(b1)))$

8.2.3.2 Detailed DataFlow Boundary

Each DfTask that is connected to more than one Flow is translated into a DfBoundary in the detailed DataFlow Model. This can be defined as: for each dataflow task that is connected to more than one flow, there exists one equivalent boundary in the detailed dataflow model. The following FOPL expresses this formally:

 $tr: DfTask \rightarrow DfBoundary$ \forall t1 : DfTask, f1, f2 : Flow, a1 : DfActor, $m1: DataFlowModel \bullet ((t1, f1, f2, a1 \in_m m1))$ $∧(f1 \neq f2) ∧ (Source(f1, t1) ⊕$ $Target(f1,t1)) \wedge (Source(f2,t1)$ $\oplus Target(f2,t1))$ −→ $\exists ! b : DfBoundary, t2 : DfTask, f3, f4 : Flow,$ $a2: DfActor, m2: DataFlowModel$ • $((b \in_m m2))$ \land NameOf(t1,b) \land (f3 = tr(f1)) \land $(f4 = tr(f2)) \wedge (t2 = tr(t1)) \wedge (a2 = tr(a1))$

8.2.3.3 Detailed DataFlow Task

As a general rule, each *DfTask* in an initial DFD model is translated into an equivalent *DfTask* in the detailed DFD one. This mapping is applicable when the DfTask is the source of either create, delete or update, or when it is target of either input or read. The following rule expresses the cases when the flow is either a *create, delete* or *update* flow:

 $tr: DfTask \rightarrow DfTask$ $\forall f1 : Flow, t1 : DfTask, m1 : DataFlowModel$ • $((t1, f1 \in_m m1) \wedge (DfCreate(f1))$ $\oplus DfDelete(f1) \oplus DfUpdate(f1)) \wedge (Source(f1, t1))$ −→

> $\exists !\ t2 : DfTask, m2 : DataFlowModel$ • $((t2 \in_m m2) \land NameOf(t1, t2))$

In addition to this, the following rule expresses the mapping when the flow is either an *input* and a read flow.

$$
tr: DfTask \rightarrow DfTask
$$

\n
$$
\forall f1: Flow, t1: DfTask, m1: DataFlowModel
$$

\n• ((t1, f1 ∈_m m1) ∧ (DfRead(f1)
\n
$$
\oplus DfInput(f1)) ∧ (Target(f1, t1)))
$$

\n
$$
\rightarrow
$$

\n∃! t2: DfTask, m2: DataFlowModel
\n• ((t2 ∈_m m2) ∧ NameOf(t1, t2))

Any *DfTask* that is connected to more than one *Flow* will be decomposed into a number of atomic tasks based on the number of flows to which it is connected. The following rule expresses this mapping when the task is a source in more than one flow:

$$
tr: DfTask \rightarrow DfTask
$$

\n
$$
\forall f1, f2: Flow, t1: DfTask, m1: DataFlowModel
$$

\n• ((t1, f1, f2 ∈_m m1) ∧ (f1 ≠ f2)
\n∧ (Source(f1, t1) ⊕ Source(f2, t1)))
\n
$$
\rightarrow
$$

\n
$$
\exists t1a, t1b: DfTask, m2: DataFlowModel
$$

\n• ((t1a, t1b ∈_m m2) ∧ AtomicTaskOf(ta, t1a)
\n∧ AtomicTaskOf(t1, t1b)
\n∧ (Source(f1, t1a) ⊕ Source(f2, t1b)))

It is worth saying that the rule expressed above might be rewritten by considering the task to be target in more than a flow, or when it is source in one flow and target in another one. The following rule expresses the second case:

 $tr: DfTask \rightarrow DfTask$ $\forall f_1, f_2 : Flow, t1 : DfTask, m1 : DataFlowModel$ • $((t1, f1, f2 \in_m m1) \wedge (f1 \neq f2))$ \land $(Target(f1, t1) \oplus Target(f2, t1))$ −→

 \exists t1a, t1b : DfTask, m2 : DataFlowModel

• $((t1a, t1b \in_m m2) \land AtomicTaskOf(t1, t1a)$

 \wedge AtomicTaskOf(t1,t1b)

 \wedge $(Target(f1, t1a) \oplus Target(f2, t1b)))$

A priority scoring mechanism is introduced for prioritising each atomic task in the detailed DFD. This technique is applied during decomposing initial DFD tasks into their atomic ones.

8.2.3.4 Detailed DataFlow Object

A DfObject is passed directly to the detailed DataFlow Model during the transformation step. This mapping is carried out by the ArcToDfObject agent which maps the DfObject at both ends both ends of each Flow. The following rule demonstrates the transformation in the case of create, delete and update flows only.

−→

−→

 $tr : DfObject \rightarrow DfObject$

 $\forall f1: Flow, obj1: DfObject, m1: DataFlowModel$ $\bullet((obj1, f1 \in_m m1) \land (DfCreate(f1) \oplus DfDelete(f1))$ $\oplus DfUpdate(f1)) \wedge (Target(f1, obj1))$

 $\exists! obj2: DfObject, m2: DataFlowModel$

• $((obj2 \in_m m2) \land NameOf(obj1, obj2))$

By considering DfReadFlow:

 $tr : DfObject \rightarrow DfObject$

 \forall f1 : Flow, obj1 : DfObject, m1 : DataFlowModel $\bullet ((obj1, f1 \in_m m1) \land (DfRead(f1) \land (Source(f1, obj1)))$

 $\exists! obj2: DfObject, m2: DataFlowModel$

• $((obj2 \in m) \land NameOf(obj1, obj2))$

8.2.3.5 Detailed DataFlow Actor

A DfActor is passed directly to the detailed DataFlow Model during the transformation step. This mapping is carried out by the $ArcToDfActor$ agent by detecting DfActor at the both ends of each Flow. The following formula illustrates the rule in the case of DfInputFlow:

 $tr : DfActor \rightarrow DfActor$ $\forall f_1 : DfInputFlow, a1 : DfAction, m1 : DataFlowModel$ $\bullet ((a1, f1 \in_m m1) \land (Source(f1, a1))$ \longrightarrow $\exists ! a2 : DfActor, m2 : DataFlowModel$ • $((a2 \in_m m2) \land (NameOf(a1, a2)))$ by considering the DfOutputFlow:

 $tr : DfActor \rightarrow DfActor$ $\forall f_1 : D fOutput Flow, a1 : D f Actor, m1 : DataFlow Model$ • $((a1, f1 \in_m m1) \wedge (Target(f1, a1))$ −→ $\exists ! \ a2 : DfActor, m2 : DataFlowModel$ • $(a2 \in m \space m2)$

8.2.3.6 Detailed DataFlow Flows

All Flow types are passed directly to the detailed DataFlow Model the transformation step except DfUpdateFlow. This can be expressed in the following general FOPL law:

```
tr : Flow \rightarrow Flow\forall f1 : Flow, m1 : DataFlowModel • ((f1 \in_m m1)
        \wedge (DfCreateFlow(f1) \oplus DfReadFlow(f1))\oplus DfDeleteFlow(f1) \oplus DfInputFlow(f1)
                               \oplus DfOutputFlow(f1))
                                                    −→
                                                        \exists ! \; f2 : Flow, m2 : DataFlowModel• ((f2 \in_m m2) \wedgeFlowTypeOf(f1, f2))
```
In the case of $DFUpdateFlow$, it is decomposed into $DfReadFlow$ and $DfWriteFlow$ and attached to the decomposed $DfTask$ connected to that $DfUpdateFlow$. This can be expressed as:

 $tr: UpdateFlow \rightarrow DfReadFlow \times DfWriteFlow$ \forall f1 : Flow, m1 : DataFlowModel • $((f1 \in_m m1) \wedge (UpdateFlow(f1))$ \longrightarrow $\exists ! \text{ } f2 : D \text{ } f ReadFlow, f3 : D \text{ } fWriteFlow,$ $m2 : DataFlowModel \bullet (f2, f3 \in_m m2)$

Priority Scoring Algorithm of Tasks

Two prioritising steps are used to extract (predict) the internal behaviour of each *DfBound*ary in the detailed DFD: analysing types of tasks and analysing data on flows. In the first step, Partial Dependency from the task execution perspective is considered a strategy for prioritising all atomic tasks at the detailed DFD model. For each DFD boundary, tasks are classified into four levels of priority based on their types. The task that has a high priority score is executed before the lower ones. It is assumed that the input tasks have the highest priority in a boundary with a score of 4. The read tasks come next with a score of 3. Whereas, create, delete and write come after with a score of 2. The output task, at the end, has a score of 1 (see Table 8.1).

Table 8.1: Priority ranking of different types of task

In the second step, data on flows is checked in order to be sure in which order tasks might be executed. For example, suppose we have a piece of data (x) on an *input* flow and the same (x) on a create flow, it can be predicted (extracted) that x has to be input before it can be used in a create task. As a result of these two steps, additional DfReadFlows are generated to connect tasks that have higher priority scores to the lower ones, for each detailed DFD boundary.

8.2.4 Translating DataFlow Model into (Screen) State Model

8.2.4.1 State Diagram (model)

A *StDiagram* element in translated from an existing *DfDiagram* node in the detailed *DataFlow* Model, in which For each detailed DFD, there exists a diagram in the State Model. This can be expressed as:

 $tr : DfDiagram \rightarrow StDiagram$ $\forall d1 : DfDiagram, b1 : DfBoundary,$ $m1 : DataFlowModel \bullet ((b1 \in_m d1))$ \wedge $(d1 \in_m m1))$ \longrightarrow ∃! d2 : StDiagram, b2 : StBoundary, m2 : StateModel • $((b2 \in_m d2) \wedge (d2 \in_m m2))$ \land $NameOf(d1, d2) \land (b2 = tr(b1)))$

8.2.4.2 Boundary

Each *StBoundary* corresponds to a *DfBoundary* in the *DataFlow Model*. This mapping is acheived using the *DfdBoundaryToStBoundary* agent that provides a one-to-one translation step. The main rule of this mapping is: For each boundary in the DataFlow Model, there exists a boundary in the State Model. The logic expression of this rule is:

 $tr: DfBoundary \rightarrow StBoundary$ \forall m1 : DataFlowModel, t : DfTask, f : Flow, $b1: DfBoundary \bullet ((t, f \in_m b1) \wedge (b1 \in_m m1))$ −→ $\exists ! \ m2 : StateModel, b2 : StBoundary, s : State,$ ts : Trainsition • $((s, ts \in_m b2) \land$ $(b2 \in_m m2) \land NameOf(b1, b2) \land$ $(s = tr(f)) \wedge (ts = tr(t)))$

Is it worth saying that each generated boundary has a number of states that describe the successful case of the business scenario. The step of adding *Error States* for handling possible failure scenarios is a separate *one-to-one* translation step. The *StBoundaryToStBoundary* is responsible for this by traversing each *State* in the *StBoundary* to produce and attach a related error reporting state to the boundary.

8.2.4.3 State

Waiting State in the State Model corresponds to an *DfInputFlow* element in the DataFlow Model, in which each input flow is translated into a *State* that is waiting for receiving some inputs. These inputs, which are equivalent to the data on DfInputFlow, are translated into internal variables (Variable) for that state. This mapping is achieved by the $ArcToWaiting State$ translator for each accepted source element in the DataFlow Model.

One the other hand, *Ready State* is derived by translating the other kinds of flows, namely: DfCreateFlow, DfReadFlow, DfUpdate, DfOutputFlow and DfWriteFlow. This translation is done by the ArcToReadyState agent for each accepted Flow in the DataFlow Model. From the above, a number of basic transformation rules are formed:

1- For each Input flow in the DataFlow Model, there exists a State in the State Model that is waiting for receiving some external inputs.

 $tr: DfInputFlow \rightarrow State$ $\forall m1 : DataFlowModel, f : DfInputFlow$ • $(f \in_m m1)$ −→ ∃ s : State, m2 : StateModel • $((s \in_m m2) \land Waiting(s))$

2- For each flow that is either Create, Read, Delete, Write or Output flow in the DataFlow Model, there exists a State in the State Model that is ready to fire an internal system operation or funtion.

 $tr : Flow \rightarrow State$ \forall m1 : DataFlowModel, f : Flow • ((f \in_m m1) $\wedge (DfCreateFlow(f) \oplus DfReadFlow(f) \oplus$ $DfDeleteFlow(f) \oplus DfWriteFlow(f) \oplus DfOutputFlow(f))$ −→ ∃ s : State, m2 : StateModel • $((s \in_m m2) \land Ready(s))$

8.2.4.4 Transition

In order to obtain Transition elements, a complex merge translation step is required. Thus, a DfTaskToTransition agent is designed to fit this need. This translator must be fired after deriving all *States* in the *State Model. DfTaskToTransition* takes each flow with its attached end Roles to create a new Transition.

Therefore, the mapping rule of this agent is: For each source in a flow in the DataFlow Model and its translated state in the State Model, there exists a transition in the State Model, where the source of the transition is the state that is pased as an argument of this agent. This can be expressed in FOPL as:

 $tr : DfTask \rightarrow Transition$ \forall m1 : DataFlowModel, t : DfTask, $m2: StateModel, f: Flow, s: State \bullet ((f \in_m m1))$ \wedge (s \in_m m2) \wedge Source(t, f) \wedge (s = tr(f)) −→

 $\exists ! t : Transition \bullet ((t \in_m m2) \land Source(t, s))$

8.2.4.5 Action

Each transition has an action attribute refering to its type. For example, if the Flow was a DfCreateFlow. then the action type will be create. This Action takes some arguments that are eqivalent to the internal Variable of the State passed to the DfTaskToTransition agent. The body of the action might have assignment statements if the Flow is either DfCreateFlow or DfWriteFlow, whereas it may have filter statements if the Flow is DfReadFlow.

8.3 Transformation Rules at the Analysis to Design Phase

8.3.1 Translating State into GUI Model

8.3.1.1 GUI Diagram

A GuiDiagram element in translated from an existing StDiagram node in the State Model, in which for each State diagram, there exists a diagram in the GUI Model. This can be expressed as:

 $tr : StDiagram \rightarrow GuiDiagram$ ∀ d1 : StDiagram, b1 : StBoundary, $m1: StateModel \bullet ((b1 \in_m d1))$ \wedge $(d1 \in_m m1)$ −→ $\exists ! d2 : GuiDiagram, b2 : GuiBoundary, m2 : GUIModel$ • $((b2 \in_m d2) \wedge (d2 \in_m m2))$ \land $NameOf(d1, d2)) \land (b2 = tr(b1))$

8.3.1.2 GUI Window

Each Window corresponds directly to a State in the State Model. This straightforward mapping is implemented in $StateToWindow$ agent. This can be expressed in logic as:

$$
tr: State \rightarrow Window
$$

\n
$$
\forall s: State, \forall v: Variable, m1: StateModel
$$

\n• ((s ∈_m m1) ∧ (c ∈_p s))
\n
\n
$$
\exists! w: Window, c:Widget, m2: GUIModel
$$

\n• ((d2 ∈_m m2) ∧ (c ∈_c w)
\n∧ NameOf(s, w) ∧ (c = tr(v)))

Variables in the State is translated into Widget by an agent related to the type of the state. If the state is *waiting*, then the variables will be handled by *svariableToTextfield* that translates these variables into entry text fields. On the other hand, If the state is ready, then the variables will be handled by *svariableToLabel* that translates these variables into displaying labels.

8.3.1.3 GUI Button Widget

Button control appears in all types of screens. The basic mapping rule for obtaining Button can be defined: Each generated window has a single button. Unlike other widgets, deriving the Button control requires such a complex translation step than other widgets, because it has an event that triggers (invokes) a particular stored procedure in the database. The following FOPL statement express the main rule for deriving *Button*:

```
tr : State \rightarrow Window\forall s: State, m1: StateModel \bullet ((s \in_m m1))−→
                                                 \exists ! w : Window, b : Button, m2 : GUIModel• ((w \in_m m2) \land (b \in_c w) \land NameOf(s, w))
```
8.3.1.4 GUI Label Widget

Label widget appears in all type of windows, displaying the title, error message, names of textfields and more. The basic mapping rule for obtaining Label can be defined: Each generated window has some label. This can be expressed as:

> $(Label \leq: Widget)$ $tr : Variable \rightarrow Label$ $\forall s : State, \forall v : Variable, m1 : StateModel$ \bullet ((s \in_m m1) \wedge (v \in_p s)) \longrightarrow $\exists ! w : Window, l : Label, m2 : GUIModel$ • $((w \in_m m2) \land (l \in_c w))$ \wedge NameOf(s, w))

8.3.1.5 GUI Textfield Widget

Textfield widget appears in waiting windows only for receiving external inputs. The transformation rule for deriving Textfield can be expressed in english as: Each generated wating window has some text field. The associated FOPL formula can be expressed as:

 $(Textfield \leq: Widget)$ $tr : Variable \rightarrow Textfield$ ∀ s : State, ∀ v : V ariable, m1 : StateModel \bullet ((s \in_m m1) \land (v \in_p s)) \longrightarrow $\exists ! w : Window, t : Textfield, m2 : GUIModel$ • $((w \in_m m2) \wedge (t \in_c w))$ \wedge $NameOf(s, w)$

8.3.2 Translating DataFlow and Data Model into the DBQ Model

8.3.2.1 Table

Deriving Table is achieve directly by translating an entity, in an Data Dependency Model, into a corresponding logical tabel, in the Database and Query Model. This translation is carried out by DEntityToTable agent. The rule of this translation is: For each entity in the Data Model, there exists a table in the Database Model. It can be expressed in FOPL as:

 $tr : DEntity \rightarrow Table$ $\forall e : DEntity, \ \forall a : DAttribute, m1 : DataModel$ • (($e \in_{m} m1$) \wedge ($a \in_{p} e$)) \longrightarrow $\exists ! t : Table, c : Column, m2 : DatabaseModel$

• ((t \in_m m2) \wedge (c \in_p t) \wedge NameOf(e, t) \wedge $(c = tr(a)))$

8.3.2.2 Column

A Column corresponds to a DAttribute of a Data Model DEntity. This mapping is performed by the DAttributeToColumn agent that translate every attribute into a column of a relational table. The main rule of this translation can be stated as: For each attribute in such an entity in the Data Dependency Model, there exists a corresponding column in the equivalent table to that entity. this can be written in FOPL as:

```
tr: DAttribute \rightarrow Column\forall e : DEntitu, \forall a : DAttribute, m1 : DataModel\forall t : Table, m2 : DatabaseModel• ((e \in_m m1) \wedge (a \in_p e) \wedge (t = tr(e)))
                                                       \longrightarrow∃ c : Column
                                                            • ((c \in_p t) \land NameOf(t, c))
```
It is significant to mention that some column might be ranged with upper and lower values. This is translated into SQL Triggers at the code generation step. There is no element in the DBQ model that refers directly to *triggers*, this because the variation of handling *range constraint* in different database vendors.

8.3.2.3 Foreign Key

In order to construct a Table with a foreign key to represent its dependency on another table, the agent DependToReferTable is used to do this task. It invokes the DEntityToTable agent, described above to translate the entity at the many side end-role of the dependency into a table. Additionally it invokes *DRoleToForeignKey* aget to translate the PK of the entity at the *one* side end-role of the dependency into a foreign key attached to the table.

8.3.2.4 Stored Procedure

A *Procedure* is derived as a result of merging concepts from *DataFlow* and *Data Model*. As the DBQ procedure is designed to bind a SQL statement (SELECT, UPDATE, DELETE and INSERT), there are four kinds of agents each one is responsible for translating a particular type of Flow in a DataFlow Model into a Procedure. These translators are: CreateFlowToStored-Procedure, DeleteFlowToStoredProcedure, WriteFlowToStoredProcedure, ReadFlowToStoredProcedure.

The CreateFlowToStoredProcedure agent takes a DfCreateFlow along with a DataModel as arguments to produce a *Procedure* that consists of a *Create* element, and possibly a *Query* one. The Data Model is used to access all *Columns* of the *DEntity* that matches the one appears at the target end of that DfCreateFlow. Then, each piece of data on the DfCreateFlow with its related *Columns* in the *DEntity* forms a single *assignment statement* using the element *Operator*.

In the same context, the ReadFlowToStoredProcedure agent takes a DfReadFlow and a DataModel as parameters of this translator. It then produces a *Procedure* that contains a *Query* and Project node. The Data Model is used to access all Columns of the DEntity that matches the one appears at the source end of that *DfReadFlow*. Then, each data item on the *DfReadFlow* is matched to its associated *Columns* in the *DEntity* to be considered a *child* of the *Project* element. (filter)

Regarding the *DfDeleteFlow*, the *DeleteFlowToStoredProcedure* agent, which is responsible for this translation, takes a DfDeleteFlow with a DataModel as arguments in order to produce a Procedure that has a Delete element, and possibly a Query one. The Data Model is used to access all Columns of the DEntity that matches the one appears at the target end of that DfDeleteFlow.

Similarly, the WriteFlowToStoredProcedure agent is responsible for translating DfWriteFlow into the related DBQ Procedure. It takes a DfWriteFlow and a Data Model to construct a stored procedure *Procedure* with an *Update* and *Query* element. Data on the *DfWriteFlow*, which is interpreted as data to be stored, is converted into an assignment statement.

8.4 Alternative Translation Step For Generating the Information Model

Apart from the chain of model transformations that are designed to produce basic information systems, there is a separate translation step that might be applied independently for constructing intermediate artefacts. The following section presents the rules for translating a pre-defined Impact Model into a partial Information Model.

8.4.1 Translating the Impact Model into an (initial) Information Model

8.4.1.1 Information Diagram (Model)

A Diagram element is translated from an existing ImpDiagram node in the Impact Model, in which for each Impact Diagram, there exists an initial Information Model (diagram). This can be expressed as:

> $tr: ImpactModel \rightarrow InformationModel$ ∀ m1 : ImpactModel −→ ∃! m2 : InformationModel

8.4.1.2 Information Entity

Entities in the Information Model are derived automatically by a direct one-to-one mapping rule. This can be stated as: For each object in every boundary in the Impact Model, there exists an entity in the Information Model. This can be expressed in FOPL as:

 $tr : ImpBoundary \rightarrow InfEntity$ $\forall b : ImpBoundary, \forall obj : ImpObject,$ $m1: ImpactModel \bullet ((obj \in_m b) \land (b \in_m m1))$ −→ ∃! e : InfEntity, m2 : InformationModel • ((e \in_m m2) \wedge (obj = tr(e)))

8.4.1.3 Information Attribute

Attributes in the Information Model are not derived automatically from the Impact Model.

8.4.1.4 Information Association

Association types are derived directly from pre-defined Flow types in the Impact Model. This can be expressed in the following general FOPL law:

−→

 $\forall b : ImpBoundary, \forall f : Flow, m1 : ImpactModel$ $\forall (obj1, obj2): ImpObject \bullet (Connect(f, obj1, obj2)) \wedge$ $((f, obj1, obj2 \in_m b) \wedge (b \in_m m1) \wedge (obj1 \neq obj2))$

> $\exists a : Association, \forall (e1, e2) : InfEntity,$ $m2: InformationModel \bullet ((a, e1, e2 \in_m m2)) \wedge$ $(obj1 = tr(e1)) \wedge (obj2 = tr(e2)) \wedge \wedge$ $(a = tr(f))$ Connect (a, e_1, e_2)

8.5 Outlook on the Chapter

The chapter has presented, in-depth, the rules of transformations that are applied to source models in each tranformation step, within the three development phases of BUILD. These rules were formalised and expressed using First-Order Predicate Logic. For each rule, source elements and the generated target ones are highlighted, including constraints that must be satisfied, in the form of unary and n-ary predicates, in order to complete the execution of that rule.

9 Case Studies

"The best time to plan an experiment is after you've done it" R. A. Fisher

9.1 Context

This chapter illustrates the results of the work presented in the previous chapters by introducing a real-world enterprise information system case study (*University Administration System*). Three experiments are designed in which each involves a particular part of the system. The construction of system models, expressed using the μ ML notation, at each development level of BUILD is highlighted. Then, significant translation decisions, which are interesting, are highlighted and discussed for each development phase.

It is also worth mentioning that each μ ML model is represented through BUILD as an abstract syntax tree (AST) , written in Java node classes, that mimics the structure of that model, allowing the creation of a variety of memory models representing different systems.

In technical detail, the current version of BUILD does not include a graphical editor to enable the automatic shift from the graphical notation of models, sketched by business end-users, into their corresponding underlying XML representations (AST). Therefore, in each experiment, all μ ML requirements models are constructed directly in memory, using the builder-API supplied by the metamodel classes. This is just for the purposes of conducting these experiments; in live usage, the initial models would be supplied as XML files, to be unmarshalled by the $JAST$ parser.

9.2 Overview of the University Administration System

The University Administration System is a common information system that might be developed within academic institutions. From the domain user perspective, the system can be divided into two major parts (sub-systems), namely, Module Management System and Student Enrolment System. The following subsections introduce each part of the system separately, including its evolution during the BUILD development stages.

9.3 Overview of the Module Management Sub-system

The Module Management System enables the users to perform real-world tasks, such as obtaining the full description of a module, updating some details of a module, adding a new module to the system, and removing an existing module. The system has two main types of actor, *Staff* and Student, where each actor role has different tasks to perform. Staff are responsible for adding, removing, and modifying modules in the system, one at a time, while *Students* are able to search by a module code to retrieve the full description of that module only. Each business task here is performed on a single module at a time; for example, staff can add one module to the system during the execution of the Add Module task, and the rest is likewise.

9.3.1 Information System Representation at the Requirements Sketching Phase

In this section, the μ ML requirements models that represent the Module Management System are demonstrated graphically. The modelling (construction) activities are then discussed, highlighting critical concepts that are captured in each model. The complete underlying XML representation of all models at this stage is listed in Appendix A.

The development process starts when a business-user expresses the structure of the *Module Management System* using the μ ML Task Model. The functions of the system can be categorised into two main jobs which are managing modules and displaying modules description. Thus, they can be expressed as two main Tasks called Manage Module and See Description respectively. The Manage Module consists of three independent business tasks namely, Add Module, Delete Module and Modify Module. A white diamond is used to express the part of relationship that has independent subtasks.

The human-computer interaction is represented via placing some Actors connected to business tasks they are involved in. The Student interacts with the See Description task by supplying it with a module code to get the full detail of the corresponding module, whereas, Staff might interact with all business tasks within the Manage Module task. The following Task Model, Figure 9.1 represents the structure of the Module Management System.

Figure 9.1: Task model. Module Management System

With regard to the behaviour of the system, the μ ML Impact Model is used to capture the internal interaction between the business tasks and the system objects. Each task, captured in the Task Model, has a different kind of impact on the Module entity. The Add Module task connects to the Module by a create flow, the Delete Module task connects to the object by a

delete flow, the *Modify Module* task connects to the object by a bi-directional *update* flow, and finally, the Read Module task connects to the object by a read flow. The following figure, Figure 9.2 is the Impact Model of the Module Management System.

Figure 9.2: Impact model. Module Management System

The third source model at this stage of development is the Information Model. The model consists of an entity called Module that has four attributes, namely, code, title, desc and credit. The following figure (Figure 9.3) illustrates this.

Module
code: Integer
title: String
credit: Integer
desc: String

Figure 9.3: Information model. Module Entity.

9.3.2 Running the Experiment on the BUILD Framework (1)

Purpose: Generate a JDBC Java Swing application with MySQL backend database. Critical Feature: Adopting appropriate decisions for constructing a menu that handles the optional functions of the system.

Input: Three μ ML models: Task, Impact and Information Model.

Output: Java classes (*.java) and a MySQL script file (*.sql).

Running Environment: Eclipse.

9.3.2.1 Construction of Requirements Models

The first requirements model considered by business-users in the development process is the Task Model. As there is a custom AST representing Task Model in BUILD, the memory model of the system Task Model is constructed out of the corresponding Java node classes. For example, the Diagram node class is used to initialise a new model, and the Task node class is used to define new all business tasks of that model. Listing 9.1 below illustrates a snapshot of the manual creation of the Task Model of the Module Management System. The complete code is presented in Appendix A.

```
Listing 9.1: Construction of the Task Model
```

```
1 Diagram taskModel = new Diagram();
2 Boundary boundary = new Boundary("Module Management");
3
4 \mid Task manage = new Task ("Manage Module");
5
6 Task add = new Task("Add Module");
7
8 Actor actor1 = new Actor("Staff");
9
10 Composition comp = new Composition();
11 comp.addRole(new mde.task.model.Role("manage", manage));
12
13 Participation link1 = new Participation();
14 link1.addRole(new mde.task.model.Role("staff",actor1));
15 link1.addRole(new mde.task.model.Role("add", add));
```
Similar to the creation of Task Model, business users design the Impact Model of the system using node classes in the *Impact Model Package*. The *ImpDiagram* class is used to initialise a new Impact Model in memory; ImpTask and ImpObject are used to define all business tasks and actors for their system respectively, and the rest in likewise. Listing 9.2 below demonstrates a snapshot of the construction of Impact Model, where the complete code is presented later in Appendix A. It can be seen that names of boundaries and tasks are equivalent to previously designed boundaries and tasks in the Task Model, Listing 9.1 above.

Listing 9.2: Construction of the Impact Model

```
1 ImpDiagram ImpactModel = new ImpDiagram();
2 ImpBoundary impboundary = new ImpBoundary("Module Management");
3
4 ImpTask impManage = new ImpTask("Manage Modules");
5 ImpTask impAdd = new mde.impact.model.ImpTask("Add Module");
6
7 ImpObject impObj1 = new ImpObject("Module");
8
9 ImpCreateFlow cf = new ImpCreateFlow();
10 | ImpRole impcf1 = new ImpRole("module", impObj1);
11 ImpRole impcf2 = new ImpRole("add", impAdd);
12 cf.addImpRole(impcf2);
13 cf.addImpRole(impcf1);
```
The final model at this stage, created manually by user, is the *Information Model*. The following listing (Listing. 9.3) illustrates a snapshot of the related part of the Module Management System. It is designed using the Information Model Package of the BUILD framework. The class Entity is used to define the back-end entities of the system.

Listing 9.3: Construction of the Information Model

```
1 mde.information.model.Diagram informationModel =
2 new mde.information.model.Diagram();
3
4 Entity moduleEntity = new Entity("Module");
5 Attribute attr12 = new Attribute("code",
6 new Type("Integer")).setIdentifier(true);
7 Attribute attr13 = new Attribute("title", new Type("String"));
8 Attribute attr14 = new Attribute("credit", new Type("Integer"));
  9 Attribute attr15 = new Attribute("desc", new Type("String"));
```
Once the initial requirements models have been constructed in memory, the class ASTWriter, the class *ASTWriter* is used for marshalling each abstract syntax tree, built in memory to an XML file. The method *writeDocument* is used to serialise each memory model into its related XML file. The idea behind having all models stored into XML files is to apply a XML model inspection as a strategy of evaluating the quality of produced artifacts, as well as verifying and checking the correctness and completeness of translation results. This is discussed later in Chapter 9.

9.3.2.2 Information System Representation at the Analysis Phase

The Requirement-to-Analysis model transformation step leads to creating a number of intermediate analysis models of the Module Management System. These models are: DataFlow, State Model and Data Dependency Model. In this section, the μ ML analysis models that express the Module Management System are demonstrated graphically, including some key automatic transformation decisions made by the approach. The complete underlying XML representation of all models at this stage is listed in Appendix A.

While the *Module Management System* deals with a single business entity, there is no significant transformation decision, from Information to Data Dependency Model, in this case study. The following figure (Figure 9.4) demonstrates the produced Data Dependency Model at the Analysis Phase of BUILD.

Module
code: Integer
title: String
credit: Integer
desc: String

Figure 9.4: Data model. Module Entity

Furthermore, terminal tasks in both Task and Impact Model appear in the initial DataFlow Model, as a result of the translation shift from the requirements phase to the analysis one. These tasks are: See Descriptions, Add Module, Modify Module and Delete Module. Moreover, business entities and system actors in both Impact and Task Model, respectively, also appear in the *DataFlow Model*, in which each complete business transaction from an actor to object is expressed in the model (Figure 9.5). These transactions can be listed as:

- The See Description task is connected to a Student actor by an input flow and to a Module object by a read flow.
- The Add Module task is connected to a Staff actor by an *input* flow and to a Module object by a create flow.
- The Modify Module task is connected to a Staff actor by an input flow and to a Module object by an update flow.
- The Delete Module task is connected to a *Staff* actor by an *input* flow and to a *Module* object by a delete flow.

From that, each generated task is either connected to an actor or an object, or to both. There are no interconnections between tasks within the boundary and no data on flows. The generated DataFlow Model can be visualised using μ ML notation, and interpreted by the user using the clear distinction between its types of flows, as well as the life (cycle) history of an object. One possible example is: a Staff inputs some data to the Add Module task, then the task creates an object of the type Module and stores it in the Module datastore. The following figure (Figure 9.5) demonstrates the initial DataFlow Model of the Module Management System.

Figure 9.5: (initial) DFD model. Manage Module System

At this stage of the development, a manual engagement, provided by business-users, is required for supplying the created *DataFlow Model* with data on flows. Then, the next translation step is fired by the user in order to complete the development process. The following figure (Figure 9.8) illustrates an initial DataFlow with data on flows, supplied.

Figure 9.6: DFD model. Manage Module Sys. (with data on flows)

From the generated *DataFlow Model*, it can be seen that each task is connected by two flows, which means each involves two atomic tasks. For instance, *Add Module* task receives the detail of a new module from the user and then (creates) inserts a new module instance into the collection of the modules in the system. Therefore, these tasks must be decomposed into their atomic tasks by translating the initial DataFlow into the detailed DataFlow Model.

Deriving Detailed DataFlow Model

A further transformation step is carried out in order to derive a complete DataFlow Model that contains atomic tasks. This transformation is known in BUILD as intitial DataFlow to detailed DataFlow Model, in which each task in the initial model is broken down into its atomic subtasks that are grouped into a logical boundary. The following figures: Figure 9.7, Figure 9.8, Figure 9.9 and Figure 9.10 demonstrate the generated logical boundaries that form a detailed DataFlow Model resulting from this transformation step.

Figure 9.7: DDFD model. Delete Module subtask.

Figure 9.7 above illustrates the internal data flow representation of the *Delete Module* task in the initial DFD model. The task is expressed, in the detailed DFD representation, as a logical boundary that consists of two atomic subtasks. The first one is a subtask to receive data from a user (Input), whereas the second task is a task to delete the corresponding module from the system. These subtasks are resulting from the decomposition step of the original Delete Module task expressed in the source DataFlow Model.

An extra Read flow between these generated tasks is added after analysing their priority scores. The source of the flow is the subtask that has a higher priority score. In this case, the Input Code task that is connected to an input flow is higher priority than the Delete Module task.

Figure 9.8 shows a logical detailed DFD Boundary that represents the initial DFD Add Module task at this level of abstraction. It consists of two decomposed atomic subtasks, a subtask to receive data from a user (Input) and another one to create a new module (insert it into the system).

Similar to the Delete Module task, presented before in Figure 9.7, an additional read flow between these subtasks is added to connect the subtask that has a higher priority score to the one that has a lower score. In this case, it is the *Input Code* task that is connected to the *Create* Module one.

Figure 9.8: DDFD model. Manage Add Module subtask.

Figure 9.9: DDFD model. Manage Modify Module subtask.

Figure 9.9 illustrates the internal data flow representation of the Modify Module task. The task is expressed, in the detailed DFD model, as a data flow boundary that contains three main subtasks. The first task is introduced for receiving data from a user (Input), the second and third task correspond to the broken down Update task, and the read and write subtask. The read task is used to retrieve a module that required to be updated, while the write one is used to update particular columns of the retrieved module.

There is a generated *read* flow between the decomposed *Update Module* subtasks (*read* and write). It is assumed that this additional flow connects the Read Module task to the Write Module task. On the other hand, another read flow is added that connects the Input Code task to the Read Module task based on their priority scores.

Figure 9.10) shows the data flow representation within the *See Description* task. A data flow boundary, with an equivalent name to the task name, is utilised at this level of abstraction to present the internal behaviour of each business task, in terms of data flow. It consists of a subtask to receive data from a user $(Input)$, a task to read a new module and insert it into the system, and a third subtask to display (output) the retrieved result to a user.

At the analysis phase, an independent mapping step is carried to derive the *State Model* from the previously generated detailed DFD model. The *Add Module DFD* boundary, for instance, is mapped directly to an *Add Module* boundary in the *State Model*, containing two states (*input* and create) that are derived from the DataFlow Input and Create flow respectively. Moreover, a new transition, with a particular action and/or condition(s), is added for each generated state

Figure 9.10: DDFD model. See Module Description subtask.

input_code_title_credit_desc_Waiting and Create_Module_Ready as a result of translating DFD Tasks.

In addition to this, a separate in-place modification is applied to the constructed Add Module boundary in order to generate an error handling (reporting) state for each State in the boundary. In this case, the *input_code_title_credit_desc_Waiting_Error* and *Create_Module_Ready_Error* are added to the *Add Module* boundary. Figure 9.11 below demonstrates the content of the *Add* Module boundary in the State Model.

Figure 9.11: State model. Add Module subtask.

Similar to translation steps and decisions mentioned previously for constructing the Add Module boundary, the content of the Delete Module and Modify Module are constructed likewise. The following figures (Figure 9.12 and Figure 9.13) illustrate the content of the Delete Module and Modify Module boundary respectively.

Figure 9.12: State model. Delete Module subtask.

At this stage, the *Data Dependency*, *DataFlow* and *State Model* are completely derived to represent a detailed picture of the system. All business tasks have been broken down into their atomic parts and relate to each other and other tasks in a particular control flow.

9.3.2.3 Information System Representation at the Design Phase

The intermediate artifacts that are generated from the analysis phase, namely, Data Dependency, DataFlow and State Model are used as source models in this stage of development. The Design phase in BUILD aims at producing two detailed models: Database and Query Model and GUI Model. The following subsections represent both models at the design level of abstraction.

While the *Module Management System* deals with a single business entity, there is no significant transformation decision, from *Data* to *Database and Query Model*, in this case study. The following figure (Figure 9.14) demonstrates the produced DBQ model at the Analysis Phase of BUILD. The Module entity expresses the structure of a relational database table in a platformindependent way rather than specialised for a particular target environment.

In addition, the Graphical User Interface (GUI) Model is the last detailed design model at this stage. Each state from the analysis State Model appears in the GUI Model as a window. For example, the *Input_regNumber_Waiting* state is translated into a window that has the same name. The generated GUI control elements (widgets) for each *Window* is based on the type of the State, either waiting, ready or error. The number of controls is determined by the number of local variables in each state. The following listing (Listing. 9.4) demonstrates a snapshot of the *GUI Model*, including the specifications of three windows of the models, namely, Input regNumber Waiting, Input regNumber Waiting Error and Read Code Ready.

Figure 9.13: State model. Modify Module subtask.

Figure 9.14: DBQ model. Module Entity

Listing 9.4: GUI Model

```
1 <gui:GuiBoundary id="1" name="Enrol">
2 <gui:Window id="2" name="Input_regNumber_Waiting" order="4">
3 <gui:Textfield id="3" name="regNumber"/>
4 <gui:Button id="4" name="Exception" event="Exception" exit="false"/>
5 <gui:Button id="5" name="Input" event="Input" exit="false"/>
6 \sim /gui:Window>
7 <gui:Window id="6" name="Input_regNumber_Waiting_Error" order="0" error="true">
8 <gui:Label id="7" name="input_regNumber_Waiting_Error_warning"
9 text="Null value not accepted"/>
10 <gui:Button id="8" name="initialise" event="initialise" exit="false"/>
11 </gui:Window>
12 <gui:Window id="9" name="Read Code_Ready" order="3">
13 <gui:Label id="10" name="regNumber" text="regNumber"/>
14 <gui:Label id="11" name="code" text="code"/>
15 <gui:Button id="12" name="Read" event="Read" exit="false"/>
16 <gui:Button id="13" name="Exception" event="Exception" exit="false"/>
17 </gui:Window>
18 </gui:GuiBoundary>
```
9.3.2.4 Executable Code from the Student Enrolment System

The framework of BUILD generates a number of Swing Java classes (*.java) files, and a single $MySQL$ dump file, $(*.sq)$ file. A full representation of code is presented in Appendix B. The following images demonstrate the running system screens under Eclipse, as well as the MySQL Workbench 6.1 compiling report after executing the generated $MySQL$ script file.

9.4 Overview of the Student Enrolment Sub-system

The Student Enrolment System enables the users to complete the enrolment process in university. The system has one main actor (*Student*). The actor is able to enrol in a module, using their registration number and the module code (for simplicity). The task is performed (executed) on a single enrolment at a time.

9.4.1 Information System Representation at the Requirements Sketching Phase

In this section the μ ML requirements models that express the enrolment system is presented graphically. The translation activities are then discussed with the highlighting of some key transformation decisions made by the transformation approach. The complete underlying XML representation of all models at this stage is listed in Appendix B.

The development process starts when a business-user expresses the structure of the system using the μ ML Task Model, as seen in Figure 9.15 below. The system does one main job, namely, enrolling a student in a module. This business process can be expressed as a single main Task called Enrol. The human-computer interaction is represented via placing a Student actor connected to the Enrol business task.

Figure 9.15: Task model. Enrol a Student sub-Sys.

Furthermore, the behaviour of the enrolment system is captured using the μ ML Impact Model. The model consists of a task that is equivalent to a business task captured in the *Enrol a* Student in the Task Model. The task, in this model, interacts internally with one object namely, Enrolment, in which the Enrol task is connected to the Enrolment entity by a create flow. The following figure (Figure 9.16) is the *Impact Model* of the *Enrol a Student* sub-system.

Figure 9.16: Impact model. Enrol a Student sub-Sys.

The following figure (Figure 9.17) is the *Information Model* of the *Enrol a Student*:

Figure 9.17: Information model. Enrolment Entity

9.4.2 Running the Experiment on the BUILD Framework (2)

Purpose: Generate a Java Swing information system with MySQL backend database. Critical Feature: Adopting appropriate decisions for handling multiple steps of inputing data to the system.

Input: Three μ ML models: Task, Impact and Information Model. Output: Java classes (*.java) and a MySQL script file (*.sql). Running Environment: Eclipse.

9.4.2.1 Construction of Requirements Models

Similar to the previous example, a memory model of the system Task Model is constructed first manually by business users. The Task class is used to initialise the $Enrol$ business task in memory, while the Actor and Paricipation are used to define the *Student* actor and its interactions to the system respectively. Listing 9.5 below demonstrates the Java code for the construction of the Task Model of the Student Enrolment System to be used in BUILD.

Listing 9.5: Construction of the Student Enrolment Task Model

```
1 Diagram taskModel = new Diagram();
 2 Boundary boundary = new Boundary ("Enrol a Student");
3
4 Actor actor1 = new Actor("Student");
5 \mid Task enrolStd = new Task ("Enrol");
6
7 Participation link4 = new Participation();
8 link4.addRole(new mde.task.model.Role("student",actor1));
9 link4.addRole(new mde.task.model.Role("enrol", enrolStd));
10 Participation link3 = new Participation();
11 link3.addRole(new mde.task.model.Role("student",actor1));
12 link3.addRole(new mde.task.model.Role("enrol", enrolStd));
```
In a similar way, the *Impact Model* of the system is created manually in memory, using node classes in the Impact Model Package. The ImpTask class is used to initialise the Enrol (Impact) task in memory, whereas the $ImpCreateFlow$ and $ImpObject$ are used to define the create impact and the *Enrolment* business object respectively. Listing 9.6 below demonstrates the construction of the Impact Model. It can be seen that names of boundaries and tasks are equivalent to previously designed boundaries and tasks in the Task Model, Listing 9.5 above.

```
Listing 9.6: Construction of the Student Enrolment Impact Model
```

```
1 ImpDiagram ImpactModel = new ImpDiagram();
2 ImpBoundary impboundary = new ImpBoundary("Enrol a Student");
3
4 ImpTask impEnrolStd = new mde.impact.model.ImpTask("Enrol");
5
6 ImpObject impObj3 = new ImpObject("Enrollment");
7
8 ImpCreateFlow cf = new ImpCreateFlow();
9 ImpRole impcf2 = new ImpRole("enrol", impEnrolStd);
10 ImpRole impcf1 = new ImpRole("enrollment", impObj3);
11 cf.addImpRole(impcf2);
12 cf.addImpRole(impcf1);
```
Furthermore, the Information Model is the final model of this stage. The following listing (Listing. 9.7) demonstrates a snapshot of the related part of the Student Enrolment System. It is designed using the Information Model Package of the BUILD framework. The class Entity is used to define the back-end entities of the system.

Listing 9.7: Construction of the Student Enrolment Information Model

```
1 mde.information.model.Diagram informationModel = new mde.information.model.Diagram();
2
3 Entity enrolEntity = new Entity("Enrollment");
4 Attribute attr10 = new Attribute("regNumber", new Type("Integer"))
5 . setIdentifier(true);
6 Attribute attr11 = new Attribute("code", new Type("String")).setIdentifier(true);
7 enrolEntity.addAttribute(attr10);
8 enrolEntity.addAttribute(attr11);
9 informationModel.addEntity(enrolEntity);
```
9.4.2.2 Information System Representation at the Analysis Phase

Similar to the experiment (1), the requirements models of the system (Task, Impact and Information Model) are produced manually at the requirement stage. A series of translations at this development layer aims at producing three analysis models, namely, the Data Dependency, DataFlow and State Models. The following sections present the creation of these models.

As this sub-system interacts with only a single entity, the constructed *Data Model* consists of one entity, Enrolment. There is no significant decision in this example. The following figure (Figure 9.18) demonstrates the structure of Enrolment entity at the Data Dependency Model.

Figure 9.18: Data model. Enrollment Entity

Besides this, as a result of merging Task and Impact Model, the initial DataFlow Model is constructed, containing a single business transaction involving a task, Enrol, to receive users inputs and for creating a new enrolment object. It can be noticed that there are two input flows from the *Student* actor to the *Enrol* task. This expresses the multiple steps of entering data into the system without any information about their order. The following figure (Figure 9.19) shows the representation of the initial DFD.

Figure 9.19: (initial) DFD model. Enrol a Student sub-Sys.

After the end-users engagement to supply the model with data on flows, the following figure (Figure 9.20) represents the initial DFD after adding appropriate data on flows.

Figure 9.20: (initial) DFD model. Enrol a Student sub-Sys. with data of flows

From the generated DataFlow Model (Figure 9.20), it can be seen that the *Enrol* task is connected by two *input* flows and one *create* flow. This requires a further decompositional step in order to generate a number of atomic tasks that perform the original Enrol one.

Deriving Detailed DataFlow Model

The intital DataFlow to detailed DataFlow Model transformation step is applied in order to derive a complete *DataFlow Model* that contains atomic tasks. The interesting transformation decision, in this case, is made to construct a logical boundary to represent the *Enrol* task that contains two Input tasks and one Create task. These three subtasks are connected to each other by a number of read flows.

The direction of flows, between subtasks, is determined by the type of flow that is connected to each task. For example, an extra read flow is generated from a task that is connected to the *input* flow (e.g. *Input Code*) An extra read flow is generated to connect *InputCode* to CreateEnrollment (the former handles the initial input flow and the latter handles the final create flow; so these must communicate). However, the transformation step that is responsible for generating flows between tasks has a default rule applied to the model when they have two or more input flows. Tasks are ordered in the DataFlow Model, based on the order in which the input flows were entered in the Task Model. In this case, the task Input RegNumber precedes

the task InputCode. The following figure (Figure 9.21) demonstrates the detailed DFD:

Figure 9.21: (detailed) DFD model. Manage Module Sys.

Similar to the previous case study (section 9.3), the *State Model* is derived from the detailed DFD model. The *Enrol* DFD boundary is mapped directly to an equivalent state boundary, called Enrol, in boundary within the target model. It contains three states (two are waiting states and one is ready) that are derived from the DataFlow Inputs and Create flow respectively.

Furthermore, a number of transitions, with appropriate actions and/or condition(s), are added to the model, reflecting the right sequence of navigation of system screens to perform the original task. Similar to a previous case study (section 9.3), a separate in-place modification is applied to the constructed *State Model* to generate error handling (reporting) states for each State in the boundary, see Figure 9.22 below.

Figure 9.22: State model. Manage Module Sys.

From the above transformations, the construction of all required analysis artifacts (Data, DataFlow and State Model) is accomplished. These models are ready to be utilised as source models in the next Design Phase of the framework.

9.4.3 Information System Representation at the Design Phase

The Design Phase in BUILD aims at producing the Database and Query and GUI Model from the intermediate analysis models, constructed above. The following subsections represent both models at the design level of abstraction.

While the system deals with a single business entity, there is no significant transformation decision, from *Data to Database and Query Model*, in this case study. The following figure (Figure 9.23) demonstrates the produced DBQ model at the *Analysis Phase* of BUILD. The *Enrolment* entity expresses the structure of a relational database table in a platform-independent way.

Enrollment
regNumber: Integer(10) code: Integer(10)
createEnrollment()

Figure 9.23: DBQ model. Enrollment Entity

In a similar way to the previous experiment, this model is constructed mainly from the analysis State Model. The following listing (Listing. 9.8) demonstrates a snapshot of the GUI Model, including the specifications of two windows, namely, Create_Enrollment_Ready and Create Enrollment Ready Error.

Listing 9.8: GUI Model

```
1 <gui:GuiBoundary id="1" name="Enrol">
2 <gui:Window id="16" name="Create Enrollment_Ready" order="2">
3 <gui:Label id="17" name="regNumber" text="regNumber"/>
4 <gui:Label id="18" name="code" text="code"/>
5 <gui:Button id="19" name="Create" event="Create" exit="false"/>
6 <gui:Button id="20" name="exit" event="exit" exit="true"/>
7 <gui:Button id="21" name="Exception" event="Exception" exit="false"/>
8 </gui:Window>
9 <gui:Window id="22" name="Create Enrollment_Ready_Error" order="0" error="true">
10 <qui:Label id="23" name="Create Enrollment Ready Error warning"
11 text="Connection to the Data source is fail"/>
12 <gui:Button id="24" name="exit" event="exit" exit="false"/>
13 </gui:Window>
14 ...
15 </gui:GuiBoundary>
```
9.4.3.1 Executable Code from the Student Enrolment System

The framework of BUILD generates a number of Swing Java classes, (*.java) files, and a single $MySQL$ dump file, $(*.$ sql) file. A full representation of code is presented in Appendix B. The following images demonstrate the running system screens under Eclipse, as well as the MySQL Workbench 6.1 compiling report after executing the generated $MySQL$ script file.

9.5 Overview of the Extended Information Model of the University Administration System

This section presents a fairly complex conceptual data model for the University Administration System case study adopted in this chapter. The model consists of a number of information entities and various types of associations, generalisations, and one kind of composition. In addition, a range constraint is applied to an attribute in order to illustrate how the transformation may react and makes alternative decisions compared to unrestricted attributes.

Two ways for constructing the Data Dependency and DBQ Models are followed and the results are compared. Starting from the μ ML *Information Model* is one possible way for starting the transformation chain till the Data and DBQ Model is achieved. The second way is by deriving automatically the initial Information Model from the Impact Model that is provided manually by business users. The following figure (Figure 9.24) illustrates the extended Information Model of the University Administration System.

Figure 9.24: Information model. University Administration Sys.

9.5.1 Running the Experiment on the BUILD Framework (3a)

Purpose: Generate a MySQL database system. **Input:** one μ ML model: Information Model. Output: a MySQL script file (*.sql). Running Environment: MySQL.

The Information Model of the University Administration System is used to demonstrate the capability of the BUILD framework to generate a complete executable M_ySQL script from a very abstracted information model. In order to achieve this, two steps of model-to-model translation are required, namely, Information to Data Dependency Model and Data Dependency to Database and Query (DBQ) Model. Moreover, a further code generation step is also applied to generate the final $MySQL$ code from the μ ML DBQ schema.

9.5.1.1 Information Model Representation at the Requirements Sketching Phase

Similar to previous experiments, the Information model is constructed manually in BUILD as an Abstract Syntax Tree (AST). A snapshot of the specification of the entities and attributes are expressed in the Listing 9.9 below (Person entity only), which represents the Java code used to create the model in memory (see Appendix A for the full code). The *Entity* and Attribute classes are used to construct the structure of the business entities and their properties, whereas the *Association*, *Generalisation* and *Composition* classes are used to define a variety of relationships between entities. The *range constraint* is specified using the $setLowerbound()$ and setUpperbound() methods of the Attribute object.

```
Listing 9.9: Construction of the Complex Information Model (Person Entity)
```

```
1 mde.information.model.Diagram informationModel = new
2 mde.information.model.Diagram();
 3
 4 Entity personEntity = new Entity("Person");
5 Attribute attr2 = new Attribute("name", new Type("String"));
 6 Attribute attr3 = new Attribute("age", new Type("String"))
7 . setLowerbound(18). setUpperbound(35);
8 Attribute attr4 = new Attribute("gender", new Type("String"));
9
10 personEntity.addAttribute(attr2);
11 | personEntity.addAttribute(attr3);
12 personEntity.addAttribute(attr4);
13
14 informationModel.addEntity(personEntity);
```
9.5.1.2 Information Model Representation at the Analysis Phase

At the analysis phase of BUILD, the *Data Model* is constructed (Figure 9.25). It includes some significant translation decisions made by the *Information to Data Model* translator. The manyto-many association between *Student* and *Module* is promoted to an entity *Enrolment* which depend on its related entities. In addition to this, the many-to-one association between *Student* and Address is resolved in the direction from the many side (Student) to the one side (Address) in the Date Model.

Furthermore, the Generalisation relationship between Person, Staff and Student is resolved by making Student and Staff (the specific entities) depends on the Person (the general entity). On the other hand, the *disjoint Generalisation* between *Assessment* and *Assessment Exam* is resolved by merging the general entity (Assessment) into each specific one (Assessment Exam and Assessment_Project). Moreover, the Total Composition relationship between Module and Assessment is resolved by making the part entity (Assessment) dependant on the whole one (Module).

9.5.1.3 Data Dependency Representation at the Design Phase

The evolution of the Data Dependency Model at the Design Phase in BUILD is the Database and Query (DBQ) Model, visualised in Figure 9.26 below. The generated DBQ model is considered a final detailed model that expresses a relational database schema in such a generic representation in a higher level of abstraction. All Dependency relationships between entities are translated into Foreign Keys (references).

There is a standard rule in the *Data Model to DBQ Schema* that translate any range constraint applied to a particular attribute into a Trigger associated to the table that holds that attribute. This is expressed in the following DBQ schema (Figure 9.26) in the third partition of the Person entity.

Figure 9.25: Data model. University Administration Sys.

9.5.1.4 Executable MySQL Code from the Database and Query (DBQ) Schema

Unlike previous experiments, the code in this example is only a single $(*)$ file. The current version of BUILD has a domain-specific code generation framework $(MySQL)$. The following code presents a snapshot of the final database schema in $MySQL$, representing the structure of the Person table and its associated Trigger. The complete schema is presented in Appendix C.

Listing 9.10: database MySQL.sql

```
1 CREATE DATABASE sysDatabase;
2 USE sysDatabase;
3
4 -- Structure for table 'Person'
5
6 CREATE TABLE Person (
7 identity INT(10) NOT NULL,
8 name VARCHAR(30).
9 age VARCHAR(30),
10 gender VARCHAR(30),
11 PRIMARY KEY(identity));
12
13 -- Trigger: Applying Checking Constraints on table 'Person'
14
15 DELIMITER //
16
17 DROP TRIGGER IF EXISTS 'personCheck' //
18 CREATE TRIGGER 'personCheck' BEFORE INSERT ON Person
19 FOR EACH ROW
20 IF (NEW.age < 18 OR NEW.age > 35) THEN
21 SET msg = 'INVALID DATA IN age';
22 SIGNAL SQLSTATE '45000' SET MESSAGE_TEXT = msg;
23 END IF; //
```


Figure 9.26: DBQ model. University Administration System

9.5.1.5 Running the Experiment on the BUILD Framework (3b)

Purpose: Generate a MySQL database system. **Input:** one μ ML model: Impact Model. Output: a MySQL script file (*.sql). Running Environment: MySQL.

A possible Impact Model of the University Administration System is designed to be used for deriving an initial *Information* and *Data Model* and then generating $MySQL$ code. In order to achieve this, a further model-to-model translating step is required, which is Impactto-Information Model.

9.5.1.6 Information Model Representation at the Requirements Sketching Phase

The Impact Model is provided manually, by users, in BUILD as an AST. The suggested model consists of a number of tasks that interact with system entities. Appendix C contains the complete Java code for creating the Impact Models using JAST[102].

Deriving the Initial Information Model

At the requirement gathering level, the initial *Information Model* might be derived from the pre-defined Impact Model. The model consists of a number of information Entities and their predicted relationships. This translation step is carried by an Impact-to-Information Model translation agent (rule). The rules are discussed previously in Chapter 7. Figure 9.27 below illustrates the generated initial Information Model.

Figure 9.27: The Initial Information Model

According to the Impact-to-Information Model step, each object in the Impact Model is used to generate a corresponding Entity in the Information Model, without any predicted details about its attributes. Business users might supply this information later to the generated model in an independent step, similar to the step of annotating the initial DFD model. Besides this, after tracing the CRUD effect for each task in the Impact Model, the relationships with appropriate multiplicities are predicted.

It is useful to compare the Information Model provided in experiment (3a) with the Information Model provided in experiment (3b) to examine the similarities and differences between them. All entities appearing in model (3a) are generated in the second one in (3b), except the *Enrolment* entity in which the many-to-many Association between *Student* and *Module* is captured as two many-to-one relationships from the Impact Model. This difference is in fact the result of normalising the *many-to-many* relationship, which would happen in the next step anyway. This shows that it is possible to derive normalised data models directly from the *Impact* Model.

Furthermore, a possible Generalisation relationship is predicted between the Student and Staff object, in the Impact Model. Consequently, a general entity is manufactured and a number of many-to-one associations between the specific entities (many sides) and the general one (one side) are generated directly between them. As before, the Total Composition relationship between the Model and its parts Exam and Project entity are also predicted as a number of many-to-one Associations, where the parts are connected on the many-side of these associations.

9.5.1.7 Information Model Representation at the Analysis Phase

At the *Analysis* stage, the typical translation rules are applied by the *Information-to-Data Model* agent to produce the *Data Dependency Model*, (Figure 9.28) below. To avoid the repetition, the generated DBQ Model and $MySQL$ script are described in Appendix C containing the representation of the (3b) DBQ Model and its generated $MUSQL$ database schema.

Figure 9.28: The derived Data Dependency Model

After compiling the generated $MySQL$ script file into the MySQL Workbench 6.1 server, the following result is obtained:

Using the reverse engineering facility in the MySQL Workbench 6.1 server, we can reconstruct the database schema of the generated and executed database system we run. From that, the following diagrams demonstrate the reverse engineered data model derived from the Impact model and the user-defined information model respectively.

CHAPTER 9. CASE STUDIES 9.5. OVERVIEW OF THE EXTENDED INFORMATION MODEL

Business User Made Information Model v.s Generated (initial) Information Model

When considering the user-defined *Information Model* (Figure 9.24) and the derived one, from the Impact Model (Figure 9.27), we can notice that the generated model contains all entities appearing in the manually constructed one. However, unlike the first *Information Model*, entities in the second model (derived) have no detail about their attributes. This occurs because there is insufficient information in the Impact Model to enable translation agents to translate or predict and then manufacture appropriate attributes for each entity. An additional entity appears in the generated Information Model, from the Impact one, as a result of capturing the many-to-many association via two impacts on the *Module* and *Student* object, in the *Impact Model*. Each one represents a single many-to-one association. In addition to this, the parts Exam and Project are not expressed in the Impact Model, see Appendix C.

Furthermore, *aggregation* relationships are predicted as *many-to-one* associations between the whole entity and its associated parts. Besides this, it can be seen that generalisation relationships are also captured via the existence of the Conjunction concept in the model and directly resolved into *many-to-one* associations between the entities involved.

After completing the whole development process that produces the final MySQL schema, the number of tables in the user-defined model is fewer than the number of tables in the derived model, due to the accuracy of resolving the *disjoint generalisation* relationship, in which the general and specific table will be merged.

9.6 Outlook on the Chapter

A real-world enterprise information system case study (University Administration System) was presented in this chapter. Three experiments were designed, each involving a particular part of the system. Requirements models, namely, Task, Impact and Information Models were expressed, manually, using the graphical notation of μ ML.

For each case study, the chapter represents the development stages of each model from the requirements sketching level to the code generation one. All significant transformation decisions are highlighted, showing how the approach generates an independent menu screen in the first experiment, whereas it manufactures a sequence of input forms (screen) for allowing a user to insert inputs step by step.

Throughout the demonstrated results, it can be seen that the designed chain of transformations is able to generate complete basic information systems that are connected to backend relational databases. The generated system is platform-specific and is tested within Eclipse. The produced information system is implemented in Java and the related database is generated in $MySQL$. The generated Java code includes a JDBC connection statement for establishing the required connection to the backend database. Full listings of these experiments are given in Appendices A, B and C.

10 Evaluation and Testing

"Discovering the unexpected is more important than confirming the known" George Box (1919- 2013)

10.1 Context

The main theme of this thesis is to introduce a simple and semantically cleaned modelling language (μML) , and a MDE framework, which is suitable for end-users, to ease and accelerate the capability of developing small/medium scale information systems. In order to achieve this ambition, the proposed language must be able to express, in a simple and clear way for business end-users, all critical concepts in the information systems domain. In addition to this, the transformation approach must be able to lead the development of information systems from requirements to the final executable code. As a result, μ ML and BUILD are introduced in the previous chapters of this thesis.

This chapter tries to emphasise the completeness and correctness of the output results from BUILD. To make sure μ ML is capable of modelling the most common enterprise system concepts we have chosen to evaluate it against the generated objects and chunks of code that appear in the final results. Model translation and code generation for producing JDBC Java Swing Applications with a MySQL backend database, in the BUILD Framework, are used to verify and inspect the evolving of models from the requirement phase to the final one that generates the executable code.

Furthermore, the *time efficiency* of transformation algorithms is evaluated to measure the quality and scalability of the proposed framework. The Big-O notation technique is adopted, in this chapter, for articulating how long an algorithm takes to run, and to examine whether or not there exist some factors or limits on inputs that affect the growth rate of time in some transformation steps. Apart from that, the suitability of the proposed μ ML is also evaluated by conducting an experiment with selected university students from different disciplines. It aims at evaluating μ ML in terms of its simplicity and ease to adopt by end-users as a modelling language for expressing basic specifications of information systems.

10.2 Evaluating The Generated Results from BUILD

10.2.1 Assuring the Determinism of the Transformation Rules

This criterion emphasises the capability of the framework to restrict transformations to the appropriate type of input element, and to produce one specific type of result. There are two types of transformation rules provided in BUILD, namely, one-to-one (forward) rule and two-toone (merge) mapping rule. When a particular concept appears in a model, it will be accepted only by one or more approved rules to produce a particular target concept(s). This means that the mapping is controlled by a precondition(s) that must be met before executing any translation rule, via the (accept()) method (see Chapter 7).

This strategy also ensures the right treatment of concepts (based on their types and semantics), in which it is impossible to misinterpret any concept chosen by a business user, and produce an unexpected target element. However, it leads to increasing the number of translation agents, as there is an agent for each concept to ensure that a correct decision is made.

For example, handling composed business tasks, in the Task Model, is different from the atomic ones, in which each task type is determined by checking whether it is attached to the head of any kind of *Composition* relationship or not. The composed task is normally ignored at the first translation step (constructing the initial DFD model), whereas atomic tasks are merged with the equivalent tasks that appear in the Impact Model, and converted into $DataFlow$ tasks, and then to State Model transitions. Table 10.1 shows translation rules that are applicable for each Task Model concept.

Table 10.1: Task Model concepts and Related Agents

In the *Impact Model*, for example, each type of *impact* is translated differently into its corresponding type of flows in the DataFlow Model. For example, Create Impact is converted into a DfCreateFlow that has more specific features (e.g. *assignment*) than the DfReadFlow, for instance, which has *filters*. Table 10.2 shows translation rules that are applicable for each Impact Model concept.

Impact Model Concept	Translator Agents
ImpDiagram	DiagramToDfDiagram
ImpBoundary	BoundaryToDfBoundary
ImpRole	ImpRoleToDfRole
ImpTask	MergeTaskToDfTask
Create Impact	CreateFlowToDfCreateFlow
Read Impact	ReadFlowToDfReadFlow
Update Impact	UpdateFlowToDfUpdateFlow
Delete Impact	DeleteFlowToDfDeleteFlow

Table 10.2: Impact Model concepts and Related Agents

When considering the translation of the Association concept in Information Model, for instance, it can be noticed that it might be treated in three different ways based on its specific meaning $(m-to-1, m-to-n, or 1-to-1)$. The $m-to-1$ Association is translated into a *Dependency* relationship in Data Model. Whereas m-to-n and 1-to-1 Associations are converted into Linker and Merged Entity respectively. Generalisation and Compostion are treated likewise. Table 10.3 shows translation rules that are applicable for each Information Model concept.

Information Model Concept	Translator Agents
InfDiagram	InfDiagramToDDiagram
Entity	InfEntityToDEntity
Association (m-to-1)	InfAssocToDDependency
Association $(1-to-1)$	InfAssocToMergedTable
Association (m-to-n)	InfAssocToLinkerTable
Generalisation (overlapping)	InfGenToDDependency
Generalisation (disjoint)	InfGenToMergedTable
Composition (total)	InfComToMergedTable
Composition	InfCompToDDependency

Table 10.3: Information Model concepts and Related Agents

These *Dependency* relationships are resolved in the next translation step into a number of *Foreign* Keys in the Database and Query (DBQ) Model. The generated FKs represents the relationship between the business entity at the Design phase for generating a relational database system. Therefore, the only meaning of Foreign keys is to represent a relationship between the entity that holds the key and the one it refers to. Table 10.4 shows translation rules that are applicable for each Data Dependency Model concept.

Data Model Concept	Translator Agents
DDiagram	DDiagramToSchema2
DEntity	DEntityToTable, DEntityToPrimaryKey
DAttribute	DAttributeToColumn
Dependency	DependToReferTable
DRole	DRoleToForeignKey

Table 10.4: Data Dependency Model concepts and Related Agents

Regarding the construction of Procedure concept in the DBQ Model, it can be seen how the importance of the type of corresponding Flows in the detailed DFD model. For example, any DFD create flow is translated into a DBQ representation of a *Procedure* that performs *insert* database operation. There is no other way to produce *insert (create)* DBQ procedure. Table 10.5 shows translation rules that are applicable for each *DataFlow Model* concept.

Table 10.5: DataFlow Model concepts and Related Agents

In addition to this, concepts in the *State Model* are treated in a direct way, in which the internal architecture of its relevant translation agent is simple. Based on the type of *State*, Ready or *Waiting*, the decision of firing either the *SvariableToLabel* or *SvariableToTextfield* is made to translate the variables of this state into suitable GUI controls. Table 10.6 shows translation rules that are applicable for each State Model concept.

Table 10.6: State Model concepts and Related Agents

In regard to the GUI Model that is used at the Code Generation phase, each type of widget element is translated into the corresponding Swing Java control for that element, by a particular code generator agent. For example, the specification of Java Labels are generated from GUI Label elements in the GUI Model by the JavaLabelGenerator, using approved specifications by the transformation approach (e.g. font type, size, colour, and text), and the rest is likewise.

Furthermore, as each generated system screen, there is a main button that performs that main task associated with that screen. The *Java* code of the button is generated via the *JavaBut*tonGenerator, in which the agent fills the body of the *actionPerformed* method of that screen by boilerplate code. This is approved by the generator agent, to establish a suitable JDBC connection, invoke the related stored procedure that performs a particular database operation, or presents user interactions by passing their inputs to the system. This strategy is applicable for all types of screens that might be generated using BUILD.

Additionally, the **boilerplate code** for error handling is also considered as a part of the actionPerformed method, using the try-catch expression and a number of appropriate Exceptions to avoid unexpected behaviour of the generated system. Table 10.7 shows generation rules that are applicable for each GUI Model concept.

GUI Model Concept	Generator Agents
GuiBoundary	JavaCodeFileGenerator
Window	JavaClassGenerator
Textfield	JavaTextFieldGenerator
Label	JavaLabelGenerator
Button	JavaButtonGenerator

Table 10.7: GUI Model concepts and Related Agents

Besides generating JDBC Swing Java code, the Code Generation framework also generates suitable $MySQL$ scripts from the DBQ model in BUILD. Each concept appearing in the DBQ model has a particular generator agent that is responsible for generating a specific portion of the $MySQL$ Data Definition Language (DDL) or Data Manipulation Language (pre-defined queries). Table 10.8 below illustrates this.

Stored Procedures in MySQL can be generated from either a DBQ Create, Update, Delete or Query concept. The difference between the procedure derived from the DBQ create and the one generated from the DBQ Query for instance, is in the $MySQL$ statement (command). Therefore, stored procedures that bind the $MySQL$ INSERT statement are generated only from the DBQ procedures that have a create concept, in contrast to the ones that bind MySQL SELECT query that are derived from the DBQ procedures that have a *Query* element. The rest are likewise.
DBQ Model Concept	Generator Agents
Schema	MySQLDumpFileGenerator, MySQLSchemaGenerator
Table	MySQLTableGenerator
Column	MySQLFieldGenerator
PrimaryKey	MySQLPKGenerator
ForeignKey	MySQLFKGenerator
Procedure	MySQLStoredProcGenerator
Range constraint	MySQLTriggerGenerator
Filter	MySQLQueryGenerator, MySQLUpdateGenerator & MySQLDeleteGenerator
Project	JavaButtonGenerator
Query	MySQLQueryGenerator
Create	MySQLCreateGenerator
Update	MySQLUpdateGenerator
Delete	MySQLDeleteGenerator

Table 10.8: DBQ Model concepts and Related Agents

10.2.2 Was it Possible to trace the Development of Models during the Development Stages?

This criterion demonstrates the capability to trace each intermediate artifact and its evolution after completing every translation and/or transformation step. One of the main claims of the μ ML is its suitability for business users, in which it is designed to enable them, with their limited technical knowledge, to express their system functionalities in a very generic and abstracted way. These semantically cleaned μ ML models are serialised into XML files, during each development stage in BUILD, resulting in a number of platform-independent views of the system. This This supports a validation strategy by Model Inspection (or Code Inspection) to identify defects in each produced artifact. Any unexpected defect might be determined by tracing the evolution of each requirement-level concept throughout the development stages, and comparing it with what concept we expect to see at the final code.

The first column in tables 10.9 and 10.10 demonstrate each visualised concept that appears in μ ML requirement models, before commencing any translating step. The rest of the columns show the evolution of each concept throughout the development stages with regard to what concept we expect to see at each level of detail till reaching the final code. The Code Generation Framework in BUILD produces two types of domain-specific code, namely, JDBC/Swing Java code and MySQL scripts. Table 10.9 considers the main requirements model concepts and traces its development till generating the final *JDBC Swing Java* code. On the other hand, Table 10.10 considers the development of the concepts till generating the executable database schema script in MySQL.

For instance, the *create* impact, in the *Impact Model* is developed to be a *Ready State*, at the end of the *Analysis* phase. This state, next, is represented as a system *Window*, at the *Design* stage. Then, each developed screen is converted into an equivalent Swing Java class, (Table 10.9). At the same time, the *create* impact is also developed to be a *create* flow, in the *DataFlow* Model supplied with some data on it, at the end of the Analysis stage. This flow is converted into a DBQ Procedure concept and then into a MySQL Stored Procedure at the Design and final $MySQL$ script respectively (Table 10.10).

In addition to this, it is worth mentioning that for each item of data on each flow in the DataFlow, we expect to see a state variable equivalent to this datum. At the Design phase, this variable is converted into an argument in such a DBQ Procedure and in a local GUI control widget of the system screen. At the end, it becomes an *argument* in a *constructor* of a related Swing Java class, as well as an argument in a relevant MySQL Stored Procedure.

Table 10.9: Evolution of Requirement concepts to Java

Furthermore, there is a step to ensure that each DBQ Table has a Primary Key in the DBQ model. This step is taken by the *DEntityToTable* translation rule (table 10.9). If there is no identifier for such a DEntity, there will be a manufactured Primary Key attached to the target DBQ Table.

Evolution of Range constraints is another example that shows an interesting transformation decision. It can be evaluated by detecting the generated MySQL schema file looking for an equivalent *Trigger* associated with the table that holds the *ranged* attributes. This decision is relied on by the domain-specific code generator, in this case MySQL Code Generator because it is known that range constraints might be handled differently by another relational data vendor, such as *Oracle*. In the *Oracle* version of SQL , this constraint is handled by the CHECK command as part of a column declaration. Therefore, the evaluation of this concept might be distinct based on the target environment of each code generator.

Initial Concept	Analysis Phase	Design Phase	MySQL Script	
Business Actor	Actor (DFD), Entity (Data)	Table (DBQ)	MySQL Table	
Impact Object	Object (DFD), Entity (Data)	Table (DBQ)	MySQL Table	
Identifier	Identifier (Data)	Primary Key (DBQ)	MySQL Primary Key	
Attribute	Attribute (Data)	Attribute(DBQ)	MySQL Column	
Range constraint	range constraint (Data)	range constraint (DBQ)	MySQL Trigger BEFORE INSERT	
Upper and Lower Bounds	Upper and Lower Bounds (Data)	Upper and Lower Bounds (DBQ)	MySQL Trigger	
Association (1-to- 1), $(m-to-n)$	Entity (Data)	Table (DBQ)	MySQL Table	
Association $(m-$ to-1)	Dependency (Data)	Foreign Key (DBQ)	MySQL Foreign Key	
Generalisation, Composition	Dependency or Entity (Data)	Foreign Key or Table (DBQ)	MySQL Foreign Key or Table	
Impact Create	Create flow (DFD)	Create Procedure (DBQ)	MySQL Stored Procedure	
Impact Read	Read flow (DFD)	Query Procedure (DBQ)	MySQL SP, MySQL Stored Procedure	
Impact Update	Update flow (DFD), Read/Write flows (DDFD)	Update Procedure (DBQ)	MySQL SP, MySQL Stored Procedure	
Impact Delete	Delete flow (DFD)	Delete Procedure (DBQ)	MySQL Stored Procedure	
- -	Datum on flow (DFD)	Argument of Procedure (DBQ)	Argument MySQL Stored Procedure	

Table 10.10: Evolution of Requirement concepts to MySQL

10.2.3 Was the Generated Code Complete with Respect to Initial Inputs at the Requirement Sketching Phase?

This criterion examines some quality dimensions of the models employed within BUILD during the whole development. This includes the completeness, consistency and clarity of the content of each model. This helps to know whether the transformation approach takes correct mapping decisions and produces complete outputs with an acceptable level of quality. All models used within BUILD must satisfy the proposed evaluation criteria in this section. It is represented, in table 10.11, as a number of goals that must be met.

According to the criteria provided, table 10.11, it can be seen that both the State and GUI Model have not satisfied the goal: Name of elements look professional. This occurs because of the strategy we follow to form names of *States*, in which some names might be long. This leads to having long names for both the generated Windows, and Java classes. A proper naming convention has to be considered at this level of transformation.

Goal	\rm{Task}	Impact	Information	DFDs	State	ら ひ	Data		Code Model
Each model is represented via a well-formed XML document.	✓	✓	✓	\checkmark					
Every element has an appropriate namespace.	✓	✓	✓	✓	✓	✓	\checkmark		
Every element has a unique <i>name</i> and <i>id</i> .	✓	✓	✓	✓	\checkmark	✓			
Inter-refering relationships appear between elements belong to the same model.	✓	√	\checkmark	✓	\checkmark	√	\checkmark		
Elements that are instances of a model belong to the associated metamodel of that it or the <i>core</i> μ ML.	✓					✓			
The structure of the model is logically correct.	✓	✓	✓	✓	✓	✓	✓	✓	
Each <i>descendant</i> of <i>Arc</i> has a source and a target.	✓	✓	✓	✓	✓	✓	✓	✓	
Each typed element has a recognised data type.	✓	✓	✓	✓	✓	✓	✓	✓	
Each Named element has a meaningful name.		✓	✓	\checkmark	✓	✓	✓	✓	
Name of elements look professional.		✓	✓	✓	X	Х	✓	✓	
Appropriate Naming convention is applied to each elements (e.g. camelCase).	✓	\checkmark	✓	\checkmark		✓			
No dublicated concepts in models	✓	\checkmark	✓	√		√			
No null elements in models									

Table 10.11: μ ML Evaluation Criteria

10.2.4 Could We Generate All That We Want?

In this section, a broader picture of the generated enterprise systems is considered, as well as some of its detailed component specifications. System windows, navigation, type of widgets, business logics and database schema are examples of main components that must be produced completely and correctly from the transformation approach. Any good (valid) system must have these components to be able to perform real-world business processes.

The BUILD framework is able to produce consistent screens with business tasks described at the Requirement Phase. This means that for each atomic task, there will be a screen (Swing Java class), that extends the *JFrame class* to perform this task and a screen to handle (report) its failure scenario. For instance, according to the *Student Enrolment System* case study presented previously in Chapter 9, three system screens are generated in correspondence to the subtasks of the system, as well as another three error reporting ones. Table 10.12 below describes this.

Atomic Task	System Screen
Input Code	Input_code_Waiting, Input_code_Waiting_Error
Input reg_Number	Input_regNumber_Waiting, Input_regNumber_Waiting_Error
Create Enrollment	Create_Enrollment_Ready, Create_Enrollment_Ready_Error

Table 10.12: Atomic Business Tasks and their Corresponding Windows

According to the introduced priority scoring strategy, which is based on the impact of each task on system entities and the participation of actors on each task, the order of executing these subtasks is determined to form the internal sequence of the main business process. As a consequence, the navigation between system windows is established, reflecting the order of subtasks execution, calculated by the transformation approach. Table 10.13 and 10.14 below demonstrate the generated sequences in both presented case studies.

Table 10.13: Order of Windows (Student Enrolment)

Table 10.14: Order of Windows (Module Management - Modify Module)

Furthermore, the framework is also able to generate a system with a number of business tasks that might be executed optionally in an independent way. It generates an appropriate menu screen to enable the end-user to select which process is required to be performed. For instance, in the Module Management System case study, chapter 9, Manage_Module_Main_Menu_Waiting is a root menu for the system that allows the selection between a number of atomic business tasks, such as Add Module, Modify Module and more.

The current version of BUILD is able to generate limited navigations (links between screens). This limitation is discussed in detail in the following section (10.8). The following table (10.15) summarises the number of completely generated links between windows in both case studies presented previously in chapter 9.

In addition to this, establishing a proper *JDBC* connection between the *presentation layer* and the *data layer* is critical. This part of the generated code is regarded **boilerplate**, which is present in every Java class that requires preparing $MySQL$ database connectivity statements. Besides this, the part of the code for catching internal errors, such as passing invalid (null) inputs to the system, is also considered **boilerplate**, in which all screens involved in the successful business scenario must have this exception handling chunk of code to report any error to the user.

According to the structure of system windows, this version of the framework is able to generate four types of Swing Java GUI controls (widgets), namely, JTextField, JButton and JLabel, all placed within the JPanel container. The JButton and JLabel controls are common for each type of screen, in which each window has a label to display its title and a button to perform its specific action. However, the JTextField widget appears only when passing external user inputs to the system.

Control properties, such as *size, text* and *data type* are extracted from both the DBQ and GUI Model. From that, it is possible to generate very basic GUI controls that enable a simple business process to be completed. It is worth mentioning that additional controls might be generated from the DBQ Model, but it is not supported yet in this version of BUILD. See next section (10.8) for further details.

With regard to the generated relational database connected to the system, it is essential to use a normalised database in order to provide an information system that works at an acceptable level of performance (Table 10.16). By using BUILD, it can be assured that the final M_ySQL database schema satisfies the requirements of the third-normal form. This degree of normalisation is achieved according to the translation chain started from the *Information Model*, which contains unnormalised relationships (e.g. generalisations, compositions, and many-to-many associations), and ending up with the normalised Database and Query Model, which contains resolved relationships.

Relationship	Engaged Entities	Rule Applied	Final Table(s)
Association $(m-to-1)$	Student, Address	Generate FKs	Student, Address
Association $(m-to-n)$	Student, Module	Generate Linker table and FKs	Student, Module, Enrollment
Generalisation (disjoint)	Assessment. Assessment_Exam	Merge	Assessment Exam
Generalisation (overlapping)	Person, Student, Staff	Generate FKs	Person, Student, Staff
Composition $_{\rm (total)}$	Module, Assessment	Generate FKs	Module, Assessment_Exam

Table 10.16: Data Model Normalisation

10.2.5 Are Transformations Adequate to Fill the Gap in Implementation?

On some occasions, during the chain of model transformations, new concepts are inferred due to applying particular mapping decisions. For instance, a unique *Identity* in the *Data Dependency* Model or an automatic *Primary Key* in the Database and Query Model is manufactured according to the translation agent decision when such an entity in the source model has no information about its *Identity*. Furthermore, the size of some attributes, that have no *size* detail given in the source model, is also manufactured based on the data type.

In addition to this, generating a menu screen that holds alternative functions of the system is another example of the design decision, taken by the transformation/code generation rules, to fill the implementation gap. The Module Management case study is an example of this. Moreover, the Code Generation approach is able to fill the implementation gap by filling the generated classes by boilerplate code for arranging the database connectivity and embedding commands that are used for executing and calling SQL stored procedures.

10.2.6 Can We Execute the Generated Code?

This criterion aims at examining the generated code/database from BUILD after running them under the related environments. This helps to answer some code evaluation questions to know whether the generated code has compiled, executed and done its desired job successfully without mistakes. As the current *Code Generation Framework* of BUILD produces only *Java* code and $MUSQL$ script, a decision was made to compile the generated Java files using the Eclipse Framework, and execute the final $MySQL$ dump file using $MySQL$ Server 6.1.

The following Table (10.17) demonstrates the criteria for evaluating the generated system and its related database. It can be seen that the generated Java code has not satisfied the goal: The generated system is able to receive multiple selected record from the database. This occurs because the code generators in BUILD are currently designed as a proof of concept, in a way that supports only passing back a single record, or part of a record. This issue might be overcome by considering retrieving multiple records in a later version of the code generation framework.

Indeed, this goal is related to another one that the generated $MySQL$ schema has not satisfied, which is: Query stored procedures are able to pass multiple selected records to the system. This happens because the stored procedures are designed, as a simple proof of concept, in a way that they bind a single SQL statement that is stored in a single-value out parameter to be passed back to the code. This issue might be also overcome by considering retrieving multiple records in a later version of the code generation framework.

Table 10.17: Criteria for Evaluating the Generated Information System

10.2.7 Things Were Wrong or Missing and Suggested Repairs

Expanding the generated type of widgets, to cover more GUI controls, is a main issue that has to be considered in later versions of BUILD, as the type of widget is based on the type of screen, whether it is for inputting or displaying information. The current translation rules provide two types of GUI control to handle this, namely, *JLabel* and *JTextField*. It is possible to improve the current version of BUILD by filling the implementation gap with more accurate widgets than the current one. For example, a *JComboBox* control might be used to restrict the input data to a particular field. This can be done by allowing only one (or many) selected value(s) to be valid inputs, from selection constraints on some attributes in the Data or DBQ Model.

In addition to this, adding the ability to display a collection of results (items) in such a JTable control is essential. The current version of the BUILD Code Generator is able to generate executable code for retrieving and displaying a single result only. In the case that the query procedure returns a collection of records, the code will present the first record only to be the retrieved result. In order to solve this issue, we must expand our μ ML notation at the requirement level to be able to express a collection of objects.

Furthermore, in the case of generating a system with multiple input steps, the current version of BUILD produces a window for each step. These windows form a chain to perform the whole input process. Adding an extra design choice to allow merging these windows into a compound one, which holds a longer *entry form* to receive all user inputs at once, offers an optimal system with a fewer number of screens. This can be achieved by enhancing the $(in-place)$ transformation agent that optimises the generated *State Model* from the detailed *DataFlow* one.

Moreover, there is an issue (limitation) in the developed algorithm for translating the detailed DFD into *State Model*, in which the current version is allowed to translate a short sequence of atomic tasks that must start (be activated) by a user input only. We are aware that some business tasks might be fired automatically, before any user interaction, when the system starts (pre-processing).

The pattern for extracting this from the detailed Data Flow Model is not yet supported. In order to overcome this dilemma, the algorithm must be improved to be capable to cover all missing patterns.

Besides this, there are a variety of business patterns (control flows) that might be covered in later versions of the proposed μ ML and related BUILD translator agents. To solve this problem without expanding the graphical constructs of the modeling language (keeping μ ML lightweight), we suggest adding the multiplicity annotation to the Task Model, particularly to the Composite relationship and improve the extraction (prediction) of its possible interpretations. Then, the relevant translator agent generates a possible control-flow construct, similar to the way we follow for dealing with the optionality of subtasks (see the related case study in chapter 9).

10.3 Analysis of Time Complexity

Analysing the *time efficiency* of transformation algorithms is considered to measure the quality and scalability of the proposed framework, including its architecture and implementation. This section aims at evaluating the performance of the model transformation approach within BUILD. The Big-O notation technique is adopted, here, for articulating how long an algorithm takes to run, and to examine whether or not there exist some factors or limits on inputs that affect the growth rate of time in some transformation steps. This section discusses the *time efficiency* of the framework from three perspectives: rule complexity, transformation step complexity and overall complexity. The section also describes how the complexity in such a translation step may be reduced by caching the results of repeated transformations in the Context.

10.3.1 Big-O Analysis of Rules

As previously presented in the thesis, the current version of BUILD consists of two types of transformation rule, namely, one-to-one translation and two-to-one merging (folding) rules. The following subsections discuss the complexity of these types. The time complexity varies from one rule to another even within a transformation step. The worst case scenario is considered for each rule, and later for each step (Section 10.3.2). This decision assures that any possible alternative scenario has less complexity than the worst case one.

For all transformation rule classes in the framework, the main computation of each rule is performed by the doTranslate method that belongs to that rule. This method runs once for each distinct execution of the rule on distinct arguments. Each rule is constructed once inside the dominant rule that calls it. The graph of all applicable rules is constructed prior to rule execution, so does not contribute to the time complexity of execution. Construction time is negligible, the linear in the number of rules constructed.

The body of the *doTranslate* method in all rules contains one or more conditions for checking any precondition of that rule, implemented using if or if. else statements. This chunk of code (construct) is executed once at each run of the doTranslate method. From that, the time complexity of this construct is $O(1)$. On some occasions, calling other methods, e.g. translate, is required, the time complexity of this call or any further method invocation is also $O(1)$.

In non-terminal rules, when translating further properties of a target element, or invoking some subrules for translating its descendants, a kind of iteration is required to complete the whole translation step. The time complexity of this part of the code is $O(N)$, where N is a number of properties that are ready for translation. Even if the loop has some if or if. else statements, object creation and method invocations, the worst case is still $O(N)$ based on the number of properties/descendants of the target element.

Complexity	Number of one-to-one Rules	Number of two-to-one Rules
O(1)	31	$\overline{2}$
O(N)	17	
$\mathbf{O}(N*M)$	θ	$\overline{2}$
$\mathbf{O}(N * M^2)$	1	θ
$\mathbf{O}(N * M * K)$	3	5
$\mathbf{O}(N * M * K * P * Q)$	1	θ
	0	3

Table 10.18: Different Complexity Types in the Transformation Approach

When nested *iterations* occur within *doTranslate* method, the computation time is increased to be $O(N*M)$, where N represents the number of target elements and M represents the number of a particular type of properties/descendants for each target element. For each type of rule in the framework, Table 1 demonstrates all possible worst case complexities and the number of rules that represents each one.

10.3.1.1 One-to-One Rule Complexity Analysis

The *one-to-one* rule complexity analysis is applicable for analysing *time efficiency* for all concrete rules that extend Translation class in the top-level framework. This includes one-to-one translation, transformation and in-place modification. It is worth emphasising that the time complexity of most terminal rules is $O(1)$, as the body of their *doTranslate* methods consists of two main parts: checking preconditions $(O(1))$ and performing the mapping $(O(1))$.

However, there are a few terminal rules that have a loop to handle a collection of properties within that same rule. The time complexity of these rules is $O(N)$, where N is number of properties/descendants of the source element. According to the implementation, one or more foreach constructs are used in some rules to perform this job. The majority of rules, in the framework, have at most one loop. As seen in Table 1, the time complexity of those rules can be either $O(1)$ or $O(N)$.

Apart from this, the usual complexity of a rule reduces to $O(1)$ if the same rule is invoked again on the same argument(s). This will happen if multiple higher-level rules depend on the same lower-level rule, which is executed multiple times. The time saving is achieved by caching the result of every rule execution in a Context, so that it need not be recomputed. However, the overall reduction in complexity depends on the degree of overlap and depth of nesting in how the rules are designed, so it cannot be systematically quantified in general.

10.3.1.2 Two-to-One Rule Complexity Analysis

The two-to-one rule complexity analysis is applicable for analysing time efficiency for all concrete rules that extend the *MergeTranslation* class in the top-level framework. The majority of rules have some *iterations*. These loops are mostly used to traverse elements in the second source model for checking or collecting further information required for generating a target element.

The complexity of these rules depends on how deeply the iteration is nested. It is found in the designed rules that the number of nested loops are either three or two. The time complexity of these constructs are either $O(N * M)$, $O(N * M * K)$ or $O(N^2)$. It is worth mentioning that a few rules have at most one loop. The time complexity of these rules can be either $O(1)$ or $O(N)$ (see Table 10.18.).

Figure 1(a) summarises the remarkable changes in growth rates of time complexity in four possible worst case scenarios: $O(N)$, $O(N*M)$, $O(N*M*K)$, $O(N*M^2)$, $O(N*M*K*P*Q)$ and $O(N^2)$. The x axis represents the number of N, while the y axis represents the computation time. It is assumed that all calculations are based on a single run. It can be seen how the scalability becomes poorer in the cases that have a larger number of nested and linear iterations.

10.3.2 Big-O Analysis of Translation Step

As previously presented in the thesis, the current version of BUILD consists of a number of model transformation steps, implemented independently in a number of *Java Packages*. These packages contain several rules (Java classes) to perform such a complete mapping step. This part of time efficiency analysis focuses on the complexity of the transformation step itself, by considering the Big-O notation of the worst case scenario of all included rules in each step. For each transformation, the total number of rules that have similar complexity are grouped together and are represented in Table 10.19.

Table 10.19 summarises the total number of each type of rule involved in each translation step. In the step of translating the detailed *DataFlow Model* into the *DBQ Model*, for instance, there are four rules that have $O(N * M * K)$ time complexity. For each flow in the DFD, N is the number of flows, M is the number of data variables on that flow and K is the number of attributes in the target end-role of that flow. The complexity of each translation step is calculated as the worst-case complexity of any rule group involved in that step. From that, the time complexity of Information to Data is $O(N)$, Data to DBQ is $O(N * M^2)$, Task & Impact to DBQ is $O(N^2)$, DFD to DDFD is $O(N * M * K * P * Q)$, DFD to DBQ is $O(N * M * K)$, DFD to State is $O(N^2)$, Task to State is $O(N)$, State to GUI and Impact to Information is $O(N * M * K)$ time.

10.3.3 The Overall Time Complexity Analysis of BUILD

This sections aims at calculating the exact time performance of BUILD, rather than the big-O estimate. As the BUILD transformation strategy is designed in a linear way, in which each step runs only once. Transformation rules within each step are designed in a way that a rule calls another rule one or more than one time (e.g. to translate an entity and its corresponding attributes). The time complexity for each transformation step is calculated in the previous section (Section 10.3.2) and summerised in Table 10.19. The highest time complexity of each step is considered the worst case (Big-O notation) of that step.

In order to calculate the overall time complexity, we need to multiply the time complexity by the number of transformation rules that has this complexity. For instance, according to Table 10.19, the time complexity of the DFD to DBQ step is $O(N * M * K)$, in which there are four rules that have this complexity. From that, the exact time performance of this transformation step is $N * M * K * 4$. The overall complexities of the rest of steps are calculated likewise.

10.4 Evaluation Experiment of μ ML Notation

In order to evaluate the suitability of μ ML, an experiment has been conducted with selected university students from different faculties and departments such as, Electrical and Communication Engineering, Information Management, Chemistry and Medicine. The main objective of this experiment is to evaluate μ ML in terms of its simplicity and ease of adoption by endusers as a modelling language for expressing basic functional requirements and specifications of a real-world and commonly used information systems. It is worth noting that none of the participants were Computer Scientists, who might have been exposed to UML notations. This was important, in order to get the responses of non-specialist business users.

Students are interviewed and the general overall idea about the task and the modelling language is explained to them before commencing the task. After that, an incomplete solved

business case study is given to them, showing what exactly is expected from them. Then, students were given a task to draw several μ ML diagrams of given business scenarios, taken from the same case study. The time taken to draw each diagram was measured manually by the students.

10.4.1 Design of the Experiment

The experiment involves an *Online Hospital Booking System (OHBS)* case study, exemplifying how a registered patient books an appointment to see a registered doctor in the hospital. A Task, Impact and Information Model to this part of the system is given to students showing how the notation of each model looks like. Then, students will be asked three questions to draw Task, Impact and Information Models for some business activities within OHBS and record the time they needed for each question (see Appendix D). In order to analyze the results of the experiment, criteria are designed, including the following points:

- The average time for drawing diagrams is less than 5 mins, between 5 mins and 8 mins or greater than 8 mins.
- Business entities is captured with correct names and notation.
- Business activities (tasks) are captured with correct names and notation.
- Business structure of task are drawn with correct notation.
- HCIs between stakeholders and business activities are captured with correct notation.
- Impacts between entities and business activities are drawn with correct notation and direction.

10.4.2 Analysis of the Results

As mentioned before, three questions are provided to draw three types of μ ML requirement models, namely, Task, Impact and Information Model. The following subsections analyse the results for each task provided.

10.4.2.1 Results Analysis for the Task Model

The expected answer has 9 elements: 6 tasks, 2 compositions and 1 actor. Figure 10.1 below demonstrates a sample of the correct solution, sketched by a student. The assumed time for completing this question is three and a half minutes.

It was observed, from student answers, that all students were able to identify the stakeholder of the system. They also used the correct notation (ellipse) for expressing business activities. 60% of the students selected the correct (*composition*) type for representing the right whole/part relationship between business tasks, whereas 20% had not completed the model and 20% failed in using the accurate notation for representing a sequence of subtasks, which is one (whole) task and some (parts) subtasks that are connected to the *whole* one by a *Composition* relationship, but instead they used a chain of *Composition* relationships from the *whole* task to the last part subtask. This gives a specific order of subtasks' execution that is not a part of the givn scenario. The average duration time of all students to complete this task was 4.6 minutes.

Figure 10.1: Student attempt at drawing a Task Model

Table 10.20 illustrates (summarised) all answers provided by the participants for the first question.

Item	Std 1	Std 2	Std 3	Std 4	Std 5
Number of Element					
Drawing Time (mins)	5 ₅				

Table 10.20: Summary of All Student Answers (Question 1)

10.4.2.2 Results Analysis for the Information Model

The expected answer has 3 elements only: 2 objects (entities) and 1 association. Figure 10.2 illustrates an example of a complete solution, drawn by a student. The assumed time for the completion of this question is just one minute.

It was observed, from the provided answers, that all students were able to identify the correct business objects (entities) of the system. 100% of the students succeeded in drawing the correct notation (rectangle) for expressing business entities. Furthermore, all of them were able to extract the relationship between these entities. It is worth mentioning that determining multiplicities is out of the scope of this task.

However, 20% of the students sketched duplicate entities (e.g. New Doctor and Existing Doctor), and added extra entities out of the scope of the given case study. The average duration time of all students to complete this task was 3.8 minutes.

Figure 10.2: Student attempt at drawing an Information Model

Table 10.21 demonstrates all answers provided by the participants for the first question.

Table 10.21: Summary of All Student Answers (Question 2)

10.4.2.3 Results Analysis for the Impact Model

The assumed answer has 12 elements: 4 tasks and 4 objects and 4 impacts. Figure 10.3 below exemplifies a partial solution, answered by a student. The expected time for completing this question is 6 minutes.

One of the main observations was that 80% of the students were able to distinguish the notation used for capturing tasks from the one that is used for capturing *entities* and draw it correctly without any confusion. Only one sample (20% of the whole number of participants) used a wrong notation (ellipse) for drawing entities and (rectangle) for drawing tasks and did not complete the model.

In addition to this, 60% of the participants used correct impacts for representing internal interactions between tasks and entities, but 40% of them have at most either one wrong impact direction or wrong *impact* type.

Figure 10.3: Student attempt at drawing an Impact Model

Table 10.22 summarised all answers provided by the participants for the first question.

Item	Std 1	Std 2	Std 3	Std 4	Std 5
Number of Element					
Drawing Time (mins)	-5				

Table 10.22: Summary of All Student Answers (Question 3)

10.4.2.4 Overall Conclusion

After analysing the experiment results, it can be noticed that about 67% of the drawn models are complete and answers are correct. Non-specialist students were able to differentiate one concept from another. Reasons for failure could include that they did not learn the notation well enough; or they failed to understand the business scenario. Better-trained users might improve this statistic; and domain-experts might understand how to express their business more effectively. This is because they are domain experts who clearly understand all business activities. From that, our strategy in constructing small models that have clean semantics and simple notation is worthwhile to be adopted by non professional developers.

10.5 Outlook on the Chapter

The overall work presented in this chapter can be divided into three main evaluation dimensions, namely, evaluating the completeness and correctness of the rules, time efficiency analysis and suitability of μ ML concepts and notations. In the first dimension, criteria are introduced for evaluating whether or not the generated code is complete with respect to initial specifications expressed at the Requirements Sketching phase. Furthermore, the ability to execute the generated system (code) and its backend database without errors is also evaluated. It shows how the transformations approach fills the gap in implementation by taking critical design decisions and introducing appropriate chunks of boilerplate code.

In addition to this, the Big-O notation technique is utilised for evaluating the *time efficiency* of the model transformation algorithms in BUILD. For each type of transformation, the worst case complexity is calculated and explained within the chapter. At the end an experiment is conducted for evaluating the suitability of μ ML in terms of its simplicity and ease of adoption by end-users as a modelling language. The majority of results shows positive feedback from participants.

11 Conclusion and Future Work

"If you don't work on important problems, it's not that you'll do important work." Richard Hamming

11.1 Context

This chapter summarises the overall contributions conducted by the research into the Model-Driven Engineering for enterprise information system development with respect to the following aspects:

(1) Business users engagement contributions related to how the thesis expands the end-users role in the development lifecycle by enabling them to act as designers and lead the development processes, using less technical knowledge to construct their desired system.

(2) Modelling language contributions related to achieving a lightweight language that is able to capture end-users logical thinking about their information system using less technical knowledge, with a higher level of abstraction than other existing UML and other DSL approaches.

(3) Development approach contributions relate to the proposed MDE method that adopts a forward model transformation strategy and employs the proposed lightweight modelling language, to produce an executable information system derived from its initial requirements.

11.2 In Support of the Thesis

By observing the current state of art in Model-Driven approaches for web applications and information systems development in Chapter 2, we noticed, on the one hand, that most of the approaches target developers who have the ability to construct rich detailed models with technical specifications, such as UML-based approaches. On the other hand, in DSL-based approaches, such as $WebML[124]$, that aim at raising the abstraction level further, we noticed that developers must fill those models with a lot of detail in order to generate a complete code at

the end. For instance, the components of each screen must be specified, as well as the navigation links between pages.

From the above, it can be concluded that both approaches rely on modelling technical concepts at a particular level of abstraction. Based on this given knowledge, the transformation rules directly map the concepts, which appear in source models, to those equivalent concepts in the target. This section highlights the contributions of this thesis that improve the development process of information systems from the perspective of business end-users.

With respect to the current analysis of the state of art, Chapter 3 introduced a novel Model-Driven Engineering approach (BUILD) that includes a more intelligent forward model transformation strategy than current approaches. It allows rules of transformation to predict and provide new knowledge at each translation step carried on within the proposed method. Unlike the existing approaches, mentioned in the survey, the BUILD method starts with simple and semantically cleaned models (μML) ; the transformation mechanism then evolves these models and introduces new and more detailed concepts.

Besides this, the literature review revealed how little attention is paid to having an appropriate modelling language that supports End-User MDE for information systems development. To address this, a lightweight modelling language (μML) was introduced to enable the end-user to specify desired systems using their conceptual and business knowledge about the domain. The proposed language achieved an acceptable degree of clarity from the perspective of end-users.

The same notation was employed in different μ ML models, giving different meanings in respect to the context of that model. Chapter 4, 5 and 6 discuss in depth the notation and semantics of each model with regard to the development stage, in BUILD, which it belongs to. A formal flavour, using First Order Predicate Logic FOPL with some extentions, mentioned previously in Chapter 3, has been added to the definition of each concept.

Chapter 7, reflected the three-phase linear composition of model translations, within BUILD, that shifts the requirement models from the Requirements Sketching phase to the Analysis phase, and from the intermediate analysis models, used in the *Analysis* phase to the *Design* phase, and lastly from design models, used in the *Design* phase, to the code generation stage. This illustrates the advantages of the intermediate layers, in which new knowledge, supplied by business users or by the transformation rules, are allowed at each layer. As a result, some of the final outcomes are more heavily influenced by useful, generated intermediate concepts, than by the initial user input, illustrating the usefulness of a multilayered approach.

In addition, Chapter 7 also discussed the efficiency of the proposed overall architecture of the transformation approach in BUILD. It demonstrates how the structure successfully orchestrates the hybrid transformation rules that are implemented using Java, and which are styled as a number of Java translation agents (class). In the main rule for each translation step, As a result, control over the execution order is achieved, by natural interpretation of the rule dependency structure. Rules are also idempotent, in that a duplicate application maps to the same result.

Chapter 9 discussed, in depth, how the evolution of models is achieved. A full list of the detailed transformation rules was described for each translation step in the framework, showing how each rule (translation agent) produces a target concept from the source, infers new knowledge that is inserted in the translated elements, and makes a critical design decision to the structure or the behaviour of the system. It is worth saying that the designed rules are formalised using FOPL in terms (Forall (\forall) , Exists (\exists)) to ensure that all rules are deterministic with a known type of target.

At the end a number of translation, transformation, and in-place modification steps are assembled together, forming a complete chain of model transformations that start from the requirements and end with the final code.

11.3 Discussion of Research Questions

Question 1 What is the required interaction mechanism between the end-user and the transformation approach to capture their desired system specifications?

The interaction is achieved via the manual creation (sketching) of the requirements models using the proposed μ ML notation that is able to capture functional requirements, using endusers logical thinking and less technical knowledge, via three system views (Chapter 3 and 4).

Question 2 What kind of abstract system views do we need to capture in order to have comprehensive knowledge and behaviour of a system?

Domain experts have conceptual knowledge about their business from three perspectives. The structure of business processes and how system actors might interact with each process, the structure of business entities and how these entities might relate to each other, and finally, it is required to express how to supply/retrieve external data into/from the system. In other words, the interaction between system actor(s) and each process must be considered. Furthermore, the types of interactions (effects) between business processes and entities must be modelled (Chapter 3 and 4).

Question 3 What transformation rules are required to fold and optimise high-level views and introduce extra detailed design information into a system?

Merging the rule of concepts that appear in the Task Model with the equivalent ones that appear in the *Impact Model* is a significant step to producing the core DataFlow Model. It is designed in BUILD via the Task-Impact-to-DataFlow agent. Furthermore, inferring an initial Information Model from a pre-defined Impact Model is a mapping step that predicts new details regarding the relationships between business entities in the *Information Model* (Chapter 7 and 8).

Question 4 What transformation rules are required to refine high-level models and introduce richer design information into a system?

At the *Analysis* phase, the step of deriving the detailed *DataFlow Model* from the initial one is one refinement step. Task decomposition approaches are applied to tasks appearing in the initial DFD to produce atomic tasks that perform a single CRUD operation. Furthermore, the in-place model modification step on the generated State Model to generate appropriate error states for handling possible failure cases of business processes is another example of this kind of transformation.

In addition to this, the translation step from *Information Model* to the normalised Data Dependency Model is a step that performs data normalisation on the source model. It might add a manufactured *identity* for each entity that has no user-defined *identifier*. Not only this, This step also derives the detailed *Database and Query Model* by merging concepts from both the Data and DataFlow Model, which consists of a complete database schema that holds business

logic as well as table definitions.

Besides this, a direct one-to-one mapping between the *Information Model* and the intermediate Data Dependency Model is another translation step that analyses the types of relationships (e.g. generalisations and compositions), as well as multiplicities on associations to produce the dependency relationships between *Data Model* entities (Chapter 7 and 8).

Question 5 At which level of development is end-user engagement required to supply new knowledge to be considered by model translators at the next translation step?

Users manual engagement is performed by annotating the generated (initial) DataFlow Model with data on each flow. This occurs during the *Analysis* phase as a critical step in order to generate the detailed DataFlow Model (Chapter 3 and 4).

Question 6 To what extent are end-users able to generate a complete system for their demand?

As a proof-of-concept, the current version of BUILD is able to generate executable information systems that are *JDBC Java Swing* applications running in an Eclipse environment, that connect via $JDBC$ to a $MySQL$ backend database (Chapter 9 and 10).

Question 7 With respect to code generation, to what extent are we able to construct and link the IS layers from generic and less technical specifications?

As the current version of BUILD successfully produces two-tier information systems, its database connection is prepared and established as a *boilarplate* code generated by the relevant code generator within the framework. This strategy links the generated classes (screens) at the presentation layer to the back-end database. As a result, a complete connection is achieved.

11.4 Summary of Findings

11.4.1 Semi-Automated Way For Producing a Rich Detail Data Dependency Model

According to the BUILD method, the *Data Dependency Model* is only derived from the predesigned *Information Model*. The *Information Model* can be constructed manually by business users who sketch all model contents from scratch, including entities, attributes and relationships. The other way is by a particular translation agent that derives automatically an initial Information Model from scratch, using the knowledge provided previously in the Impact Model to predict entities and resolved relationships only.

An interesting exploration is concluded from the similarity between both generated Data Dependency Models resulting from experiment 3a and 3b, Chapter 9. The default transformation rules within the Impact-to-Information Model translator successfully produced a number of many-to-one associations that lead to derive identical dependencies between entities in the generated Data Model, from the user-defined Information Model. Under few circumstances, the translator agent cannot do more due to the lack of knowledge in expressing whether the predicted generalisation is *overlapping* or *disjoint*. Therefore, a default rule that ignores this issue

treats any possible generalisation as an overlapping one.

This finding raises a clear idea for introducing an alternative approach for constructing a rich detailed data model that combines benefits from both development ways mentioned above. It aims to achieve new generated Data Dependency Models from an Information Model that contains both the manually specified entities and their attributes, and automatically inferred and resolved relationships between entities, from an Impact Model.

This approach is beneficial because it narrows and defines the manual engagement carried out by end-users for constructing Information Models. This significant idea reduces possible manual errors, performed by business users, in designing relationships between entities. This current work can be continued to provide extra support to information system development approaches, led by end-users. This is described later in the future work section 11.5.

11.4.2 Alternative Strategy For Modelling Rich Information Models at the Requirement Phase

After the successful generation of $MySQL$ set constraints and the $PRIMARY$ KEY from requirement level concepts, we discovered that handling constraints from that level of abstraction is possible in a straightforward way. It does not require vast end-users technical knowledge, and substantial computation by the related translators. Unlike the manufactured and generated constraints, such as FOREIGN KEY, UNIQUE, and NOT NULL that requires checking other intermediate elements in the source and/or target models, during the creation of the Database and Query Model, these trivial constraints are directly passed from the source to the target till the development approach generates an equivalent script at the code generation stage.

Given the above, if emphasis is laid more on supplying trivial attribute constraints in an initial Information Model, this will lead to producing trivial attribute constraints at this level, which enriches the *Information Model*, alongside predicting associations automatically, with less end-user engagement, leads to producing a more sophisticated data model. The detailed description of entities and attributes, including structure and constraints may be extracted from the Information Model, whereas the knowledge about relationships between business entities may be predicted from the *Impact Model*. This strategy avoids the repetition of concepts or information in both the Impact and Information Model.

11.5 Future Work

It can be said that the introduced Micro-Modelling Language (μML) and the proposed Model-Driven Engineering approach (BUILD) improve the information systems development process from the perspective of business end-users. But there are a number of ways in which the current work can be extended.

11.5.1 Designing Additional Types of Generic Transformation Rules At the Top-Level Architecture

The current version of BUILD has two types of model transformation rules at its top-level framework. Other model transformations such as, in-place model modification and one-to-two forward translation must be added to the family of generic translation rules designed in BUILD. Currently, the in-place modification is treated as a one-to-one model transformation, in which the source and target have the same type, whereas the one-to-two translation is implemented via two one-to-one transformation rules.

11.5.2 Developing a Merging Translation Agent for Data Dependency Model Construction

This piece of future work is related to the findings discussed in section 11.4; it aims at enhancing the approach for developing the *Data Model* by enabling the BUILD method to avoid possible accidental end-user mistakes during designing relationships between entities or during annotating (assigning) their multiplicities. It is assumed that defining the structure of entities is considered a straightforward task for domain experts. In contrast, constructing accurate types of relationships, with correct multiplicities might be error-prone from the perspective of na¨ıve business users, which requires a further treatment.

This issue can be tackled by minimising the usage, and possibly notation, of the Information Model to cover only business entities and their structure (attributes and some constraints on them), and letting the relationships between these entities be predicted from the *Impact Model*. A slightly alternative approach might be adopted here, by including all partial relationships, about which the designer (end-user)has total confidence, in the initial information model.

In both suggested solutions, the knowledge in an (initial) Information Model is accumulated together with the predicted one from the Impact Model in an independent merging step. From that, designing the Information-Impact-To-Data Model translation agent is a sensible solution to be in charge of carrying out this translation step.

11.5.3 Enrich the Information Model by Supporting More Types of Constraints

In data modelling, various types of constraints are applied to fields in order to enforce particular input, restrict values, prevent empty values, and more. Some of these constraints are naturally known by domain experts without the demand to have additional technical data modelling skills. From their prespective, it is assumed that not null, unique, set constraint, range constraint, and default value are examples of trivial constraints that might be determined by business users at the requirement level.

Set constraint, for instance, is a type of constraint that applies to values of attributes within an entity/table. It occurs when acceptable values of a particular field are restricted by a limited known set of values. This constraint is expressed in MySQL using an Enumerator, appearing as an enum field within the generated table, which rejects any attempt to input any invalid data. As we successfully generated an equivalent $MySQL$ code for our early work in ReMoDeL[101] DBQ model and Database Generation Framework, introduced in [108] and [109].

As we have not paid adequate attention to these concepts during the design of the *Infor*mation Model, adding extra concepts that are able to express these constraints on attributes is, therefore, worth considering. In real-world business, this information is regarded as trivial from the perspective of domain experts (end-users). The evolution strategy of these concepts must be determined to enable passing these specifications throughout development stages of BUILD.

11.5.4 Covering More Business Workflow Patterns

The current version of BUILD is able to extract knowledge from various μ ML models to construct two types of business workflow patterns, namely, Sequence and Exclusive Choice. These supported patterns were demonstrated previously throughout Chapter 9 case studies. Covering additional workflow patterns is an interesting future research question raised by this thesis. It includes expanding the concepts and notation of some μ ML models, and designing extra translation agents to deal with these new concepts.

One possible suggestion is expanding the Task Model concepts to allow the expression of multiplicities in *Composition* relationships between tasks. This indicates the number of executions for each subtask involved in a *Composition*. As a result of this, an iteration control flow pattern might be detected from this concept. Besides this, adding conditions (filters) on the branches of a Composition allows for the possible detection of sophisticated conditional branching of business tasks. For both suggestions, an appropriate translation chain has to be defined in order to get the benefits of these patterns later in the design and code generation phases.

11.5.5 Optimising the DataFlow-to-State-Model Translation Algorithm

As mentioned in the limitation of the work in Chapter 10, section 10.8, the current version of the DataFlow-to-State-Model agent only supports the translation of a particular business process pattern that is expressed in the detailed DFD model into the associated State Model. It is desirable to enhance the current algorithm in order to construct more complex business tasks (such as a task that consists of a collection of input and output subtasks along with one or more subtasks that performs a single CRUD operation).

To achieve different pattern coverage, a further consideration might be required in the proposed priority scoring algorithm, in order to enable having more detailed scores to cover any possible repetition of task type. For example, in the current version, the score of read is 3 and the score of input is 4. On some occasions, when having a system that requires more than one input step, the *scoring algorithm* gives both (input) steps 4, and a default rule is applied to consider the first input task appearing in the model to be the first one in the sequence of input and the rest is likewise. An additional step for analysing the data on flow together with the flow-type is needed to extract complex (longer) sequences.

11.5.6 Constructing a Separated Business Logic Layer

The current version of BUILD is able to construct only a 2-tier information systems design, which consists of a *presentation* and *data storage (database)* tier. The specifications of these tiers are expressed via a number of μ ML models. In the generated systems from BUILD, the business logic is embedded in the *data storage* tier, and invoked from UIs, the *presentation* tier.

The logic is implemented via a number of stored procedures, each of which performs a specific CRUD operation on the database. This would not be applicable for computational business tasks that perform different jobs rather than CRUD operations.

From that, on some occasions, having a separate business logic tier brings benefits to the design, in which it enables dealing with some types of *computational* business tasks. As a consequence, the proposed method (BUILD) might be used on a large scale, dealing with systems that require a further security and performance of the overall information system design.

Constructing a 3-tier architecture might be achieved by modelling business logic in a separate model apart from the DBQ model that represents pure OO code. The μ ML Code Model serves this purpose. Chapter 6 introduces the notation and semantics of some of its core concepts and the rules-structure and definition of the translation step for deriving the Code Model from the DataFlow and DBQ model.

In order to continue the work in this direction, the notation of the Task Model must be extended to support new types of business activities, e.g. *arithmetic* in addition to the existing types of tasks. Choosing the right level of abstraction is important; for example, high-level tasks representing predefined accounting operations may be preferable to low-level arithmetic. Moreover, appropriate transformation rules must be designed to derive the equivalent code from related models.

11.5.7 Integrating BUILD with Eclipse Graphical Modeling Framework (GMF)

In order to improve the interoperability of the BUILD framework, designing a graphical editor tool is an essential improvement of the work. Eclipse, via the Graphical Modeling Project[37] offered a set of generative components and runtime infrastructures for creating graphical editors based on the Eclipse EMF[36] artefact and the GEF[35] framework. Using the technology provided by GEF[35] on the top of our approach, to invent rich graphical editors, would allow end-users to sketch their system models in a convenient way. In the literature survey (Chapter 2), many model transformation approached and development methodologies support their work with an Eclipse-based CASE tool, such as, ATL[33], Epsilon[34] and WebML[124] via its WebRation[4] CASE tool.

11.6 Final Remarks

In summary, the thesis introduced a user-friendly modelling language, called Micro-Modelling Language (μML) , aimed at supporting end-user model-driven engineering by raising the level of abstraction rather than using other UML approaches, which requires more technical knowledge. The main purpose of μ ML is to tackle some issues regarding the ambiguity of UML semantics and the complexity of its models. This is because the Models in UML are too complex and eclectic to be given a single, clear interpretation, while paradoxically not covering all of the views that are needed to completely specify a software system.

Designing this lightweight language (μML) that has a simpler notation with cleaner semantics enabled capturing the functional requirements of systems in a generic way, using familiar concepts to a business domain. This was demonstrated via examples and case studies throughout the thesis. The individual μ ML model is smaller and more restricted; but more types of

models were used together to cover the different interlinking aspects (views) of an information system.

In addition, possible total and partial forward transformations, between different kinds of model, were identified. The rules of transformation are designed in a hybrid style, gathering benefits from the known declarative and imperative approaches. Each rule is implemented as a Java class (agent), forming a layered forward model transformation strategy to generate an executable 2-tier information system from requirements models.

From that, a novel model-driven engineering approach (BUILD) was established. The BUILD method introduced a development methodology that contributes in the area of modeldriven information systems engineering. The method introduces a development technique that allows domain experts to lead the development process, and participate more at various development stages, by constructing initial MDE artifacts, and annotating some of the intermediate models at a later phase.

The engagement of business users is achieved via the proposed μ ML. Through BUILD, the development lifecycle starts with models, which hold business knowledge, defined manually by end-users. More views of the desired system are constructed automatically during the development process, in which the transformation rules evolve the captured requirements to predict, and produce, more detailed concepts, at the design and code generation phases.

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Appendices
A Appendix: Models and Executable Code of Experiment (1)

A.1 Complete Models and Full results of Experiment (1)

This Appendix presents the XML representations of all μ ML models of the Module Management System case study, and the completed Java and MySQL code generated from these models via BUILD. The experiment presented previously in chapter 10.

Experiment (1): Requirement Models Construction

Listing A.1: CaseStudy.java

38 import mde task model Composition;
39 import mde task model Diagram; 40 import mde.task.model.**Participation;
41 import** mde.task.model.**Task;**
42 import mde.taskimpact2dataflow.rule**.DiagramToDfDiagram;
43 import org.jast.ast.ASTWriter;** $\begin{array}{c} 44 \\ 45 \end{array}$ public class CaseStudy { $\frac{46}{47}$ 47 public static void main (String [] args) throws IOException,
48 TreeException, mde.gui.gen.TreeException {
49 // TODO Auto−generated method stub $\begin{array}{c} 50 \\ 51 \end{array}$ 51 // C o n s t r u c t t h e Task Mode l 52 Diagram taskMod e l = new Diagram () ; 53 Boundary boundary = new Boundary("Module Management") ; $\begin{array}{c} 54 \\ 55 \\ 56 \end{array}$ 55 Task manage = new Task("Manage Module");
56 Task add = new Task("Add Module"); $\begin{tabular}{ll} 57 & \textbf{Task} delete = new \textbf{Task}("Delete Module"); \\ 58 & \textbf{Task} modify = new \textbf{Task}(" Modify Module"); \\ 59 & \textbf{Task} see = new \textbf{Task}("See Description"); \\ 60 & \textbf{Actor actor 1 = new \textbf{Actor}("Staff"); \\ 61 & \textbf{Actor actor 2 = new \textbf{Actor}("Student"); \\ \end{tabular}$ 62
 63 63

64
 $\begin{array}{ll} \textbf{Composition} & \textbf{comp} = \textbf{new} \textbf{ Composition } () \, ; \\ \textbf{comp. addRole (new mde. task. model.} \textbf{Role (" manage", manage}) \, ; \\ \textbf{comp. addRole (new mde. task. model.} \textbf{Role (} "add", add)) \, ; \\ \textbf{comp. addRole (new mde. task. model.} \textbf{Role (} "delete", delete)) \, ; \\ \textbf{66} & \textbf{comp. addRole (new mde. task. model.} \textbf{Role (} "notify", modify)) \, ; \\ \end{array}$ 68 $\begin{array}{c|c} 69 & \textbf{Partition link1 = new Partition();} \hline 70 & \textbf{link1.addRole(new med. task model.Role("std")} \end{array}$ 70 link1.addRole(**new** mde.task.model.**Role**("staff",actor1));
71 link1.addRole(**new** mde.task.model.**Role**("add", add)); Participation link2 = new Participation ();

173 link2.addRole (new mde.task.model.Role ("staff", actor1));

174 link2.addRole (new mde.task.model.Role ("modify", modify));

175 Participation link3 = new Participation (); 84 85

86 boundary . addActor (actor 2) ;

87 boundary . addTask (manage) ;

88 boundary . addTask (add) ;

89 boundary . addTask (delete) ; 90 boundary . addTask (modify) ;
91 boundary . addTask (see) ; $\frac{92}{93}$ $\begin{bmatrix} 93 \\ 94 \end{bmatrix}$ boundary . add Participation (link1);
add Participation (link2); 94 boundary . add Participation (link2);
95 boundary . add Participation (link3); 96 boundary . add Participation (link4);
97 boundary . add Participation (link5); 98 boundary . addParticipation (link
98 boundary . addParticipation (link
98 boundary . addComposition (comp); $\frac{99}{100}$ taskModel.addBoundary (boundary); $\frac{101}{102}$ $\begin{array}{c|c}\n 102 \\
 103\n \end{array}$ ASTWriter writer = new

ASTWriter (new File 103 ASTWriter(new File("Manage_Module_taskModel.xml"));
104 writer.usePackage("mde.task.model", "xmlns:task");
105 writer.writeDocument(taskModel); 106 writer.close(); $\frac{107}{108}$ 108 System.out.println $("1)$ Task Model is Created by user.");
 $\frac{109}{ }$ 109 // −− // $\begin{array}{c} 110 \\ 111 \end{array}$ $\begin{tabular}{ll} 111 & // Construct the Impact Model \\ InpDiagram Inpatched = new ImpDiagram(); \\ InpBounding inpbound = new ImpBoundary("Module Management"); \\ InpTask impManager = new ImpTask("Manager Modules"); \\ InpTask, impAdd = new mde import model. ImpTask("Add Module"); \end{tabular}$ 116 ImpTask impD e l et e = new mde . impact . model . ImpTask(" D e l e t e Module ") ; 117 ImpTask impModify = new mde . impact . model . ImpTask("Mod ify Module ") ; 118 ImpTask impSee = new mde . impact . model . ImpTask(" See D e s c r i p t i o n ") ; 119 120 ImpObject impObil = new ImpObject ("Module"): $\frac{121}{122}$ 122 **ImpCreateFlow** cf = **new ImpCreateFlow** ();
123 **ImpRole** impcf1 = **new ImpRole**(" add", impAdd);
124 **ImpRole** impcf2 = **new ImpRole**(" add", impAdd); 125 cf. addImpRole(impcf2);
126 cf. addImpRole(impcf1); $cf. addImpRole (impc f1)$ $\frac{127}{128}$ 128 **ImpDeleteFlow df = new ImpDeleteFlow** ();
129 **ImpRole** impdf1 = **new ImpRole**("module", impObj1);
ImpRole impdf2 = **new ImpRole**("delete", impDelete); 131 df. addImpRole (impdf2);

```
132 df.addImpRole(impdf1);<br>
133 ImpUpdateFlow uf = new ImpUpdateFlow();<br>
134 ImpRole impuf1 = new ImpRole("module", impObj1);<br>
135 ImpRole impuf2 = new ImpRole("modify", impModify);
136 uf.addImpRole(impuf2);<br>137 uf.addImpRole(impuf1);
138 ImpReadFlow rf = new ImpReadFlow ();<br>139 ImpRole imprf1 = new ImpRole("nodule", impObj1);<br>140 ImpRole imprf2 = new ImpRole("see", impSee);
141 rf. addImpRole (imprf1);<br>142 rf. addImpRole (imprf2):
142 r f . addImpRole ( imp r f 2 ) ;
143 impboundary . addImpTask ( impManage ) ;
144 impboundary . addImpTask ( impAdd ) ;
145 impboundary . addImpTask ( impD e l et e ) ;
146 impboundary . addImpTask ( impModify ) ;
147 impboundary .addImpTask(impSee);<br>148 impboundary .addImpObject(impObj1);<br>150 impboundary .addImpCreateFlow(cf);<br>150 impboundary .addImpDeleteFlow(df);
151 impboundary . addImpUpdateFlow ( u f ) ;
152 impboundary . addImpReadFlow ( r f ) ;
153
                            154 ImpactModel . addImpBoundary ( impboundary ) ;
\frac{155}{156}156 ASTWriter writer 1 = new<br>157 ASTWriter (new File)
157 ASTWriter(new File ("Manage_Module_impactModel .xml"));<br>158 writer1 . usePackage ("mde.impact . model" , "xmlns:imp");
159 | writer1.writeDocument (ImpactModel);<br>160 | writer1.close ();
\frac{161}{162}162 System . out . p r i n t l n ( " ( 2 ) Impact Model i s Cr eat ed by u s e r . " ) ;
163 // −−−−−−−−−−−−−−−−−−−−−−−−−−−−−−−−−−−−−−−−− //
164<br>165
165 // Construct the Information Model<br>166 mde.information.model.Diagram informationModel = new
167 mde. information. model. Diagram ();
168<br>169
169 Entity modu l eEnt ity = new Entity ( "Module " ) ;
170 Attribute a t t r 1 2 = new Attribute ( " cod e " , new Type( " I n t e g e r " ) )
171 . s e t I d e n t i f i e r ( true ) ;
172 Attribute a t t r 1 3 = new Attribute ( " t i t l e " , new Type( " S t r i n g " ) ) ;
173 Attribute a t t r 1 4 = new Attribute ( " c r e d i t " , new Type( " I n t e g e r " ) ) ;
174 Attribute a t t r 1 5 = new Attribute ( " d e s c " , new Type( " S t r i n g " ) ) ;
175
176 module Entity . add Attribute (attr12);<br>177 module Entity . add Attribute (attr13);
177 modu l eEnt ity . a d dA t t r i b u t e ( a t t r 1 3 ) ;
178 modu l eEnt ity . a d dA t t r i b u t e ( a t t r 1 4 ) ;
179 modu l eEnt ity . a d dA t t r i b u t e ( a t t r 1 5 ) ;
180
181 informationModel.addEntity (moduleEntity):
\frac{182}{183}183 informationModel.addAssociation (new Association ()<br>184 . addRole (new Role ("module", informationModel
184 .addRole(new Role("module", informationModel<br>185 .getEntity("Module")).setMultiple(true))<br>186 .addRole(new Role("student", informationModel<br>187 .getEntity("Student")).setMultiple(true)));
188
189
<br>
ASTWriter writer2 = new<br>
ASTWriter(new File("Manage_Module_informationModel.xml"));<br>
writer2.usePackage("mde.information.model", "xmlns:inf");<br>
writer2.close();<br>
193
<br>
writer2.close();
\frac{194}{195}195 System . out . println ("(3) Information Model is Created by user.");<br>
\frac{196}{ }196 // −−−−−−−−−−−−−−−−−−−−−−−−−−−−−−−−−−−−−−−−−−−−−− //
\frac{197}{198}198 // Generate\ DataFlow\ Model<br>199 DiagramToDfDiagram topRule = new DiagramToDfDiagram ();
200 DfDiagram dataflowModel = topRule.translate(taskModel, ImpactModel);
\frac{201}{202}202<br>
ASTWriter writer3 = new<br>
ASTWriter (new File ("Manage_Module_DataFlowModel.xml"));<br>
writer3.usePackage ("mde.dataflow.model", "xmlns:dfd");<br>
writer3.close ();<br>
writer3.close ();
\frac{207}{208}208 System.out.println("(4) DataFlow Model is Created"+<br>209 "(Task + Impact −−> DataFlow).");
210 // −−−−−−−−−−−−−−−−−−−−−−−−−−−−−−−−−−−−−−−−− //
\frac{211}{212}1/|A| dding some data on flows in dataflow model.
213
\begin{array}{c|c}\n 214 \\
 \hline\n 215\n \end{array} // input to add<br>
\begin{array}{c}\n 214 \\
 \hline\n 215\n \end{array}214 data flow Model . get DfB ound aries ( ) . get (0 ) . get DfInputFlows ( )<br>215 data flow Model . get (0 ) . addDataonflow ("code" );
\begin{tabular}{ll} 216 & . get (0). addDataonflow("code"); \\ 217 & dataflowModel. get 01. get 010. get 00. get 01. get 011. get 011. get (0). addDataonflow("title"); \\ 219 & dataflowModel. get 01. addDataonflow("credit"); \\ 220 & . get (0). addDataonflow("credit"); \\ 221 & dataflowModel. get 01. addDataonflow("credit"); \\ 222 & . get (0). addDataonflow("desc"); \\ 223 & . get (0). addDataonflow("desc"); \\ 224 & dataflowModel. get 01. addDataonflow("code"); \\ 225 & . get (1). addDataonflow("code"); \\
```
226 dataflowModel.getDfBoundaries ().get (0).getDfInputFlows ()
227 .get (1).addDataonflow ("credit"); 228 // input to delete
239 dataflowModel.getDfBoundaries ().get (0).getDfInputFlows ()
231 // input to see
231 // input to see 232 dataflowModel.getDfBoundaries().get(0).getDfInputFlows()
233 .get(3).addDataonflow("code"); $\begin{tabular}{ll} 234 & // \textit{create assignment} & \\ \textit{dataflowModel} & \textit{get0IDBoundaries()}), get (0), getDfCreateFlows() \\ 236 & .get (0).addAssignment("@code = code"); \\ 237 & dataflowModel. getDfBoundaries() , get (0). getDfCreateFlows() \\ & . get (0).addAssignment("@title = title"); \\ 238 & dataflowModel. getDfBoundaries() , get (0). getDfCreateFlows() \\ & . get (0).addAssignment("@circle = credit"); \\ 240 & . get (0).addAssignment("@circle = credit"); \\ & . get (0).getDfCreateFlows() \\ & . get (0).getDfCreateFlows() \\ \end{tabular}$ 242 . $\frac{242}{243}$. The substitution of the set (0.000 addAssignment ("@desc = desc"); $\begin{tabular}{ll} 243 & // \,\, create & \,\, dataflow Model.\, getDfBoundaries\,().\, get \,0). \, get \,0'. \,get \,0'. \,get \,0'. \,set \,0'. \,$ 253 dataflowModel .getDfBoundaries ().get (0) .getDfDeleteFlows ()
254 .get (0) .set Filter ("[@code = code] @code");
// delete 256 dataflowModel .getDfBoundaries ().get (0) .getDfDeleteFlows ()
257 . get (0) .addDataonflow ("code");
// read filter 259 dataflowModel.getDfBoundaries ().get (0).getDfReadFlows ()
260 . get (0).setFilter ("[@code = code] @desc"); 261 // read
262 dataflowModel.getDfBoundaries ().get (0).getDfReadFlows ()
263 (apet (0).addDataonflow ("desc");
// update filter 265 dataflowModel.getDfBoundaries ().get (0).getDfUpdateFlows ()
266 . get (0).setFilter ("[@code = code] @code, @credit"); 267 // update assignment
268 dataflowModel.getDfB 268 dataflowModel.getDfBoundaries().get (0).getDfUpdateFlows()
269 .get (0).addAssignment ("@credit = credit");
// update // update 271 dataflowModel.getDfBoundaries().get(0).getDfUpdateFlows()
272 dataflowModel.getDfBoundaries().get(0).addDataflow("code");
273 dataflowModel.getDfDoundaries().get(0).getDfUpdateFlows()
274 get(0).addDataonflow("credit") $\frac{275}{276}$ 276 System . out . p r i n t l n (" (4) Data on f l o w s a r e added i n th e DataFlow model . ") ; 277 // −− // $\frac{278}{279}$ 279 $\frac{1}{279}$ // Generate Detailed DatFlow Model 280 DfDiagramToDfDiagram topRu l e6 = new DfDiagramToDfDiagram () ; 281 DfDiagram d e t a i l e d d a t a f l o wM o d e l = topRu l e6 . t r a n s l a t e (data f lowMod e l) ; $\frac{282}{283}$ 283
 ASTWriter writer6 = **new ASTWriter** (new

File ("Manage_Module_Detailed_DataFlowModel.xml"));

writer6.usePackage ("mde.dataflow.model", "xmlns:dfd");

writer6.writeDocument(detailed_dataflowModel);

287

writer6.cl 288 289 System.out.println("(5) Detailed DataFlow Model is Created"+
290 "(DataFlow Model −−> DataFlow Model)."); 291 // −− // $\begin{array}{c} 292 \\ 293 \end{array}$ $\begin{array}{c|c} 293 & // Generate Data Model \\ 294 & InfDiagram ToDDiagram \end{array}$ 294 InfDiagramToDDiagram topRule4 = new InfDiagramToDDiagram();
295 DDiagram dataModel = topRule4.translate(informationModel); $\frac{296}{297}$ 297
 ASTWriter writer4 = **new ASTWriter** (**new**
 File ("Manage_Module_DataModel.xml"));

writer4 .usePackage ("mde.data.model", "xmlns:data");

writer4 .usePackage ("mde.model", "xmlns:mod");

writer4 .writeDocument 302 writer 4. close (); $\frac{303}{304}$ 304 System . out . p r i n t l n (" (5) Data Model i s Cr eat ed "+ 305 " (I n f o rm a t i o n Model −−> Data Model) . ") ; 306 // −−− // 307 308 //Generate State Model for Screens Navigation
309 //(DataFlow Model ---> State Model)
310 **DidDiagramToStDiagram** topRule7 = **new DidDiagramToStDiagram**();
StDiagram stateModel=topRule7.translate(detailed_dataflowMode $\frac{311}{312}$ 313
 ASTWriter writer7 = **new**
 ASTWriter(**new File**("Manage_Module_StateModel.xml"));

315

writer7.writeDocument(stateModel);

writer7.writeDocument(stateModel); $\begin{array}{c} 317 \\ 317 \end{array}$ writer 7 . close (); $\frac{318}{319}$ $System.out.println("6) State Model is Create "+$

Task model: Module Management System

Impact model: Module Management System


```
34 <imp:ImpRole id="18" name="modify"><br>35 <imp:ImpTask ref="5"/>
36 </imp:ImpRole><br>
37 </imp:ImpRole id="19" name="module"><br>
38 </imp:ImpObject ref="7"/><br>
40 </imp:ImpUpdateFlow><br>
41 </imp:ImpUpdateFlow><br>
42 </imp:ImpBoundary>
```
Inofrmation model: Module Management System

```
Listing A.4: M anage Module informationModel.xml
```

```
1 </a><br>
2) <inf:Diagram xmlns:inf="me.information.model" id="0"><br>
3) <inf:Entity id="1" mame="Module"><br>
4) <inf:Attribute id="2" name="code" identifier="true" size="0"><br>
4) <fr(2) inf:Attribute id="2" name="code" identifi
\begin{array}{c}\n16 \\
16 \\
17\n\end{array} \begin{array}{c}\n\diagup \text{limit of} \\
\diagup \text{limit of} \\
\diagup \text{limit of} \\
\diagup \text{limit of} \\
\diagdown \text{matrix of} \\
\diagdown \text{matrix
```
DataFlow model (initial): Module Management System

```
Listing A.5: M anage Module DataFlowModel.xml
```


DataFlow model (detailed): Module Management System

Data Dependency model: Module Management System

Listing A.7: M anage Module DataModel.xml

```
1 <?xml version="1.0" encoding="UTF-8"?><br>
2 <data:DDiagram xmlns:data="mde.data:D="m"><br>
3 <data:D="m"> <data:D="mde.data"="mde.data:D="" ="" <data:DAttribute id="2" name="code" identifier="true" size="10"><br>
4 <mod:Type xm
  7 <data:DAttribute i d="4" name=" t i t l e " s i z e=" 30 ">
8 <mod:Type i d="5" name=" S t r i n g "/>
9 </data:DAttribute>
10 <data:DAttribute id="6" name=" credit" size="10"><br>11 <mod:Type id="7" name=" Integer"/><br></data:DAttribute>
13 <data:DAttribute id="8" name="desc" size="30"><br>14 <mod:Type id="9" name="String"/><br>15 </data:DAttribute>
\begin{array}{c|c}\n 16 & \checkmark / \text{data:} \text{DEntity} \\
 \hline\n 17 & \checkmark / \text{data:} \text{DDi,} \text{gram}\n \end{array}1/ data:DDiagram>
```
Screen State model: Module Management System


```
343 </state:StRole><br>344 <state:StRole i
\frac{344}{344} \left\{\begin{array}{c}\n\text{state:State} \\
\text{state:State} \\
\text{state:ref} \\
\text{state:ref} \\
\end{array} \right.\n\left\{\n\begin{array}{c}\n\text{name}^= \text{main} \cdot \text{task} \cdot \text{Modify} \\
\text{model}\n\end{array}\n\right\}345 <s ta t e : S ta t e r e f=" 35 "/>
346 </ sta te :S tRo le>
 \begin{array}{c|c}\n 347 & \text{\textendash} & \text{\textendash} & \text{\textendash}<br>
348 <br>
\text{\textendash} <br>
\text{\text\begin{array}{c} 348 \overline{) } \\ 349 \overline{) } \\ 349 \overline{) } \end{array} \begin{array}{c} \sqrt{,} \text{ 2141} \text{ } \text{ 2141} \text{ } \text{ 2241} \text{ } \text{ 2341} \text{ } \text{ 241} \text{ } \text{ 251} \text{ } \text{ 261} \text{ } \text{ 271} \text{ } \text{ 281} \text{ } \text{ 281} \text{ } \text{ 291} \text{ } \text{ 201} \text{ } \text{ 201} \text{ } \\langle/state:StDiagram>
```
GUI model: Module Management System


```
\begin{array}{c|c} 1 & \text{<} \text{equiv:GuIDiagram } x \text{ m} \text{ln} s : \text{gui:} \text{m} \text{d} \text{e} \text{m} \text{d} \text{d} \text{d} \text{m} \text{d} \text{d} \text{d} \text{m} \text{d} \text{2 <gui:GuiBoundary id="1" name="Add Module"><br>2 <gui:Window id="2" name="Input_code_title_credit_desc_Waiting" order="4"><br>4 <gui:Textfield id="3" name="code"/><br>2 <gui:Textfield id="4" name="title"/>
   6 <g u i :T e x t f i e l d i d="5" name=" c r e d i t "/>
7 <g u i :T e x t f i e l d i d="6" name=" d e s c "/>
8 <gui:Button i d="7" name=" Ex c ep t i on " e v en t=" Ex c ep t i on " e x i t=" f a l s e "/>
9 <gui:Button i d="8" name=" Input " e v en t=" Input " e x i t=" f a l s e " />
10 </gui:Window>
11 <gui:Window i d="9" name=" I n p u t c o d e t i t l e c r e d i t d e s c W a i t i n g E r r o r " o r d e r="0" e r r o r=" t r u e ">
12 <gu i :Labe l i d=" 10 " name=" i n p u t c o d e t i t l e c r e d i t d e s c W a i t i n g E r r o r w a r n i n g "
13 t e x t=" Nu l l v a l u e not a c c e p t e d " />
14 <gui:Button i d=" 11 " name=" i n i t i a l i s e " e v en t=" i n i t i a l i s e " e x i t=" f a l s e "/>
15 <br>
\langle \text{gui:Window} \rangle <br>
\langle \text{gui:Window} \rangle <br>
\langle \text{gui:Window} \rangle <br>
\langle \text{gui:Value} \rangle <br>
\langle \text{gui:Label id} = "13" \text{ name} = "code" \text{ text} = "code" / \rangle<br>
\langle \text{gui:Label id} = "14" \text{ name} = "relit" \text{ text} = "tilte" / \rangle<br>
\langle \text{gui:Label id} = "14" \text{ name} = "relit" \text{ text} = "relit" / \rangle<br>
\langle \text{gui:Label id} = "15" \text{ name24 </gui:Window<br>25 < gui:Window id="20" name="Create_Module_Ready_Error" order="0" error="true"><br>26 <gui:Label id="21" name="Create_Module_Ready_Error_warning"<br>27 text="Connection to the Data source is fail"/><br>28 <
\begin{array}{c}\n 29 \\
 30\n \end{array} \begin{array}{c}\n 29 \\
 7 \text{gui:Window}\n\end{array}\begin{array}{c|c}\n 30 & \lt/\n  \text{gui:GuiBoundary}\n  \hline\n  \text{20:Gui:GuiBoundary}\n \end{array}31 <gui:GuiBoundary i d=" 23 " name="Mod ify Module ">
32 <gui:Window i d=" 24 " name=" I n p u t c o d e c r e d i t W a i t i n g " o r d e r="4">
33 <g u i :T e x t f i e l d i d=" 25 " name=" cod e " />
34 <g u i :T e x t f i e l d i d=" 26 " name=" c r e d i t "/>
35 <gui:Button i d=" 27 " name=" Ex c ep t i on " e v en t=" Ex c ep t i on " e x i t=" f a l s e "/>
36 <gui:Button i d=" 28 " name=" Input " ev en t=" Input " e x i t=" f a l s e "/>
37 </gui:Window>
38 <gui:Window i d=" 29 " name=" I n p u t c o d e c r e d i t W a i t i n g E r r o r " o r d e r="0" e r r o r=" t r u e ">
39 <gu i :Labe l i d=" 30 " name=" i n p u t c o d e c r e d i t W a i t i n g E r r o r w a r n i n g "
40 t e x t=" Nu l l v a l u e not a c c e p t e d "/>
41 <gui:Button i d=" 31 " name=" i n i t i a l i s e " e v en t=" i n i t i a l i s e " e x i t=" f a l s e "/>
42 </gui:Window>
43 <gui:Window id="32" name="Read_Module_Code_Credit_Ready" order="3"><br>
<gui:Label ref="13"/><br>
<gui:Label ref="13"/><br>
<gui:Label ref="13"/><br>
<gui:Label ref="13"/><br>
<gui:Button id="33" name="Read" event="Read" exit="false"/
</gui:Window><br>
cgui:Window id="35" name="Read_Module_Code_Credit_Ready_Error" order="0" error="true"><br>
<gui:Label id="36" name="Read_Module_Code_Credit_Ready_Error_warning"<br>
51 text="Connection to the Data source is fail"
53 </gui:Window>
54 <gui:Window i d=" 38 " name="Wr ite Modu le Ready " o r d e r="2 ">
55 <gu i :Labe l r e f=" 13 "/>
56 <gu i :Labe l r e f=" 15 "/>
57 <gui:Button i d=" 39 " name="Wr ite " e v en t="Wr ite " e x i t=" f a l s e "/>
58 <gui:Button i d=" 40 " name=" e x i t " e v en t=" e x i t " e x i t=" t r u e " />
59 <gui:Button i d=" 41 " name=" Ex c ep t i on " e v en t=" Ex c ep t i on " e x i t=" f a l s e "/>
\begin{array}{c|c}\n 59 & \text{Squi:Button} \\
 60 & \text{Squi:Window}\n\end{array}%61<br>
<gui:Window id="42" name="Write_Module_Ready_Error" order="0" error="true"><br>
<gui:Label id="43" name="Write_Module_Ready_Error_warning"<br>
ext="Connection to the Data source is fail"/><br>
(anie and text="exit" exit="exit"
\begin{array}{c|c}\n 65 & \times \text{gui:Window}\n 66 & & \times/\text{gui:Guiboundary}\n\end{array}66 <br>
\langle \text{gui:Guiboundary} \rangle<br>
\langle \text{gui:Guiboundary} \rangle id="45" name="Delete Module"><br>
68 \langle \text{gui:Guiboundary} \rangle id="46" name="Input_code_Waiting" order="4"><br>
\langle \text{gui:Textfield} \text{ id} = "46" \text{ name="code"}/\rangle<br>
\langle \text{gui:Button id="48" name="Exception" event="Exception" exit="false"}/\rangle<br>
\langle \text{gui:Button72 </gui:Window>
73 <gui:Window i d=" 50 " name=" I n p u t c o d e W a i t i n g E r r o r " o r d e r=" 0" e r r o r=" t r u e ">
74 <gu i :Labe l i d=" 51 " name=" i n p u t c o d e W a i t i n g E r r o r w a r n i n g "
75 t e x t=" Nu l l v a l u e not a c c e p t e d "/>
76 <gui:Button i d=" 52 " name=" i n i t i a l i s e " e v en t=" i n i t i a l i s e " e x i t=" f a l s e "/>
```

```
77 </gui:Window><br>78 </gui:Window
    %;<br>
<gui:Window id="53" name="Delete_Module_Ready" order="2"><br>
- coui:Label ref="13"/>
     79 <gu i :Labe l r e f=" 13 "/>
80 <gui:Button i d=" 54 " name=" D e l e t e " e v en t=" D e l e t e " e x i t=" f a l s e "/>
81 <gui:Button i d=" 55 " name=" e x i t " e v en t=" e x i t " e x i t=" t r u e " />
82 <gui:Button i d=" 56 " name=" Ex c ep t i on " e v en t=" Ex c ep t i on " e x i t=" f a l s e "/>
    \begin{array}{c|c}\n 83 & & \text{\textless}\xspace & \text{\textless}\x84 <gui:Window id="57" name="Delete_Module_Ready_Error" order="0" error="true"><br>
85 <gui:Label id="58" name="Delete_Module_Ready_Error_warning"<br>
86 text="Connection to the Data source is fail"/><br>
87 <gui:Button id="59" nam
     \begin{array}{ll} 88 & </\text{gui:Window>}\\ 89 & </\text{gui:GuiBoundary} \\ \text{Equ:GuiBoundary id="60" name="See Description">\\  \end{array}\\ \begin{array}{ll} \text{equ:GuiBoundary id="60" name="Spec Description">\\ \text{equ:GuiBoundary id="61" name="Inputcode, \textit{Code}.\textit{Waiting" order="4">}\\ 91 & <\text{gui:Textfield id="63" name="Code"/\\ & <\text{gui:Butif } \textit{Textfield} \textit{if } \textit{See} \textit{end} \end{array}\\ \begin{array}{ll} 84 & <\text{gui:Uniform id="63" name="\begin{array}{c|c}\n 95 \\
 95 \\
 96 \\
 \hline\n 87\n \end{array}
<br>
\begin{array}{c}\n 2 \text{ gain:} \text{arrows}\n  \times \text{ g} \text{ air:} \text{ bar} \times \text{ m} \times \text{96   |   Gui:Window id="65" name="Input_code_Waiting_Error" order="0" error="true"><br>97   |   Gui:Label id="66" name="input_code_Waiting_Error_warning"<br>98   |   Gext="Null value not accepted"/><br>29   Gui:Button id="67" name=
 100 </gui:Window>
101 <gui:Window i d=" 68 " name=" Read Desc Ready " o r d e r="3">
 102 <gu i :Labe l r e f=" 13 "/>
103 <gui:Button i d=" 69 " name="Read" e v en t="Read" e x i t=" f a l s e " />
104 <gui:Button i d=" 70 " name=" e x i t " e v en t=" e x i t " e x i t=" t r u e " />
105 <gui:Button i d=" 71 " name=" Ex c ep t i on " e v en t=" Ex c ep t i on " e x i t=" f a l s e "/>
\begin{array}{c|c}\n 106 \\
 107 \\
 \hline\n \end{array} \begin{array}{c}\n 2 \text{ units} \\
 3 \text{ units} \\
 4 \text{ units} \\
 \end{array}107 <gui:Window i d=" 72 " name=" R ead D es c R eady Error " o r d e r=" 0" e r r o r=" t r u e ">
108 <gu i :Labe l i d=" 73 " name=" R e ad D e s c R e ady E r r o r w a rn in g "
 109 text="Connection to the Data source is fail"/><br>110 <gui:Button id="74" name="exit" event="exit" exit="false"/>
\begin{array}{c}\n 111 \\
 111 \\
 \hline\n 112\n \end{array} \begin{array}{c}\n  \sqrt{g} \text{an} \text{B} \text{arivon} \\
  \sqrt{g} \text{an} \text{B} \text{div} \text{S} \\
  \sqrt{g} \text{mi} \cdot \text{G} \text{ui} \text{B} \text{ounda}\n \end{array}112 \langlegui:GuiBoundary><br>
113 \langlegui:GuiBoundary i
 (113) \leq (113) \leq (113) \leq (113) \leq (113) \leq (114) 
\begin{array}{c}\n\sqrt{g} \text{ and } \\
119 \\
120\n\end{array} \begin{array}{c}\n\sqrt{g} \text{ and } \\
\sqrt{g} \text{ will} \\
\\begin{array}{c|c}\n 120 & \checkmark/\text{gui:GuiBoundary}\n \hline\n 121 & \checkmark/\text{gui:GuiDigram}\n \end{array}\langle/gui:GuiDiagram>
```
Database and Query model: Module Management System

Listing A.10: Manage_Module_SchemaModel.xml

```
1 <?xml version="1.0" encoding="UTF-8"?><br>2 <dbs:Schema xmlns:dbs="mde.dbs.model" id="0" name="database"><br>3 <dbs:Table id="1" name="Module"><br>4 <dbs:Column id="2" name="code" size="10">
   5 <mod:Type xmlns:mod="mde.model" id="3" name="INTEGER"/> </dbs:Column>
   7 <dbs:Column i d="4" name=" t i t l e " s i z e=" 30 ">
8 <mod:Type i d="5" name="VARCHAR"/>
 9 </dbs:Column><br>10 <dbs:Column id="6" name="credit" size="10"><br>11 <mod:Type id="7" name="INTEGER"/>
 22 </dbs:Column><br>
21 </dbs:Column id="8" name="desc" size="30"><br>
14 </mod:Type id="9" name="VARCHAR"/><br>
21 </dbs:Column><br>
20 </dbs:PrimaryKey id="10">
10<br>
17 <dbs:Column r e f="2"/><br>
\langledbs:Column r e f="2"/><br>
\langle/dbs:FrimaryKey>
 \begin{tabular}{c|c|c} $18$ & $\sensureensure{\fbox{-}}$ & $\zeta$/dbs:PrimaryKey> \\ $19$ & $\zeta$/dbs:Table20\\ $20$ & $\zeta$/dbs:Problemi$ & $11" name="readCode_Credit" table="Modue">\n21 & $\zeta$/dbs:Argument id="12" type="Integer" name="code" in="true"/\n22 & $\zeta$/dbs:Argument id="13" type="Integer" name="code" in="true"/\n23 & $\zeta$/dbs:Argument id="13" name="SetLECT.Code, Credit" & $out="true"/\n24 & $\zeta$/dbs:Project id="15" name="proj">\n25 & $\zeta$28 </ dbs:Project>
29 <dbs:Relation i d=" 18 " name="Module ">
30 <d b s :F i l t e r i d=" 19 " name=" f i l t e r 1 ">
31 <dbs:Relation i d=" 20 " name="Module "/>
32 <dbs:Operator i d=" 21 " typ e=" b o o l e an " symbo l=" e q u a l s ">
33 <dbs:Column i d=" 22 " name=" cod e " s i z e=" 10 ">
\frac{\text{mod:Type id}}{\text{mod:Type id}} = \frac{\text{name="interior: Time}}{\text{name="interior: Time}} = \frac{\text{source: 1}}{\text{sqrt} \cdot \text{6}}\begin{array}{c|c}\n 35 & & \quad \langle / \text{dbs:Column} \rangle \\
 36 & & \quad \langle \text{dbs:Variable} \end{array}36<br>36 <dbs:Variable id="24" type="INTEGER" name=" code" /><br>37 </dbs:Variable id="24" type="INTEGER" name=" code" /><br>38 </dbs:Filter>
 37 </dbs:Operator><br>38 </dbs:Filter>
```

```
\begin{array}{c}\n 39 \\
 40\n \end{array} \begin{array}{c}\n  </\text{dbs:Relation}\n  <\n  / \text{dbs:Query}\n \end{array}\begin{array}{c}\n 40 \\
 41 \\
 \hline\n 41\n \end{array} \begin{array}{c}\n 7 \\
 7 \text{dbs:Query} \\
 6\n \end{array}\begin{tabular}{c|c|c} $\sphericalangle$ & $\sphericalangle\begin{array}{ll} \text{50} & < \!\! \text{/dbs:Project>} \\ \text{51} & < \!\! \text{dbs:Relation id="32" name="Module"} \\ \text{52} & < \!\! \text{dbs:Relation id="33" name="filter 1"'} \\ \text{53} & < \!\! \text{dbs:Relation id="34" name="Module"} \text{/>} \\ \text{54} & < \!\! \text{dbs:Operator id="35" type="boolean" symbol="equals"} \\ \text{55} & < \!\! \text{dbs:Column id="36" name="code" size="10">} \\ \text{56} & < \!\! \text{mod:Type ref="33"}/ \\ \text{57} & < \58 <dbs:Variable id="37" type="INTEGER" name=" code" />
</dbs:Operator>
   \left\{\n \begin{array}{c}\n \sqrt{d}b \sin \theta \\
 \sqrt{d}b \sin \theta\n \end{array}\n \right\}\langle d bs: Filter>
   \begin{array}{c|c}\n\hline\n61 & & \langle \text{dbs:Relation} \rangle \\
62 & & \langle \text{dbs:Query} \rangle\n\end{array}\begin{array}{c|c}\n 62 & & \sqrt{\text{abs}}:\text{Query}\n 63 & & \sqrt{\text{abs}}:\text{Proceedure}\n\end{array}%<br/>\begin{array}{ll} &<\!\!\!/\mathrm{dbs: Procedure}\!\> \mathrm{&<}\!\!\mathrm{&<}\!\!\mathrm{dbs: Procedure}\!\> \mathrm{&<}\!\!\mathrm{&<}\!\!\mathrm{dbs:Procedure}}\!\!\!\mathrm{&<}\!\!\mathrm{&<}\!\!\mathrm{dbs:Procedure\;id=``38" name="createModule" table="Module">\n66 &<\!\!\! \mathrm{dbs:Argument\;id=``39" type="Integer" name="code" in="true"}/\!\!>\!\! \mathrm{&<}\!\!\! \mathrm{&<}\!\!\! \mathrm{&<}\!\!\! \mathrm{&<}\!\!\!\! \mathrm{&<}\!\!\! \mathrm{&<}\!\!\! \mathrm{\begin{array}{rcl} \text{72} & \text{<mod:Type id="46" name="NTEGER" />} \ \text{73} & <\text{/dbs:Column}> \ \text{74} & <\text{dbs:Variable id="47" type="NTEGER" name="code" />} \ \text{75} & <\text{dbs:Operator}> \ \text{76} & <\text{dbs:Operator} \ \text{78} & <\text{dbs:Operator} \ \text{79} & <\text{dbs:Operator} \ \text{70} & <\text{dbs:Operator} \ \text{71} & <\text{dbs:Operator} \ \text{72} & <\text{dbs:Operator} \ \text{73} & <\text{dbs:Operator75 </dbs:Operator>
76 <dbs:Operator i d=" 48 " typ e=" b o o l e an " symbo l=" A s s i gn ">
77 <dbs:Column i d=" 49 " name=" t i t l e " s i z e=" 30 ">
78 <mod:Type i d=" 50 " name="VARCHAR"/>
    79 </dbs:Column>
80 <dbs:Variable i d=" 51 " typ e="VARCHAR" name=" t i t l e " />
    81<br>
</dbs:Operator><br>
<dbs:Ocerator id="52" type="boolean" symbol="Assign"><br>
83<br>
<dbs:Column id="53" name="credit" size="10"><br>
84<br>
<mod:Type id="54" name="NTEGER"/><br>
</dbs:Column> id="54" name="NTEGER"/><br>
</dbs:Column> id="
   87 \langle dbs:0perator>
    88 <dbs:Operator id="56" type="boolean" symbol="Assign"><br>89 <dbs:Column id="57" name="desc" size="30"><br>90 <mod:Type id="58" name="VARCHAR"/>
    91 </dbs:Column>
92 <dbs:Variable i d=" 59 " typ e="VARCHAR" name=" d e s c " />
   |93| \langle dbs:Operator>
   \begin{array}{c|c} 94 & <|{\rm\bf{dbs:Create}} > \ 95 & <|{\rm\bf{dbs: Procedure}} > \end{array}95 <br>
\langle/dbs:Procedure><br>
36 <br>
\langledbs:Argument id="60" name="deleteModule" table="Module"><br>
37 <br>
\langledbs:Delete id="62" name="DELETE_Code"><br>
38 <br>
\langledbs:Relation id="63" name="Delete Module"><br>
39 <br>
\langledbs:Relation id="
 104 <mod:Type i d=" 68 " name="INTEGER"/>
105 </dbs:Column>
106 (dbs:Variable id="69" type="INTEGER" name="code" /><br>107 (dbs:Operator)
 107 </dbs:Operator>
108 </ d b s :F i l t e r>
\begin{array}{c|c}\n 109 \\
 110 \\
 \hline\n 110\n \end{array} \begin{array}{c}\n 74.44 \text{ Hz} \\
 74.4 \text{ Hz} \\
 110 \text{ keV} \\
 110\n \end{array}\begin{array}{c}\n 110 \\
 111 \\
 \hline\n 6/ \text{dbs:Decedure}\n \end{array}111 <br>
\langle/dbs:Procedure><br>
112 <br>
\langledbs:Procedure id="70" name="writeModule" table="Module"><br>
113 <br>
\langledbs:Argument id="71" type="NYEGER" name="code" in="true"/><br>
\langledbs:Argument id="73" name="NYEGER" name="credit" ino
\frac{120}{120} <dbs:Relation id="77" name="Module"/>
 \begin{array}{ll} \text{{\bf 121}} & < \text{{\bf /dbs:Project>}}\\ \text{{\bf 122}} & < \text{{\bf dbs:Relation} id="78" name="Module"}> \\ \text{{\bf 123}} & < \text{{\bf dbs:Relation} id="78" name="Module"}> \\ \text{{\bf 124}} & < \text{{\bf dbs:Relation} id="79"}> \\ \text{{\bf 125}} & < \text{{\bf dbs:Relation} id="78" name="Module"}/> \\ \text{{\bf 126}} & < \text{{\bf dbs:Operator} id="81" type="boolean" symbol="equals"}> \\ \text{{\bf 126}} & < \text128 </dbs:Column>
129 <dbs:Variable i d=" 84 " typ e="INTEGER" name=" cod e " />
 130 </dbs:Operator>
131 </ d b s :F i l t e r>
\frac{132}{132} </dbs:Relation>
```

```
\begin{array}{ll} \text{133} & <\!\!\!\!/\text{dbs:Query}\!\!\!&<\!\!\!\! \text{dbs:Query}\!\!\!&<\!\!\!\! \text{dbs:Update id="85"&> \\ \text{135} & <\!\!\!\! \text{dbs:Update id="86" &<\!\!\!\! \text{type="boolean" symbol="assign"&> \\ \text{136} & <\!\!\!\! \text{dbs:Column id="87" name="credit" size="10"&> \\ \text{137} & <\!\!\!\! \text{dbs:Column} &<\!\!\!\! \text{t= "88" name="INTER"&> \\ \text{138} & <\!\!\!\! \text{139<br>
149 </dbs:Variable id="89" type="INTEGER" name=" credit"/><br>
</dbs:Variable id="89" type="INTEGER" name=" credit"/>
                                                 1/dbs:Operator>
\begin{array}{c|c}\n 141 \\
 142 \\
 \hline\n 142\n \end{array} \begin{array}{c}\n 2 \text{ hours per}-1 \\
 7 \text{ obs:Update}\n \end{array}\begin{array}{c|c}\n 142 & \lt/dbs: \textbf{Procedure} \\
 \hline\n 143 & \lt/dbs: \textbf{Schema}\n \end{array}\langle/dbs:Schema>
```
OOP Code model: Module Management System

Listing A.11: Manage_Module_CodeModel.xml

```
1 <?xml version="1.0" encoding="UTF-8"?><br>
<code:Diagram xmlns:code="mde.code.moc
     2 | <code:Diagram xmlns:code="mde.code.model" id="0"><br>3 | <code:Clazz id="1" name="Module" visible="public"><br>4 | <code:Attribute id="2" name="code" visible="private">
     5 <mod:Type xmlns:mod="mde.model" id="3" name="int"/><br>
</code:Attribute><br>
7 </code:Attribute id="4" name=" title" visible=" private"><br>
8 <mod:Type id="5" name=" String"/><br>
</code:Attribute>
\begin{array}{lll} \text{10} & \text{2} & \text{2} & \text{2} \\ & \text{3} & \text{2} & \text{2} \\ 11 & \text{3} & \text{3} & \text{2} \\ \end{array} wisible=" private">
 11 \langle \text{mod:Type id="7" name="int" / \rangle<br>
2 \langle \text{code:Attribute} \rangle<br>
3 \langle \text{code:Attribute} \rangle<br>
4 \langle \text{code:Method} \rangle<br>
4 \langle \text{21<br>
\left\{\begin{array}{r}\n\text{codewenoud.}\n\text{code:Method id="14" name="readDesc" visible="public">\n\end{array}\right.\n\left.\begin{array}{r}\n\text{code:Method id="15" name="void"}/\n\end{array}\right.\n\left.\begin{array}{r}\n\text{code:Method id="15" name="void"}/\n\end{array}\right.\n\left.\begin{array}{r}\n\text{value:} \n\end{array}\right.\n\left.\begin{array}{r}\n\text{value:} \n\end{array}\right.\n\left.\begin{array}{r}\n\text{value:} \n\end{array}\right.\n\left.\begin{array}{r}\n\text{value:}22 <mod:Type id="15" name="void"/><br>23 <code:Argument id="16" type="int" name="code"/><br>24 <code:Call id="17" name="readDesc"/>
 \begin{tabular}{ll} $25$ & <code>&lt; /\text{code:Method}\text{&lt;}\\ & &lt; \text{code.Method}\text{&lt;}\\ & &lt; \text{code.Method} \text{id} = "18" \text{ name} = " \text{createModule" } \text{visible} = " \text{public"}\text{&lt;}\\ & &lt; \text{mod}: \text{Type id} = "19" \text{ name} = " \text{void"} / \text{&lt;}\\ & &lt; \text{code.Mgrument id} = "20" \text{ type} = " \text{int" name} = " \text{code"} / \text{&lt;}\\ & &lt; \text{code.Argument id} = "21" \text{ type} = " \33 </code:Method>
34 <code:Method i d=" 25 " name=" d e l e t eM odu l e " v i s i b l e=" p u b l i c ">
 35 <mod:Type i d=" 26 " name=" v o id "/>
36 <code:Argument i d=" 27 " typ e=" i n t " name=" cod e "/>
37 <code :Ca l l i d=" 28 " name=" d e l e t eM odu l e "/>
 38 <br>
\left\{\begin{array}{ll}\n<\rangle \text{code:Method}\n> \quad \text{&lt;code:Method id="29" name="writeModule" visible="public"\n> \quad \text{&lt;node:Method id="30" name="void"\n> \quad \text{&lt;code:Argument id="31" name="code"\n> &lt; \quad \text{&lt;code:Call id="32" name="writeModule"\n}\n\end{array}\right.43 </code:Method>
44 <code:Constructor i d=" 33 " name="Module " v i s i b l e=" p u b l i c ">
45 <code:Operator i d=" 34 " typ e=" b o o l e an " symbo l=" a s s i g n ">
46 <c o d e : I d e n t i f i e r i d=" 35 " typ e=" i n t " name=" cod e " s c op e=" Ob j e ct "/>
47 <c o d e : I d e n t i f i e r i d=" 36 " typ e=" i n t " name=" cod "/>
 48 </code:Operator>
49 <code:Operator i d=" 37 " typ e=" b o o l e an " symbo l=" a s s i g n ">
50 <c o d e : I d e n t i f i e r i d=" 38 " typ e=" S t r i n g " name=" t i t l e " s c op e=" Ob j e ct "/>
51 <c o d e : I d e n t i f i e r i d=" 39 " typ e=" S t r i n g " name=" t i t "/>
 52 </code:Operator>
53 <code:Operator i d=" 40 " typ e=" b o o l e an " symbo l=" a s s i g n ">
54 <c o d e : I d e n t i f i e r i d=" 41 " typ e=" i n t " name=" c r e d i t " s c op e=" Ob j e ct "/>
55 <c o d e : I d e n t i f i e r i d=" 42 " typ e=" i n t " name=" c r e d "/>
 56 </code:Operator>
57 <code:Operator i d=" 43 " typ e=" b o o l e an " symbo l=" a s s i g n ">
58 <c o d e : I d e n t i f i e r i d=" 44 " typ e=" S t r i n g " name=" d e s c " s c op e=" Ob j e ct "/>
59 <c o d e : I d e n t i f i e r i d=" 45 " typ e=" S t r i n g " name=" d e s "/>
60<br>61 </code:Operator><br>61 </code:Constructor>
 61 <br>
\left\{\begin{array}{ll} <\!\!\ensuremath{\langle} \text{code:} \text{Consider}} \right. \ \text{code:} \ \text{code:}
```

```
73 \left\{\n \begin{array}{c}\n \left\{\n \begin{array}{c}\n \leftarrow \text{code:Constructor} \quad id = "53" \quad name = "Add Module" \right|\n \left\{\n \end{array}\n \right\}\n \left\{\n \begin{array}{c}\n \left\{\n \begin{array}{c}\n \leftarrow \text{code:Class}\n \end{array}\n \right\}\n \left\{\n \begin{array}{c}\n \leftarrow \text{code:Class}\n \end{array}\n \right\}\n \left\{\n \begin{array}{c}\n \leftarrow \text{code:Class}\n \end{array}\n \right\}\n \left\{\n \begin{array}{c}\n \leftarrow \text{code:Class}\n \end\begin{array}{r} 74 \overline{74} \overline{75} \overline{75} \overline{75} \overline{76} \overline{76} \overline{76} \overline{76} \overline{76} \overline{76} \overline{76} \overline{76} \overline{76} \overline{77} \overline{78} \overline{78} \overline{78} \overline{78} \overline{78} \overline{78} \overline{78} \overline{78} \end{array}%75 <code:Clazz id="54" name="Modify Module" visible="public"><br>76 <code:Attribute id="55" name="Module" visible="private"><br>77 <modi:Type id="56" name="Module"/><br>8 </code:Attribute>
    \begin{array}{ll} \text{{\tt code:Attribute id="57" name=" Staff" visible=" private"}\!\!>\!80 & \text{{\tt xmod:Type id="58" name=" Staff"} }\\ \text{81} & \text{{\tt xmod:Type id="58" name=" Startf"}}\!\!>\!81 & \text{{\tt xmod:Type id="58" name=" Startf"}}\!\!>\!82 & \text{{\tt code:Method id="59" name="Run" visible=" public" staticm=" true"}\!\!>\!83 & \text{{\tt xmod:Type id="60" name="void"}}\!\!>\!84 & \text{{\tt xmod:Type id"}}\!\!>\!88 & \text{{\tt xmod:Type id="60"}}\!\!>\!8\begin{tabular}{ll} \textbf{85} & & & & & & & & & \\ \textbf{86} & & & & & & & & \\ \textbf{87} & & & & & & & & & \\ \textbf{88} & & & & & & & & \\ \textbf{89} & & & & & & & & \\ \textbf{80} & & & & & & & & \\ \textbf{81} & & & & & & & & \\ \textbf{82} & & & & & & & & \\ \textbf{83} & & & & & & & & \\ \textbf{84} & & & & & & & & \\ \textbf{85} & & & & & & & & \\ \textbf{86} & & & & & & & & \\ \textbf{87} & & & & & & & & \\ \text96 </code:Method><br>97 </code:Construc
  97 <code:Constructor id="69" name="Delete Module"/><br>98 </code:Clazz><br>99 <code:Clazz id="70" name="See Description" visible
 98 <br>
\leq /code:Clazz id="70" name="See Description" visible="public"><br>
\leq code:Clazz id="70" name="See Description" visible="public"><br>
\leq code:Attribute id="71" name="Module" visible="private"><br>
\leq (mod:Type id
 102 </code:Attribute><br>103 <code:Attribute id="73" name="Student" visible="private"><br>104 <eode:Attribute id="74" name="Student"/>
 105 </code:Attribute><br>106 <code:Method id="75" name="Run" visible="public" staticm="true"><br>107 <emod:Type id="76" name="void"/>
\begin{array}{c|c}\n 108 & & \checkmark \text{code:Method}\n \end{array}<br>
109 \checkmark \text{code:Construct}109 <code:Constructor id="77" name="See Description"/><br>110 </code:Clazz>
\begin{array}{c} 110 \\ 111 \end{array} </code:Clazz>
                    \langle/code:Diagram>
```
MySQL Script: Module Management System


```
1 \vert -- \; Database \; \; Creation\frac{2}{3}CREATE DATABASE sysDatabase;
       USE sysDatabase;
 \frac{4}{5}6 −− S t r u c t u r e f o r t a b l e ' Modu le '
  \frac{7}{8}\begin{array}{c|c} 8 & \textbf{CREATE TABLE Module} \\ \hline 9 & \textbf{Code INT}(10) \textbf{ NOT} \end{array}\begin{bmatrix} 9 \\ 9 \\ 10 \end{bmatrix} code INT(10) NOT NULL\begin{array}{c|c} \text{10} & \text{title} & \text{VARCHAR}(30) \,, \ \text{credit INT}(10) \,, \end{array}12 desc VARCHAR(30)<br>13 PRIMARY KEY(code
           PRIMARY KEY(code) );
\begin{array}{c} 14 \\ 15 \end{array}-- Structure of Procedure 'readCode_Credit'
\frac{16}{17}17 DELIMITER //
\frac{18}{19}19 CREATE PROCEDURE readCode_Credit (IN code INT(10), OUT credit INT(10)) BEGIN
\begin{array}{c|c} 20 & \text{BEGIN} \ 21 & \text{SELJ} \ 22 & \text{WHE} \end{array}SELECT Module.credit INTO credit FROM Module<br>WHERE
\begin{array}{c|c} \textbf{22} & \textbf{WHERE} \ \textbf{23} & \text{Module. code = code ;} \ \textbf{24} & \textbf{END} \end{array}\frac{25}{26}26 DELIMITER //
\frac{27}{28}−− Structure of Procedure 'readDesc'
\frac{29}{30}30 DELIMITER //
\frac{31}{32}32 CREATE PROCEDURE readDesc (IN code INT(10), OUT desc VARCHAR(30))<br>33 BEGIN SELECT Module desc INTO desc FROM Module
           BEGIN
34 SELECT Module . desc INTO desc FROM Module
35 WHERE<br>36 Module
                \overline{\text{Module} \cdot \text{code}} = \text{code};END //
\frac{37}{38}<br>38
       39 DELIMITER //
rac{40}{41}−− Structure of Procedure 'createModule'
\frac{42}{43}43 DELIMITER //
\frac{18}{44}
```

```
45 CREATE PROCEDURE createModule (IN code INT(10), IN title VARCHAR(30), IN credit INT(10), IN desc VARCHAR(30))
\begin{array}{c|c}\n 46 & \text{BEGIN} \\
 47 & \text{INSI} \\
 48 & \text{END}\n\end{array}INSERT INTO Module VALUES (code, title, credit, desc);
           END //
\frac{49}{50}50 DELIMITER //
 51
        −− Structure of Procedure 'deleteModule'
 53
       54 DELIMITER //
\begin{array}{c} 53 \\ 54 \\ 55 \\ 56 \end{array}56 CREATE PROCEDURE deleteModule (IN code INT(10))<br>57 BEGIN<br>DELETE FROM Module<br>59 WHERE
           BEGIN
 58 DELETE FROM Module<br>59 WHERE
 \begin{array}{c|c} 60 & \text{code} = \text{code}; \ \hline 61 & \text{END} \end{array}62<br>6363 DELIMITER //
64<br>65−− Structure of Procedure 'writeModule'
\frac{66}{67}67 DELIMITER //
68<br>69CREATE PROCEDURE writeModule (IN code INT(10), INOUT credit INT(10))<br>
70 BEGIN<br>
CREATE Module<br>
CREATE Module<br>
CREATE Module credit<br>
TA MHERE<br>
Module code = code;<br>
TAND //
           BEGIN
               JOHN<br>UPDATE Module
               \operatorname{\bf SET}Module. credit = credit\begin{array}{c|c} \text{74} & \text{WHERE} \ \text{75} & \text{Module. code} = \text{code}; \ \text{76} & \text{END} \end{array}^{+77}_{-78}78 DELIMITER //
```
Java Swing Source Code: Module Management System

Class: Manage Module Main Menu Waiting

Listing A.13: Manage_Module_Main_Menu_Waiting.java

1	import java.awt.GridLayout;
$\mathbf{2}$	import java.awt.event.ActionEvent;
3	import java.awt.event.ActionListener;
4	import javax.swing.*;
5	import java.sql. CallableStatement;
6	import java.sql. Connection;
7	import java.sql.DriverManager;
8	import java.sql.ResultSet;
9	import java.sql.SQLException;
10	import java.sql.Statement;
11	import java.sql.Types;
12	
13	public class Manage_Module_Main_Menu_Waiting extends JFrame {
14	public Manage Module Main Menu Waiting (String win Title) {
15	
16	super(winTitle);
17	
18	// Create Swing Components
19	JPanel main_menu_panel = new JPanel();
20	final JLabel main_menuLab = new JLabel();
21	main_menu_panel.add(main_menuLab);
22	JPanel message_panel = new JPanel();
23	final JLabel messageLab = new JLabel("Main Menu");
24	message_panel.add(messageLab);
25	JPanel $add_model = new JPanel()$;
26	JButton add_module = new JButton ("Add Module");
27	add_module_panel.add(add_module);
28	
29	// Add Action to the button
30	add_module.addActionListener(new ActionListener() {
31	public void actionPerformed (ActionEvent e) {
32	$Input_code_title_credit_desc_Waiting$ next Window = new
33	Input_code_title_credit_desc_Waiting("Input_code_title_
34	credit_desc_Waiting");
35	nextWindow.setSize(400, 300);
36	nextWindow.setDefaultCloseOperation(JFrame.EXIT_ON_CLOSE);
37	nextWindow.setLocationRelativeTo(null);
38	$nextWindow.setVisible(\textbf{true});$
39	dispose()
40	ł
41	$\}$);
42	
43	JPanel delete_module_panel = new JPanel();
44	JButton delete_module = new JButton ("Delete Module");
45	delete_module_panel.add(delete_module);

```
\frac{46}{47}\begin{array}{c|c|c|c|c} \hline \textbf{47} & & \textbf{48} & \textbf{49} & \textbf{48} & \textbf{49} & \textbf{40} & \textbf{40} & \textbf{40} & \textbf{40} & \textbf{40} & \48 delete_module.addActionListener(new ActionListener(){<br>public void actionPerformed(ActionEvent e){<br>50 mput_code_Waiting nextWindow = new<br>finput_code_Waiting("Input_code_Waiting");
 52 nextWindow . s e t S i z e ( 4 0 0 , 3 0 0 ) ;
53 nextWindow . s e t D e f a u l t C l o s eO p e r a t i o n (JFrame . EXIT ON CLOSE ) ;
54 nextWindow . s e t L o c a t i o n R e l a t i v e T o ( nu l l ) ;
\begin{array}{c|c} 55 \\ 55 \\ 15008 \end{array} nextWindow setVisible (true);
                                      {\rm dispose~()~;}\frac{57}{58}58 } ) ;
\frac{59}{60}60 JPanel modify_module_panel = new JPanel ();<br>61 JButton modify_module = new JButton ("Modify Module");<br>62 modify_module_panel.add(modify_module);
63<br>6464 // Add Action to the button<br>65 modify_module.addActionListe<br>66 public void actionPerform
 65<br>
modify_module.addActionListener(new ActionListener(){<br>
public void actionPerformed(ActionEvent e){<br>
Input_code_credit_Waiting nextWindow = new<br>
Input_code_credit_Waiting("Input_code_
 69 credit_Waiting");<br>
70 nextWindow.setSize(400, 300);
 71 nextWindow . s e t D e f a u l t C l o s eO p e r a t i o n (JFrame . EXIT ON CLOSE ) ;
72 nextWindow . s e t L o c a t i o n R e l a t i v e T o ( nu l l ) ;
73 nextWindow . s e t V i s i b l e ( true ) ;
                                      {\frac{1}{4} dispose ( );
\frac{75}{76}<br>\frac{77}{78}76 } ) ;
77 // S e t L a y ou t Manager
                             \frac{7}{7} s of Eugenberg manager<br>set Layout (new GridLayout (5, 0));
79
\begin{array}{c|cc}\n 80 & // Add Panels to the Window \n\end{array}81<br>
this.add (main_menu_panel);<br>
this.add (message_panel);<br>
this.add (add_module_panel);<br>
this.add (delete_module_panel);
85 this . add (modify-module-panel);
8687 } // END OF CLASS
```
Input code title credit desc Waiting

Listing A.14: Input_code_title_credit_desc_Waiting.java

```
1 import java.awt.GridLayout;<br>2 import java.awt.event.ActionEvent;
  \begin{array}{c|c} 3 & \textbf{import java.awt.event} \end{array} : \begin{array}{c} \textbf{import java.awt.event} \end{array}4 import javax.swing.*;<br>5 import java.sql.CallableStatement;<br>6 import java.sql.Connection;
   7 | import java.sql.DriverManager;<br>8 | import java.sql.ResultSet;
 9 import java.sql.SQLException;<br>10 import java.sql.Statement;
11 import java.sql. Types:
\frac{12}{13}13 public class Input code_title_credit_desc_Waiting extends JFrame {<br>14 public Input_code_title_credit_desc_Waiting(String winTitle) {
                         nd state in passession is exceeding to the credit in g calculate the public Input code_title_credit_desc_Waiting (String win Title) {
\begin{array}{c} 15 \\ 16 \end{array}super(winTitle);
\begin{array}{c} 17 \\ 18 \\ 19 \end{array}\begin{tabular}{ll} $18$ & // Create Swing Components \\ 19 & JPanel message.panel = new JPanel(); \\ 20 & final JLabel message\_panel = new JLabel("Input your Data"); \\ message.panel.add(messagelab);\quad JPanel code.panel = new JPanel(); \\ 23 & JLabel code_label = new JLabel("code"); \\ \end{tabular}\begin{array}{lll} 24 & \text{final JTextField codeTxt} = \text{new JTextField}(30); \ 25 & \text{code\_panel.add(code\_label)}; \end{array}\begin{tabular}{ll} 26 & code\_panel.add(coder{\tt Cat}); \\ 27 & JPanel title\_panel = new JPanel(); \\ 28 & JLabel title\_label = new JLabel("title"); \\ 29 & final JTextField title{\tt Txt = new JTextField(30);\newline} 30 & title\_panel.add(title\_label);\newline 31 & title\_panel.add(title{\tt Txt}); \\ 32 & JPanel credit\_panel = new JPanel(); \\ 33 & JLabel credit\_label = new JLabel("credit"); \\ \end{tabular}34<br>
final JTextField creditTxt = new JTextField(30);<br>
credit_panel.add(creditTxt);<br>
credit_panel.add(creditTxt);<br>
JPanel desc_panel = new JPanel();<br>
JLabel desc_label = new JLabel("desc");<br>
final JTextField descTxt = new J
 40 desc_panel.add (desc_label);<br>41 desc_panel.add (descTxt);
```

```
\begin{array}{lll} \text{42} & \text{JPanel input panel = new JPanel();} \end{array}<br>43 JButton input = new JButton("Input");
                             input\_panel . add (input);
\frac{44}{45}<br>\frac{46}{46}46 // Add Action to the button<br>47 input.addActionListener (new
47 input add ActionListener (new ActionListener ()<br>application action Performed (ActionEvent e
 48<br>
public void action Performed (Action Event e) {<br>
$String code = codeTxt.getText();<br>
$String title = titleTxt.getText();<br>
$String create = reditTxt.getText();<br>
$String desc = descTxt.getText();\frac{53}{54}54 i f ( ( codeTxt . g etT ext ( ) . isEmpty ( ) ) | | ( t i t l e T x t . g etT ext ( ) . isEmpty ( ) )
55 | | ( c r e d i t T x t . g etT ext ( ) . isEmpty ( ) ) | | ( d es cTxt . g etT ext ( ) . isEmpty ( ) ) ) {
56 I n p u t c o d e t i t l e c r e d i t d e s c W a i t i n g E r r o r nextErrWindow = new
57 I n p u t c o d e t i t l e c r e d i t d e s c W a i t i n g E r r o r ( " I n p u t c o d e
58 t i t l e c r e d i t d e s c W a i t i n g E r r o r " ,
59 " I n v a l i d Input . P l e a s e Try Again " ) ;
60 nextErrWindow . s e t S i z e ( 4 0 0 , 3 0 0 ) ;
 61 nextErrWindow . s e t D e f a u l t C l o s eO p e r a t i o n (JFrame . EXIT ON CLOSE ) ;
62 nextErrWindow . s e t L o c a t i o n R e l a t i v e T o ( nu l l ) ;
63 | nextErrWindow . set Visible (true);
64 dispose ();
 65 }
65<br>66<br>67\begin{array}{c|c} 67 \ 68 \end{array} else
67 else {<br>
Create_Module_Ready nextWindow = new<br>
Create_Module_Ready ("Create_Modu<br>
credit, desc);
 69<br>
Create_Module_Ready("Create_Module_Ready", code, title,<br>
70<br>
nextWindow.setSize(400, 300);<br>
nextWindow.setDefaultCloseOperation(JFrame.EXIT_ON_CLOSE);<br>
nextWindow.setLocationRelativeTo(null);
74<br>
74<br>
77 nextWindow . setLocation Relative<br>
77 nextWindow . setVisible (true);<br>
77 is pose ();
                                                             \begin{array}{c} \texttt{dispose}() \, ; \ \texttt{\emph{b}} \end{array}\begin{array}{c|c}\n 76 \\
 77 \\
 78\n\end{array} } }
77<br>78<br>79<br>80\begin{array}{c|c} 80 & // Set \; Layout \; Manager \\ 81 & set \; Lavout (new \; GridLaw) \end{array}\text{setLayout}(\text{new GridLayout}(6, 0));\frac{01}{82}\begin{array}{c|c} 83 \\ 83 \\ 84 \end{array} // Add Panels to the Window<br>
this . add (message_panel);
84 this . add (message_panel);<br>85 this . add (code_panel);
 85 this . add ( code_panel );<br>86 this . add ( title_panel );
 87 this . add ( credit_panel );<br>88 this . add ( desc_panel );<br>89 this . add ( input_panel );
 90 }
91 } // END OF CLASS
```
Input code title credit desc Waiting Error

Listing A.15: Input_code_title_credit_desc_Waiting_Error.java

```
import java.awt. GridLayout;
 2 import java.awt.event.ActionEvent ;<br>3 import java.awt.event.ActionLister
 3 import java. awt. event. ActionListener;<br>4 import javax. swing. *;<br>5 import iava. sgl. CallableStatement:
  4 import javax.swing.*;<br>5 import java.sql.CallableStatement;<br>6 import java.sql.Connection;
  7 import java.sql.DriverManager;<br>8 import java.sql.ResultSet;
9 import java.sql.SQLException;<br>
10 import java.sql.Statement;<br>
11 import java.sql.Types;
12
13 public class Input_code_title_credit_desc_Waiting_Error extends JFrame {<br>14 public Input_code_title_credit_desc_Waiting_Error(String winTitle, String err) {
\frac{15}{16}super( win Title );\frac{17}{18}\begin{array}{c|c}\n 18 \\
 19\n \end{array} // Create Swing Components<br>
19<br>
JPanel input_code_title_cre
19 JPanel input_code_title_credit_desc_waiting_error_<br>20 warning_panel = new JPanel ();<br>final JLabel input_code_title_credit_desc_waiting_error
22 warningLab = new JLabel ();<br>23 input_code_title_credit_desc_waiting_error_warning_panel
24 . add (input_code_title_credit_desc_waiting_error_warningLab);
\frac{25}{26}<br>\frac{26}{27}{\bf JPanel \; message\_panel \; = \; new \; JPanel} \;(\;)\begin{tabular}{c} 27 \\ 28 \\ 29 \\ 30 \\ 31 \\ 32 \\ 33 \\ 34 \\ 35 \\ 36 \\ 37 \\ 38 \\ 39 \\ 30 \\ 31 \\ 321 \\ 332 \\ 342 \\ 35 \\ 36 \\ 37 \\ 38 \\ 39 \\ 30 \\ 31 \\ 321 \\ 332 \\ 343 \\ 35 \\ 36 \\ 37 \\ 38 \\ 39 \\ 30 \\ 30 \\ 31 \\ 321 \\ 332 \\ 343 \\ 35 \\ 36 \\ 37 \\ 38 \\ 39 \\ 30 \\ 30 \\ 31 \\ 32 \\ 32 \\ 33 \\ 34\frac{32}{33}\frac{1}{4} Add Action to the button
```

```
34 initialise.addActionListener(new ActionListener(){<br>35 public void actionPerformed(ActionEvent e){
\begin{array}{c|c} 36 & \text{Input-code.title_credit-desc. } \text{Waiting backWindow = new} \\ 37 & \text{Input-code.title.createit-desc. } \\ 38 & \text{Waiting "Naiting" } \text{[input_code.title: credit-desc. } \text{Waiting" } \text{);} \\ 39 & \text{backWindow.setSize (400, 300);} \end{array}40 backWindow . setDefaultCloseOperation (JFrame .EXIT_ON_CLOSE);<br>
backWindow . setLocation RelativeTo (null );<br>
backWindow . setVisible (true );
\begin{array}{c|c}\n 43 \\
 44\n \end{array} dispose ();
\begin{array}{c|c}\n 44 & & \rightarrow \\
 45 & & \rightarrow \\
 \end{array}\frac{46}{47}47 // Set Layout Manager<br>48 setLayout(new GridLayout(3, 0));
4950 // Add Panels to the Window<br>
this . add(input_code_title_ci<br>
this . add(message_panel);
51 this .add (input_code_title_credit_desc_waiting_error_warning_panel);<br>52 this .add ( message_panel );
53 this . add (initialise_panel);
\begin{array}{c} 54 \\ 55 \end{array} } // END OF CLASS
```
Input code credit Waiting


```
import java.awt.GridLayout;
    2 import java.awt.event.ActionEvent;<br>3 import java.awt.event.ActionListener;
    4 import javax.swing.*;<br>5 import java.sql.CallableStatement;<br>6 import java.sql.Connection;
   7 import java.sql.DriverManager;<br>8 import java.sql.ResultSet:
   \begin{array}{c|c} 8 & \text{import java.sql. ResultSet} \\ 9 & \text{import java sell. SOLEvent} \end{array}9 import java.sql. SQLException;<br>10 import java.sql. Statement;
11 import java.sql. Types;
\frac{12}{12}13 public class Input_code_credit_Waiting extends JFrame {<br>14 public Input_code_credit_Waiting(String winTitle) {
\frac{15}{16}super(winTitle);\begin{array}{c} 17 \\ 18 \end{array}18 18 // Create Swing Components<br>19 JPanel message_panel = new JPanel();<br>20 final JLabel messageLab = new JLabel("Input your Data");
 \begin{array}{ll} \texttt{21} & \texttt{message\_panel}.\ \texttt{add}(\texttt{message\_Lab}); \ \texttt{32} & \texttt{JPanel}(\texttt{33}) \ \texttt{33} & \texttt{JLabel}~ \texttt{code\_panel} = \texttt{new} \ \texttt{JLabel}(\texttt{"code"}); \ \texttt{23} & \texttt{final} \ \texttt{JTextField}~ \texttt{code}^\mathsf{T} \ \texttt{xt} = \texttt{new} \ \texttt{JTextField}(\texttt{30}); \ \texttt{code\_panel}.\ \texttt{add}(\texttt{code\_Path}); \ \texttt{code\_panel}.\ \texttt{add}(\texttt{code\begin{tabular}{ll} \bf 27 & \bf JPanel credit\_panel = new JPanel(); \\ \bf 28 & \bf ILabel credit\_label = new JLabel("credit"); \\ \bf 29 & \bf final JTextField credit\_new JTextField(30); \\ \bf 30 & \bf credit\_panel.add(\textit{credit\_label}); \\ \bf 31 & \bf credit\_panel.add(\textit{credit.txt}); \\ \bf JPanel input = new JPanel(); \\ \bf 32 & \bf JButton input = new JButton("Input"); \\ \end{tabular}\begin{array}{c|c} \text{34} & \text{52} \text{24} \\ \text{534} & \text{534} \\ \text{534} & \text{534} \end{array}\begin{array}{c} 35 \\ 36 \end{array}\begin{array}{c|c|c|c|c} \hline 36 & // Add Action to the button \\ \hline 37 & \text{input.addActionListener( new)} \end{array}37<br>
\left.\begin{matrix} \text{input.addActionListener} \ (\text{ new ActionListener}) \ \ \{ \\ \text{public void actionPerformed(ActionEvent e) } \ \ \{ \\ \text{String code = codif x.getText();} \\ \text{String credit = creditText.getText();} \end{matrix}\right.\frac{40}{41}42 if ((codeTxt.getText().isEmpty()) || (creditTxt<br>43 .getText().isEmpty())) {
 \begin{array}{c} \text{Input-code-credit-Waiting-Error nextErrWindow = new \\ \text{Input-code-credit-Waiting-Error {\color{red} {\texttt{Pror}}}, \text{Input-code-credit-LWaiting-Error("Input-code-credit-)} \\ \text{Matrix} & \text{Matrix} & \text{Input}.\text{ Please Try Again" } \text{)}; \\ \text{nextErrWindow.setSize (400, 300)} \\ \text{nextErwindow.setDecimal (JFrame. EXIT-ON.CLOSE) } \\ \text{nextErwindow.setLocationRelativeTo (null) }; \end{array}50 nextErrWindow.set Visible (true);<br>
dispose ():
                                                                                 dispose();
 52 }
\begin{array}{c} 52 \\ 53 \\ 54 \\ 55 \end{array}\begin{array}{ll} \texttt{54} & \texttt{else} \; \{ \\ 56 & \texttt{Read}.Module\_Code\_Credit\_Ready \; \texttt{nextWindow} = \texttt{new} \\ 56 & \texttt{Read}.Module\_Code\_Credit\_Read\_Module\_Code\_ \\ 57 & \texttt{Credit}.\texttt{Ready", code, credit}); \\ 58 & \texttt{nextWindow}.\texttt{setSize(400, 300)}; \\ 60 & \texttt{nextWindow}.\texttt{setLocationRelativeTo} (\texttt{JFrame}.\texttt{EXT\_ON\_CLOSE}); \end{array}60 nextWindow . setLocationRelativeTo(null);<br>
61 nextWindow . setVisible(true);
```
 $\begin{array}{c|c} 62 & \text{dispose} (); \end{array}$ 63 } $\begin{bmatrix} 64 \\ 65 \end{bmatrix}$ } } \rightarrow) ; 66 67 // Set Layout Manager
68 setLayout (new GridLay $\text{setLayout}(\text{new GridLayout}(4, 0));$ 69
70 70
 71
 72
 73
 74
 75
 76
 78
 79
 71
 78
 79
 71
 78
 79
 71 this . add (message_panel);

73 this . add (code_panel);

73 this . add (credit_panel); 74 **this** . add (in put panel);

76 $\left\{\n\begin{array}{ccc}\n\text{this} & \text{add (input panel)} \\
\text{?76}\n\end{array}\n\right\}$ $\begin{array}{c} 75 \ 76 \end{array}$ } // END OF CLASS

Input code credit Waiting Error

Listing A.17: Input_code_credit_Waiting_Error.java

```
1 import java.awt.GridLayout;<br>2 import java.awt.event.ActionEvent;<br>3 import java.awt.event.ActionListener;
   4 import javax.swing.*;<br>5 import java.sql.CallableStatement;<br>6 import java.sql.Connection;
 7 import java.sql.DriverManager;<br>8 import java.sql.ResultSet;<br>9 import java.sql.SQLException;<br>10 import java.sql.Statement;
11 import java.sql.Types;
\frac{12}{13}13 | public class Input_code_credit_Waiting_Error extends JFrame {<br>14 |       public Input_code_credit_Waiting_Error(String winTitle, String err) {
\frac{14}{15}<br>\frac{15}{16}super(winTitle);\begin{array}{c} 17 \\ 18 \\ 19 \end{array}18 |/ Create Swing Components<br>19 | JPanel input_code_credit_waiting_error_warning_panel = new JPanel();
 20 final JLabel input_code_credit_waiting_error_warningLab = new JLabel ();<br>21 input_code_credit_waiting_error_warning_panel
 add(input_code_credit_waiting_error_warningLab);<br>
1 Danel message_panel = new JPanel();<br>
1 Inal JLabel messageLab = new JLabel(err);<br>
25 message_panel.add(messageLab);<br>
1 Danel initialise_panel = new JPanel();<br>
1 Danel ini
\frac{29}{30}<br>30<br>31
 \begin{array}{ll} \textbf{[3]} & // \textit{Add Action to the button} \\ \textbf{[3]} & \textit{initialise. addAction} \textit{Action} \\ \textbf{[3]} & \textit{public void action} \textit{BeforeInmed}(\textit{ActionEvent e}) \{ \\ \textbf{[3]} & \textit{Input-code.create} \textit{it}(\textit{NationalEvent e}) \} \\ \textbf{[3]} & \textit{Input-code.create} \textit{it}(\textit{Waiting makeWith} \textit{Window} = \textit{new} \\ & \textit{Input-code.create} \textit{it}(\textit{Nonallow = new} \textit{Window.setSize (400, 300)}); \\ & \textit{backWindow.setDefaultClose\begin{array}{c|c}\n 39 \\
 40\n \end{array} dispose ();
                                         40 }
                                \} ) ;
41<br>42<br>4343 // Set Layout Manager<br>44 setLayout (new GridLayout (3, 0));
\frac{45}{46}46 // Add Panels to the Window<br>47 this. add (input_code_credit_
47<br>
this . add (input code credit waiting error warning panel);<br>
this . add (message panel) :
 48 this . add (message_panel);<br>
49 this . add (initialise_panel);
 \begin{array}{c} 50 \\ 51 \end{array} } // END OF CLASS
```
Input code Waiting

Listing A.18: $Input_code_Waiting.java$

```
\begin{array}{c|c|c|c|c} 1 & \text{import java.awt} \end{array} import java awt event Action
```

```
2 import java.awt.event.ActionEvent;<br>3 import java.awt.event.ActionLister
3<br>3 import java.awt.event.ActionListener;<br>4 import java.swing.*;<br>5 import iava.sgl.CallableStatement:
```

```
4 import javax.swing.*;<br>5 import java.sql.CallableStatement;
```

```
6 import java.sql.Connection;<br>7 import java.sql.DriverManager;
  8 import java.sql.ResultSet;<br>9 import java.sql.SQLException;
10 import java.sql.Statement;<br>11 import java.sql.Types;
\begin{array}{c} 12 \\ 13 \end{array}13 public class Input_code_Waiting extends JFrame {<br>14 public Input_code_Waiting (String winTitle) {
\frac{15}{16}super(winTitle):\frac{17}{18}18 \begin{array}{c|c} 18 \end{array} // Create Swing Components<br>19 JPanel message-panel = new<br>20 final JLabel messageLab = 1
19 JPanel message_panel = new JPanel();<br>20 final JLabel messageLab = new JLabel("Input your Data");
21 message_panel . add (messageLab);<br>22 message_panel . add (messageLab);
22 JPanel code<sub>-P</sub>anel = new JPanel();<br>23 JLabel code<sub>-</sub>label = new JLabel("code");<br>24 final JTextField codeTxt = new JTextField(30);
\begin{array}{ll} 25 \ 26 \ 27 \ 28 \ 29 \ 21 \ 22 \ 23 \ 24 \ 25 \ 26 \ 27 \ 28 \ 28 \ 29 \ 20 \ 21 \ 22 \ 23 \ 24 \ 25 \ 26 \ 27 \ 28 \ 29 \ 20 \ 21 \ 21 \ 22 \ 23 \ 24 \ 25 \ 26 \ 27 \ 28 \ 29 \ 20 \ 21 \ 21 \ 22 \ 23 \ 24 \ 25 \ 26 \ 27 \ 28 \ 29 \ 20 \ 21 \ 22 \ 23 \ 24 \ 25 \ \frac{25}{30}<br>\frac{31}{32}31 // Add Action to the button<br>
32 input addActionListener (new ActionListener () {<br>
34 public void actionPerformed (ActionEvent e) {<br>
34 String code = codeTxt .getText ();
\frac{35}{36}36 if ((\text{codeTxt}.\text{getText}~().\text{isEmpty}~()) {<br>37 Input_code_Waiting_Error nextErrWindow = new
38 Input_code_Waiting_Error ("Input_code_Waiting_Error",<br>39 mextErrWindow .setSize (400, 300);<br>nextErrWindow .setSize (400, 300);
41 nextErrWindow . setDefaultCloseOperation (JFrame .EXIT_ON_CLOSE ) ;<br>nextErrWindow . setLocationRelativeTo (null ) ;<br>nextErrWindow . setVisible (true ) ;
\begin{array}{c} 44 \\ 45 \end{array} dispose ();
45 }
\frac{46}{47}\begin{array}{c|c}\n 47 \\
 48\n\end{array} e l s e \begin{array}{c}\n 8\n\end{array}48 Read_Desc_Ready nextWindow = new Read_Desc_Ready ("Read_Desc")<br>
49 Ready", code);<br>
50 nextWindow .setSize (400, 300);
51 nextWindow .setDefaultCloseOperation (JFrame .EXIT_ON_CLOSE);<br>52 nextWindow .setUocationRelativeTo (null );<br>nextWindow .setVisible (true );
\begin{bmatrix} 54 \\ 55 \\ 50 \end{bmatrix} dispose ();
55 }
\begin{bmatrix} 56 \\ 57 \end{bmatrix} } }
                           \rightarrow );
\frac{58}{59}\begin{array}{c|c} 59 & // Set \; Layout \; Manager \end{array}<br>
60 & \text{setLayout(new GridLayout}(3, 0));\frac{61}{62}\left\{\begin{matrix}62 \ 63 \end{matrix}\right\} \left\{\begin{matrix}7 \ 4 \text{dd} \quad P\text{anels to the Window} \\ \text{this.add(message\_panel)} \end{matrix}\right\}64 this . add ( code_panel );<br>
this . add ( input_panel );
\begin{array}{c} 66 \\ 67 \end{array} } // END OF CLASS
```
Input code Waiting Error


```
\begin{array}{c|c} 1 & \text{import java.awt.}\end{array} import java.awt.event. Action
   2 import java.awt.event.ActionEvent;<br>3 import java.awt.event.ActionListener;
   \footnotesize \begin{array}{c} 4 \quad \text{import java.swing.} \pm \footnotesize \text{import java.sql. CallableStatement;} \ 6 \quad \text{import java.sql. Consection;} \ 7 \quad \text{import java.sql. DriverManager;} \ 7 \quad \text{import java.sql. ResultSet;} \ 8 \quad \text{import java.sql. SQLException;} \end{array}10 import java.sql.Statement;<br>11 import java.sql.Types;
\begin{array}{c} 12 \\ 13 \end{array}13 public class Input code Waiting Error extends JFrame {<br>14 public Input code Waiting Error (String win Title, }<br>15 super (win Title);
 14 public Input_code_Waiting_Error(String winTitle, String err) {<br>15 super(winTitle);
\frac{16}{17}17 // Create Swing Components
 18 JPanel input_code_waiting_error_warning_panel = new JPanel();<br>19 final JLabel input_code_waiting_error_warningLab = new JLabel();
 20 in put_code_waiting_error_warning_panel<br>21 .add(input_code_waiting_error_warningLab);
```

```
22 \n\Big\} JPanel message_panel = new JPanel();<br>23 final JLabel messageLab = new JLabel(err);
 24<br>
message_panel.add(messageLab);<br>
19 JPanel initialise_panel = new JPanel();<br>
36<br>
JButton initialise = new JButton("initialise");<br>
27 initialise_panel.add(initialise);
\begin{array}{c} 28 \\ 29 \\ 30 \end{array}// Add Action to the button<br>initialise.addActionListener(new ActionListener() {
 \begin{tabular}{c} 30 \\ 31 \\ 32 \\ 33 \\ 34 \\ 35 \\ 36 \\ 37 \\ 38 \\ 39 \\ 30 \\ 34 \\ 35 \\ 36 \\ 37 \\ 38 \\ 39 \\ 30 \\ 31 \\ 34 \\ 35 \\ 36 \\ 37 \\ 38 \\ 39 \\ 30 \\ 31 \\ 34 \\ 35 \\ 36 \\ 37 \\ 38 \\ 39 \\ 30 \\ 30 \\ 31 \\ 34 \\ 35 \\ 36 \\ 37 \\ 38 \\ 39 \\ 30 \\ 30 \\ 31 \\ 34 \\ 35 \\ 36 \\ 37 \\ 38 \\ 39 \\ 30 \\ 30 \\39 }
\begin{array}{c|c}\n 40 & \rightarrow \\
 \end{array}\frac{41}{42}42 // Set Layout Manager<br>42 setLayout (pow GridLay
                              \text{setLayout}(\text{new GridLayout}(3, 0));\frac{44}{45}45 // Add Panels to the Window<br>46 this add (input_code_waiting_error_warning_panel);
 47 this . add (message_panel);<br>
48 this . add (initialise_panel);
 49 }
\begin{array}{c} 150 \\ 50 \end{array} \begin{array}{c} \uparrow \\ \downarrow \end{array} // END OF CLASS
```
Create Module Ready


```
1 import java.awt.GridLayout;<br>2 import java.awt.event.Actio
   2 | import java.awt.event.ActionEvent;<br>3 | import java.awt.event.ActionListener;
   4 import javax swing *;<br>5 import java sql. CallableStatement;<br>7 import java sql. Connection;<br>7 import java sql. DriverManager;
   8 import java.sql.ResultSet;<br>9 import java.sql.SQLException;
 10 import java.sql.Statement;<br>11 import java.sql.Types;
\begin{array}{c} 12 \\ 13 \end{array}13 public class Create Module Ready extends JFrame {<br>14 public Create Module Ready (String win Title, final String code,<br>15 final String title, final String credit, final String desc) {<br>super(win Title);
\frac{17}{18}18 // Create Swing Components<br>
19 JPanel code<sub>-</sub>panel = new JPanel ();<br>
21 code_panel add (codeLab );<br>
22 JPanel title<sub>-</sub>panel = new JPanel ();
 134<br>
134<br>
134<br>
14 itile. panel .add (itileLab );<br>
159<br>
17 Panel credit.panel = new JPanel ();<br>
17 irredit.panel = new JPanel ();<br>
17 credit.panel creditLab = new JLabel ();<br>
17 Panel desc.panel = new JPanel ();
\frac{37}{38}\begin{array}{c|c|c|c|c} \hline 39 & // Add Action to the button \cr \hline \end{array}40 create.addActionListener( new ActionListener() {<br>41 public void actionPerformed(ActionEvent e) {<br>boolean recordInserted = false;
\frac{43}{44}\frac{44}{45} // Declare Connection properties<br>
Connection conn = null;<br>
Statement stmt = null;<br>
int recAffected = 0;
\frac{47}{48}49 // Declare Connection properties<br>50 try {
 50 \hbox{try }\{ \atop 51\ }<br>
\hbox{Class. forName ("com.mysql.jdbc. Driver"). newInstance ();\n52\ } \hbox{String connectionUrl = "jdbc:mysql://localhost:3306/testdbl";\n53\ } \hbox{String connectionUser = "root";\n54\ }
```
55 conn = **DriverManager** .getConnection (connection Url ,
56 connection User , connection Password) ;
stmt = conn .createStatement () ; $\begin{array}{c} 57 \\ 58 \\ 59 \end{array}$ 59 *// Execute a query statement or a stored procedure*
60 **CallableStatement** procnone =
61 **Connectial Connective Call createModule**(?, ?, ?, ?)}"); 62
 63 $\begin{array}{c|c|c|c|c|c} 63 & \times & \times & \times & \times & \textit{Sets} & \textit{the input parameter} \\ 64 & & \textit{procnone} & \textit{setString(2, title)} \end{array}$ 64

From the set String (4, desc);

prochable desc);

prochable description to the set int (1, Integer parseInt (code));

prochable setInt (3, Integer parseInt (credit)); 68
 69 $//$ Registers the out parameters $r \cdot r \cdot r$ recaffected = procnone.executeUpdate(); 70
 71
 72
 73 // Check if there is an updated record if (recAffected > 0) 73 if (recAffected > 0)
74 recordInserted = **true**;
75 else $r \in \text{condInverted} = \text{false};$ $\begin{array}{r} 74 \\ 75 \\ 76 \\ 77 \\ 78 \\ 79 \end{array}$ 78 } catch (Exception e1) {
 $\begin{array}{c} \text{et.} \text{er } \text{et.} \text{$ $\begin{tabular}{c|c|c|c} 80 & & & & & & & & \\ 81 & & & & & & & & \\ 82 & & & & & & & & & \\ 83 & & & & & & & & & \\ 84 & & & & & & & & & \\ 85 & & & & & & & & & \\ 86 & & & & & & & & & \\ 87 & & & & & & & & & & \\ 88 & & & & & & & & & \\ 89 & & & & & & & & & & \\ 84 & & & & & & & & & & \\ 86 & & & & & & & & & & \\ 87 & & & & & & & & & & \\ 89 & & & & & & & & & & & \\ 80 & & & & & & & & & & & \\ 81 &$ 85

86 if (recordInserted) {

36 if (recordInserted) {

36 system.exit(0);

89 System.exit(0); 90 } $\begin{array}{c|c}\n 90 & & \rightarrow \\
 91 & & \text{else} \n\end{array}$ 92 Create_Module_Ready_Error nextErrWindow = new

93 Create_Module_Ready_Error ("Create_Module 93

Create_Module_Ready_Error ("Create_Module_Ready_Error ",

"Can not be Created. Please Try Again");

nextErrWindow.setDefaultCloseOperation (**JFrame**.EXIT_ON_CLOSE);

nextErrWindow.setDefaultCloseOperation (**JFrame**.EXI 97 nextErrWindow setLocationRelativeTo(null);
98 nextErrWindow setVisible(true): 98 nextErrWindow . s et Visible (true);
 $\begin{array}{c} \n\text{99} \\
\text{90}\n\end{array}$ $\begin{array}{c|c} 99 \\ 100 \end{array}$ dispose (); 100 } $\begin{array}{c|c}\n 101 \\
 102\n \end{array}$ } λ) ; 103 104 *// Set Layout Manager*
105 setLayout (new GridLayout (6, 0)); $\frac{105}{106}$ 107 // Add Panels to the Window

108 **this.add** (code_panel);

109 **this.add** (title_panel);
 this.add (credit_panel);
 this.add (desc_panel);
 this.add (message_panel); 113 \vert this . add (create_panel); $\frac{114}{115}$ 1 // END OF CLASS

Create Module Ready Error


```
\begin{tabular}{c} 23 \\ 24 \\ 25 \\ 26 \\ 27 \\ 28 \\ 29 \\ 20 \\ 20 \\ 21 \\ 22 \\ 23 \\ 24 \\ 25 \\ 26 \\ 27 \\ 28 \\ 29 \\ 20 \\ 21 \\ 22 \\ 23 \\ 24 \\ 25 \\ 26 \\ 27 \\ 28 \\ 29 \\ 20 \\ 21 \\ 23 \\ 24 \\ 25 \\ 26 \\ 27 \\ 28 \\ 29 \\ 20 \\ 20 \\ 21 \\ 23 \\ 24 \\ 25 \\ 26 \\ 27 \\ 28 \\ 29 \\ 20 \\ 20 \\ 21 \\ 23 \\ 24 \\ 25 \\\frac{28}{29}\begin{tabular}{ll} $29$ & // Add Action to the button \newline exit.addActionListener() {\newline 31 & public void actionPerformed (ActionEvent e) {\newline 32 & Input_code.title.create() {\newline 33 & Input_code.title.createdict-desc-Waiting backWindow = new {\newline 34 & Input_code.title.createdict-desc-Waiting("Input-code.title.createdict-desc-Waiting");\newline 35 & backWindow.setSize(400, 300);\newline backWindow.setDefaultCloseOperation (JFrame.EXTTON.CLOSE);\n    backWindow.setLocationRelativeTo (null);\n    backWindow.setLocationRelativeTo (null);\n    backWindow.setVisible (true);\n    classpose(); \end{tabular}\begin{array}{c} 39 \\ 40 \end{array} } }
40<br>41<br>42<br>4342 // Set Layout Manager<br>43 setLayout (new GridLayout (3, 0));
44
45 // Add Panels to the Window<br>46 this. add (create_module_read
 46 this . add ( create_module_ready_error_warning_panel );<br>47 this . add ( message_panel ) ;
\begin{array}{c|c} 48 & \text{this.add} (exit\_panel) ; \end{array}\begin{array}{c} 49 \\ 50 \end{array} } // END OF CLASS
```
Delete Module Ready


```
import java.awt.GridLayout;
 2 import java.awt.event.ActionEvent;<br>3 import java.awt.event.ActionLister<br>4 import javax.swing.*;
         import java.awt.event.ActionListener;
 4 import javax.swing.*;<br>5 import java.sql.Calla<br>6 import java.sql.Conn<br>7 import java.sql.Drive
  5 import java.sql.CallableStatement;<br>6 import java.sql.Connection;<br>7 import java.sql.DriverManager;
 \begin{array}{c|c|c|c|c} 8 & \text{import java} & \text{sql} & \text{ResultSet} \\ 9 & \text{import java} & \text{sol} & \text{SOLEvent} \end{array}9 import java.sql. SQLException;<br>10 import java.sql. Statement;
11 import java \cdot sq1. Types;
\frac{12}{13}13 public class Delete_Module_Ready extends JFrame {<br>14 public Delete_Module_Ready(String winTitle, final String code) {
15 super ( win Title );
\frac{16}{17}\begin{array}{c|c} 17 \\ 17 \\ 18 \end{array} // Create Swing Components
18 JPanel code_panel = new JPanel ();<br>19 final JLabel codeLab = new JLabel ();<br>code_panel add (codeLab);
21 JPanel message_panel = new JPanel();<br>22 final JLabel messageLab = new JLabel("Are you sure to proceed Deleting ");<br>message_panel.add(messageLab);
24 JPanel delete_panel = new JPanel ();<br>25 JButton delete = new JButton("Delete");<br>delete_panel.add(delete);
\frac{26}{27}<br>28
28 // Add Action to the button<br>29 delete.addActionListener (new ActionListener () {<br>30 public void actionPerformed(ActionEvent e) {
31 boolean recordDeleted = false;
\frac{32}{33}\begin{array}{c|c|c|c|c} \text{33} & \text{72} & \text{53} & \text{64} \ \text{34} & \text{75} & \text{66} & \text{67} & \text{78} & \text{79} & \text{88} \ \text{50} & \text{70} & \text{70} & \text{70} & \text{88} & \text{70} & \text{70} & \text{88} & \text{70} & \text34 Connection conn = null;<br>35 Statement stmt = null;<br>int recAffected = 0;
36<br>37<br>38<br>39\begin{array}{ll} &// \textit{Declare Connection properties} \\ \text{try} \left\{ \begin{array}{ll} &\text{Class. forName("com.mysql.jdbc. Driver"). newInstance();\\ 41 & \text{Class. forName("com.mysql.jdbc. Driver"). newInstance();\\ &\text{String connection Usr = "root"};\\ &\text{String connectionPassword = "";} \end{array} \right. \end{array}44 conn = DriverManager .getConnection (connection Url , connection User ,<br>
connection Password ) ;
44<br>
46 stemmer = DriverManager.getConne<br>
5 stmt = conn.createStatement ();<br>
46 stmt = conn.createStatement ();
\frac{47}{48}48 // Execute a query statement or a stored procedure<br>49 CallableStatement procnone =<br>50 conn.prepareCall("{call deleteModule(?}}");
\frac{51}{52}52 // Sets the input parameter<br>53 procnone.setInt(1, Integer.parseInt(code));
\frac{54}{55}\text{recAffected} = procnone. executeUpdate();
```

```
\frac{56}{57}57 // Check if there is an updated record<br>58 if (recAffected > 0)<br>59 cordDeleted = true;
\begin{bmatrix} 60 \\ 61 \end{bmatrix} else
                                                            recordDelete d = false;62 \overline{63}63 } catch (Exception e1) {<br>
e1.printStackTrace();
65 } f i n a l l y {
66 try { i f ( stmt != nu l l ) stmt . c l o s e ( ) ; } catch (SQLException e1 )
67 { e1 . p r i n t S t a c k T r a c e ( ) ; }
68 try { i f ( conn != nu l l ) conn . c l o s e ( ) ; } catch (SQLException e1 )
order the contracteristic equality of the set of \begin{array}{c} \text{if } (\text{conn } != \text{ null}) \text{ conn } \text{ close}(); \\ \text{if } (\text{conn } != \text{ null}) \text{ conn } \text{ close}(); \} \text{ catch } (\text{SQLException } \text{el}) \\ \text{if } \text{el.printStackTrace}(); \end{array}70 }<br>71 if (record Deleted) {
72 JOptionPane . showM essag eD ia log ( nu l l ,
73 "The b u s i n e s s t a s k i s comp l et ed s u c c e s s f u l l y . " ) ;
71 \begin{array}{ccc} \text{if (recordDeleteed)} & \{ \\ \text{72} & & \text{JOptionPane} \text{.} \ \text{.} \\ \text{73} & & \text{"The busine} \\ \text{74} & & \text{?} \\ \end{array}rac{75}{76}\begin{array}{c} 75 \\ 76 \\ 77 \\ 78 \end{array} else {
77<br>Delete_Module_Ready_Error nextErrWindow = new<br>Delete_Module_Ready_Error("Delete_Module_Ready_Error ",<br>"Can not be Deleted. Please Try Again");<br>nextErrWindow.setSize(400, 300);
81 nextErrWindow .setDefaultCloseOperation (JFrame .EXIT_ON_CLOSE ) ;<br>
nextErrWindow .setUcsationRelativeTo (null ) ;<br>
nextErrWindow .setVisible (true ) ;
\begin{array}{c|c}\n 84 \\
 85 \\
 \hline\n 85\n \end{array} dispose ();
85 }
\begin{array}{c|c}\n 86 & & \rightarrow \\
 87 & & \rightarrow \\
 \end{array}rac{88}{89}89 <br>
90 SetLayout (new GridLayout (3, 0));
\begin{array}{c} 91 \\ 92 \end{array}\begin{array}{c|cc}\n 92 & // Add Panels to the Window\n\end{array}<br>this add (code panel):
93 this.add(code_panel);<br>94 this.add(message_panel);
95 this . add ( delete<sub>-panel</sub> );
96 }
97 } // END OF CLASS
```
Delete Module Ready Error


```
1 import java.awt.GridLayout;<br>2 import java.awt.event.ActionEvent;
  \begin{array}{c|c} 3 & \textbf{import java.awt.event.ActionListener} \\ 4 & \textbf{import javax.swing.*}; \end{array}4 import javax.swing.*;<br>5 import java.sql.CallableStatement;<br>6 import java.sql.Connection;
 7 import java.sql.DriverManager;<br>8 import java.sql.ResultSet;<br>9 import java.sql.SQLException;<br>10 import java.sql.Statement;
11 | import java.sql.Types;
\frac{12}{13}13 public class Delete_Module_Ready_Error extends JFrame {<br>14 public Delete_Module_Ready_Error(String winTitle, S
14 public Delete_Module_Ready_Error (String winTitle, String err) {<br>15 super(winTitle):
                                  super(winTitle);\frac{16}{17}17 // Create Swing Components<br>18 JPanel delete_module_ready<br>19 final JLabel delete_module
 18 JPanel delete_module_ready_error_warning_panel = new JPanel();<br>19 final JLabel delete_module_ready_error_warningLab = new JLabel();
 ablette_module_ready_error_warning_panel<br>
21 add(delete_module_ready_error_warningLab);<br>
22 JPanel message_panel = new JPanel();<br>
32 final JLabel messageLab = new JLabel(err);
 24 message_panel.add(messageLab);<br>25 JPanel exit_panel = new JPanel();<br>26 JButton exit = new JButton("exit");
27 exit_panel.add(exit);
\frac{28}{29}29 // Add Action to the button<br>30 exit.addActionListener(new)<br>31 public void actionPerform
 \begin{tabular}{c|c|c} 30 & exit.\:addActionListener( \textbf{ new ActionListener() } {\small \texttt{public void actionPerformed(ActionEvent) } {\small \texttt{public void actionPerformed(ActionEvent) } {\small \texttt{public void } \texttt{input_code}.\texttt{Waiting} } } \texttt{Input} \texttt{code} \texttt{= new} } \\ 32 & \texttt{Input-code}.\texttt{Waiting('Input-code}.\texttt{Waiting''}); \\\ 33 & \texttt{backWindow}.\:set \texttt{set@4010, 300}; \\\ 34 & \texttt{backWindow}.\:set \texttt{DefaulCloseOperation(JFrame.EXT\_ON\_CLOSE)}; \\\end{tabular}36 backWindow . setLocation Relative To (null);<br>37 backWindow . setVisible (true);
\begin{array}{c|c} \text{37} \\ \text{38} \\ \text{39} \\ \text{300} \\ \text{318} \\ \text{320} \\ \text{331} \\ \text{342} \\ \text{351} \\ \text{362} \\ \text{373} \\ \text{384} \\ \text{395} \\ \text{306} \\ \text{307} \\ \text{318} \\ \text{329} \\ \text{339} \\ \text{340} \\ \text{350} \\ \text{360} \\ \text{371} \\ \text{382} \\ \text{393} \\ \text{306} \\ \text{307} \\ \text{318dispose();
\begin{bmatrix} 39 \\ 40 \end{bmatrix} }
                                  40 } ) ;
 41
```
 $\begin{array}{c|c} 42 & // Set \; Layout \; Manager \end{array}$
 $\begin{array}{c} 43 \end{array}$ setLayout (new GridLayout $(3, 0)$); $\begin{array}{c} 44 \\ 45 \end{array}$ 45 // Add Panels to the Window
46 this.add(delete_module_ready_error_warning_panel);
this.add(message_panel); $\begin{array}{c} 48 \\ 49 \end{array}$ this . add (exit -panel); $\begin{array}{c} 49 \\ 50 \end{array}$ } // END OF CLASS

Read Desc Ready

Listing A.24: Read_Desc_Ready.java

```
1 import java.awt.GridLayout;<br>2 import java.awt.event.ActionEvent;<br>3 import java.awt.event.ActionListener;
   4 import javax.swing.*;<br>5 import java.sql.CallableStatement;<br>6 import java.sql.Connection;
   7 | import java.sql.DriverManager;<br>8 | import java.sql.ResultSet;
 9 import java.sql.SQLException;<br>
10 import java.sql.Statement;<br>
11 import java.sql.Types;
\frac{12}{13}13 public class Read Desc-Ready extends JFrame {14}
                  public Read-Desc-Ready extends Strime 1 public Read-Desc-Ready (String winTitle, final String code) {
15 super (win Title):
\frac{16}{17}17 17 (Treate Swing Components)<br>18 JPanel code-panel = new JPanel ();<br>19 final JLabel codeLab = new JLabel ();
 20 code_panel.add(codeLab);<br>21 JPanel message_panel = new JPanel();<br>22 final JLabel messageLab = new JLabel("Are you sure to proceed Reading ");<br>23 message_panel.add(messageLab);
 24 JPanel read panel = new JPanel ();<br>25 JButton read = new JButton("Read");<br>read panel .add (read);
\frac{27}{28}\begin{array}{c|c} 28 & // \textit{Add Action to the button} \\ 29 & \textit{read.addActionListener( new ActionListener() } \end{array}30 public void actionPerformed(ActionEvent e) {<br>31 boolean recordReturned = false;
\frac{32}{33}33 // Declare Connection properties<br>
34 Connection conn = null;<br>
35 Statement stmt = null;<br>
ResultSet rs = null;
\frac{37}{38}\begin{array}{c|c|c|c} \hline 38 & & \end{array} \begin{array}{c} \hline \\begin{array}{ll} \textbf{139} & \textbf{try} \; \{ \\ \textbf{140} & \textbf{Class. forName("com.mysql.jdbc. Driver"). newInstance();} \\ \textbf{15} & \textbf{String connection} \; \textbf{Unres} = \texttt{``jobc:mysql: // localhost:3306/testdbl''}; \\ \textbf{16} & \textbf{String connection} \; \texttt{User} = \texttt{``root''}; \\ \textbf{17} & \textbf{String connection} \; \texttt{Password = "''}; \\ \textbf{18} & \textbf{String connection} \; \texttt{Password = "''}; \\ \end{array}44 conn = DriverManager .getConnection ( connection Url ,   connection User ,<br>45 connection Password ) ;<br>stmt = conn .createStatement ( ) ;
\frac{47}{48}48 // Execute a query statement or a stored procedure<br>49 CallableStatement procnone =<br>50 conn.prepareCall("{call readDesc(?, ?)}");
\frac{51}{52}52 // Sets the input parameter<br>53 procnone.setInt(1, Integer.parseInt(code));
\frac{53}{54}<br>\frac{54}{55}55 // Registers the out parameters<br>56 procnone.registerOutParameter(2, Types.VARCHAR);
\frac{57}{58}procnone.executeUpdate ();
\frac{59}{60}60 \begin{array}{l|l} \hline \end{array} // Declare variable(s) to stored the result(s) of the query<br>62 String descVar = procnone.getString(2);<br>62 System.out.println(descVar +" ");
63
 64 // Check if there is a returned record<br>65 if (descVar != null)<br>66 cordReturned = true;
\begin{array}{c|c} 67 & \text{else} \end{array}recordReturned = false;69
70 } catch (Exception e1)<br>
\begin{array}{ccc}\n71 \\
71 \\
72\n\end{array} } finally {
 71 e1.printStackTrace();<br>
73 finally {<br>
try { if (rs != null) rs.close(); } catch (SQLException e1)<br>
74 { if (rs != null) rs.close(); }
```

```
75 try { if (stmt != null) stmt.close(); } catch (SQLException e1)<br>76 { e1.printStackTrace(); }
 77 try { if (conn != null) conn close(); } catch (SQLException e1)<br>78 { e1.printStackTrace(); }
 79 }
 80<br>81 if (record Returned) {<br>10ption Pane.
 81 JOptionPane showMessageDialog(null,<br>82 * The business task is completed successfully.");
 \begin{array}{c|c}\n 83 & \text{Sterm. exit (0)}\n \end{array}\frac{84}{85}\begin{array}{c|c}\n 85 \\
 86\n \end{array} else {
 86 Read_Desc_Ready_Error nextErrWindow = new<br>
87 Read_Desc_Ready_Error ("Read_Desc_Read
 87<br>Read_Desc_Ready_Error ("Read_Desc_Ready_<br>Brror ", "Can not be retrieved. Please Try Again");<br>89 nextErrWindow.setDefaultCloseOperation (JFrame.EXIT_ON_CLOSE);<br>nextErrWindow.setDefaultCloseOperation (JFrame.EXIT_ON_C
 91 nextErrWindow . setLocation Relative To (\text{null});<br>92 nextErrWindow . setVisible (\text{true}):
 92 nextErrWindow . s et Visible (true);<br>93 dispose ();
\begin{array}{c|c} 93 & 93 \\ 94 & 15 \\ 95 & 1 \end{array}94 | }
 \begin{array}{c} 95 \\ 96 \end{array} } }
 97<br>98
 \begin{array}{c|c} 98 & // Set \; Layout \; Manager \cr 99 & \; setLayout(new \; GridLayout(3, 0)); \end{array}\frac{100}{101}\begin{array}{c|c|c|c|c|c} \hline 101 & & & \\ 101 & & & \text{this add (code panel)} \end{array}102 this.add(code_panel);<br>103 this.add(message_panel);
104 this . add (read_panel);
\frac{105}{106}\frac{1}{2} // END OF CLASS
```
Read Desc Ready Error

Listing A.25: Read Desc Ready Error.java

```
1 import java.awt.GridLayout;<br>
2 import java.awt.event.ActionEvent;<br>
3 import java.awt.event.ActionListener;<br>
4 import javax.swing.*;<br>
5 import java.sql.CallableStatement;<br>
6 import java.sql.Connection;
  7 | import java.sql.DriverManager;<br>8 | import java.sql.ResultSet;
9 import java.sql.SQLException;<br>
10 import java.sql.Statement;<br>
11 import java.sql.Types;
\frac{12}{13}13 public class Read_Desc_Ready_Error extends JFrame {<br>14 public Read_Desc_Ready_Error(String winTitle, String err) {
15 | super ( win Title );
16
17 17 / Create Swing Components<br>18 1 JPanel read_desc_ready_error_warning_panel = new JPanel();<br>19 1 Final JLabel read_desc_ready_error_warningLab = new JLabel();
20 read_desc_ready_error_warning_panel.add (read_desc_ready_error_warningLab);
\frac{21}{22}<br>22<br>23
22 \n\Big\} JPanel message_panel = new JPanel();<br>23 final JLabel messageLab = new JLabel(err);
24 message_panel . add (messageLab);
rac{25}{26}26 JPanel exit_panel = new JPanel();<br>27 JButton exit = new JButton("exit");<br>exit_panel.add(exit);
\frac{29}{30}30 \vert // Add Action to the button<br>31 exit.addActionListener(new ActionListener() {
32<br>
Public void actionPerformed(ActionEvent e) {<br>
Input_code_Waiting \backWindow = new<br>
\begin{array}{r} \text{Input code} - \text{Waiting} \text{ backWindow} = \text{new} \\ \text{Input code} - \text{Waiting} \text{ "Input} \text{ code} - \text{Waiting} \text{ " } \text{)}; \\ 35 \\ \text{backWindow} \text{ setSize (400, 300)}; \\ 36 \\ \text{backWindow} \text{ setLocationRelativeTo} \text{ (JFrame} \text{EXT} \text{\begin{array}{c|c} 39 & \text{dispose ()}; \\ 40 & \text{ } \end{array}40 }
\begin{array}{c|c}\n 41 & \rightarrow \\
 \end{array}\begin{array}{c} 42 \\ 43 \end{array}43 // Set Layout Manager<br>44 setLayout (new GridLayout (3, 0));
\frac{45}{46}46 // Add Panels to the Window<br>
47 this.add (read_desc_ready_eri<br>
48 this.add (message_panel);
47 this . add ( read_desc_ready_error_warning_panel );<br>
48 this . add ( message_panel ) ;
49 this . add ( exit panel );
\begin{array}{c} 50 \\ 51 \end{array} } // END OF CLASS
```
Read Module Code Credit Ready

Listing A.26: Read_Module_Code_Credit_Ready.java

```
1 import java.awt.GridLayout;<br>
2 import java.awt.event.Actio<br>
3 import java.awt.event.Actio
  2 import java.awt.event.ActionEvent;<br>3 import java.awt.event.ActionListener;
 \begin{array}{c|c} 4 & \text{import java} \\ 5 & \text{import java} \\ 6 & \text{import java} \\ 7 & \text{if } \\ 8 & \text{if } \\ 9 & \text{if } \\ 1 & \text{if } \\ 5 import java.sql.CallableStatement;<br>6 import java.sql.Connection;<br>7 import java.sql.DriverManager;
8 import java.sql.ResultSet;<br>
9 import java.sql.SQLException;<br>
10 import java.sql.Statement;<br>
11 import java.sql.Types;
\begin{array}{c} 12 \\ 13 \end{array}13 public class Read_Module_Code_Credit_Ready extends JFrame {<br>14 public Read_Module_Code_Credit_Ready(String winTitle, final String code,<br>15 final String credit) {<br>super(winTitle);
\frac{16}{17}\begin{array}{c|c} 18 & // Create Swing Components \\ 19 & J\textbf{Panel code\_panel} = \textbf{new JP} \end{array}\begin{array}{lll} 19 & \textbf{J Panel code\_panel = new JPanel();} \ 20 & \textbf{final JLabel codeLab = new JLabel();} \end{array}21 code_panel. add ( codeLab );<br>22 . Panel credit panel = new
22 JPanel credit_panel = new JPanel();<br>23 final JLabel creditLab = new JLabel();<br>credit_panel.add(creditLab);
35<br>
1993 JPanel message_panel = new JPanel();<br>
final JLabel messageLab = new JLabel("Are you sure to proceed Reading");<br>
27 message_panel .add(messageLab);<br>
1994 JPanel read_panel = new JPanel();<br>
1994 JPanel read = new JP
\frac{31}{32}32 32 // Add Action to the button<br>33 read.addActionListener(new ActionListener(){<br>34 public void actionPerformed(ActionEvent e){
\begin{array}{c|c|c|c|c} \hline \text{51} & \text{62} & \text{63} & \text{64} & \text{65} & \text{66} \\ \hline \end{array}\begin{array}{c} 36 \\ 37 \\ 38 \end{array}37 // Declare Connection properties<br>38 Connection conn = null;
\begin{array}{c|c|c} 39 & \textbf{Statement} & \text{stmt} = \texttt{null}; \end{array}<br>
\begin{array}{c|c} 40 & \textbf{ResultSet} & \text{rs} = \texttt{null}; \end{array}\frac{41}{42}\begin{array}{c|c|c|c} \hline 42 & & \end{array} // Declare Connection properties
43 degree try {<br>
44 Class .forName("com . mysql .jdbc . Driver") . newInstance ();
45<br>
String connectionUrl = "jdbc:mysql://localhost:3306/testdb1";<br>
String connectionUser = "root";<br>
String connectionPassword = ";<br>
conn = DriverManager.getConnection(connectionUrl, connectionUser,
49 connectionPassword );<br>
50 stmt = conn.createStatement ();
\frac{51}{52}\begin{array}{c|c|c|c|c|c|c|c} \hline 52 & //~Execute~a~query~statement~or~a~stored~procedure \ \hline \end{array}53 CallableStatement procnone =<br>
54 conn.prepareCall("{call readCode_Credit(?, ?)}");
^{\rm 55}_{\rm 56}56 \vert // Sets the input parameter<br>57 procnone.setInt(1, Integer.parseInt(code));
58
59 // Registers the out parameters<br>60 procnone.registerOutParameter(2, Types.INTEGER);
\frac{61}{62}pro cnone . execute Update ();
63
\begin{array}{c} \text{64} \\ \text{65} \end{array} // Declare variable(s) to stored the result(s) of the query int creditVar = procnone. getInt(2):
\begin{array}{lll} 65 & \text{int creditVar = procone.getInt (2);} \\ 66 & \text{System.out.println (creditVar +" ')} \end{array}66<br>67<br>6868 // Check if there is a returned record<br>69 if (creditVar != 0)<br>70 recordReturned = true;
\begin{bmatrix} 70 \\ 71 \\ 71 \end{bmatrix} else
                                                            recordReturned = false;72<br>73<br>7474 } catch (Exception e1) {<br>
\begin{array}{c} \text{75} \\ \text{1.} \text{ 1.} \end{array} (1) {
\begin{array}{c|c} 75 \\ 76 \\ 77 \\ 77 \end{array} e1.printStackTrace();<br>
\begin{array}{c} 27 \\ 1 \end{array} finally {<br>
\begin{array}{c} \text{if } \text{f} \text{ is } != \text{ null} \end{array}76 } finally {<br>
77 try { if (rs != null) rs.close(); } catch (SQLException e1)
78 { e1.printStackTrace(); }<br>79 try { if (stmt != null) stmt.close(); } catch (SQLException e1)
80 { e1.printStackTrace(); }<br>81 try { if (conn != null) conn.close(); } catch (SQLException e1)
\begin{array}{c|c|c|c|c|c|c} \hline \text{82} & \text{61. printStackTrace} & \text{52.} \end{array}83 }
\begin{array}{c} 84 \\ 85 \end{array}85 if (record Returned) {<br>86 Write Module
86 Write_Module_Ready nextWindow = new<br>87 Write_Module_Ready ("Write_Module_Ready", codeLab.getText(),
88 c r e d it Lab.get Text () );<br>
89 next Window.set Size (400, 300);
```

```
90 | nextWindow . setDefaultCloseOperation (JFrame .EXIT_ON_CLOSE);<br>91 | nextWindow . setLocationRelativeTo (null );
  92 nextWindow.setVisible(true);<br>
dispose();
                                                                \tt{dispose} ( ) ;
  94
  95 }
  \begin{array}{c|c} 96 & \text{else} \end{array} \begin{array}{c} 97 \end{array}\begin{array}{ll} \textbf{97} & \textbf{Read.Module\_Code\_Credit\_Ready\_Error \ next ErrWindow = new} \\ \textbf{98} & \textbf{Read.Module\_Code\_Predict\_Ready\_Error("Read.Module\_Code\_ Credit\_Read\_Error",} \\ \textbf{99} & \textbf{Credit\_Read\_Error",} \\ \textbf{100} & \textbf{102} & \textbf{103} \\ \textbf{101} & \textbf{nextErrWindow}.\ \textbf{setSize (400, 300)}; \\ \textbf{102} & \textbf{nextErrWindow}.\ \textbf{setDataICloseOperation (JFrame. EXT-ON\_CLOSE)}; \\ \textbf{103} & \textbf{nextErrWindow}.\ \textbf\begin{array}{c|c|c|c} \hline \text{105} & \text{dispose} & \text{)} \end{array}106 }
\begin{bmatrix} 107 \\ 108 \end{bmatrix} } }
                               \rightarrow );
109
110 // Set Layout Manager<br>111 setLayout (new GridLay
                                 \text{setLayout}(\text{new GridLayout}(4, 0));\frac{112}{113}113 \begin{array}{c|c}\n113 \\
\text{this.add}(\text{code\_panel});\n\end{array}114 \begin{array}{c} \text{this. add} \text{ (code panel)}; \\ \text{this. add} \text{ (credit panel)}; \\ \text{this. add} \text{ (message panel)} \end{array}115 this.add(credit_panel);<br>116 this.add(message_panel);
117 \left\{\n \begin{array}{c}\n \text{this } \text{add}(\text{read\_panel}); \\
 \text{118}\n \end{array}\n\right\}\begin{array}{c} 118 \\ 119 \end{array} } // END OF CLASS
```
Read Module Code Credit Ready Error

Listing A.27: Read_Module_Code_Credit_Ready_Error.java

```
\begin{array}{c|c} 1 & \text{import java.awt.}\end{array} import java.awt.event. Action
   2 import java.awt.event.ActionEvent;<br>3 import java.awt.event.ActionListener;
   4 import javax.swing.*;<br>5 import java.sql.CallableStatement;<br>6 import java.sql.Connection;
  5 import java.sql.CallableStatem<br>6 import java.sql.Connection;<br>7 import java.sql.DriverManager;
 8 import java.sql.ResultSet;<br>9 import java.sql.SQLException;<br>10 import java.sql.Statement;
11 import java.sql.Types;
\frac{12}{13}13 public class Read_Module_Code_Credit_Ready_Error extends JFrame {<br>14 public Read_Module_Code_Credit_Ready_Error(String winTitle, String err) {
                                  super(winTitle)\frac{15}{16}<br>\frac{1}{17}17<br>
18 19 17 The same of the code-credit_ready_error_warning_panel = new JPanel();<br>
19 11 11 The same of the code-credit_ready_error_warningLab = new JLabel();<br>
19 12 The code_credit_ready_error_warning_panel<br>
19 12 The co
 add(read_module_code_credit_ready_error_warningLab);<br>
21<br>
JPanel message_panel = new JPanel();<br>
Final JLabel messageLab = <b>new JLabel(err);<br>
24<br>
message_panel.add(messageLab);<br>
JPanel input = pare JPanel();<br>
JPa
27 | input_panel.add(input);
\frac{28}{29}\begin{array}{c|c} \textbf{29} & // \textit{Add Action to the button} \\ \textbf{input.addAction} & \textbf{int} & \textbf{action} \\ \textbf{31} & \textbf{public void actionPerformed (ActionEvent e) } \{ \\ \textbf{32} & \textbf{Input-code.create} \\ \textbf{33} & \textbf{Input-code.create} \\ \textbf{34} & \textbf{Input-code.create} \\ \textbf{35} & \textbf{backWindow.setSize (400, 300)}; \\ \textbf{backWindow.setDecation} & \textbf{learning \textbf{V}} \\ \textbf{36} & \textbf{backWindow.setLocationRelativeTo (null)}; \end{array}37 backWindow . s et V i sible (true);<br>38 dispose ();
\begin{array}{c|c}\n 38 & \text{dispose } ( \text{ )}; \\
 39 & \text{ } \n \end{array}39 }
\begin{array}{c|c}\n 40 & \rightarrow \\
 \end{array}\begin{array}{c} 41 \\ 42 \\ 43 \end{array}\begin{array}{c|c} 42 & // Set \; Layout \; Manager \end{array}<br>
\begin{array}{c} 43 \end{array} setLayout (new GridLayout (3, 0));
\frac{44}{45}\begin{array}{c|cc}\n45 & // Add Panels to the Window\n\end{array}<br>
this add (read module code c
 46 t h i s . add ( r e a d m o d u l e c o d e c r e d i t r e a d y e r r o r w a r n i n g p a n e l ) ;
47 t h i s . add ( m e s s a g e p a n e l ) ;
48 t h i s . add ( i n p u t p a n e l ) ;
 \begin{array}{c} 49 \\ 50 \end{array} } // END OF CLASS
```
Write Module Ready


```
1 import java.awt.GridLayout;<br>
2 import java.awt.event.Actio<br>
3 import java.awt.event.Actio
   2 import java.awt.event.ActionEvent;<br>3 import java.awt.event.ActionListener;
  \begin{array}{c|c} 4 & \text{import java.swing.} \\ 5 & \text{import java.sal.} \end{array}5 import java.sql.CallableStatement;<br>6 import java.sql.Connection;<br>7 import java.sql.DriverManager;
 8 import java.sql.ResultSet;<br>
9 import java.sql.SQLException;<br>
10 import java.sql.Statement;<br>
11 import java.sql.Types;
\begin{array}{c} 12 \\ 13 \end{array}13 public class Write Module Ready extends JFrame {<br>
14 public Write Module Ready (String win Title, final String code,<br>
15 final String credit) {<br>
super (win Title);
\frac{16}{17}\begin{array}{c|c} 18 & // Create Swing Components \\ 19 & J\textbf{Panel code\_panel} = \textbf{new JP} \end{array}19 JPanel code_panel = new JPanel();<br>20 final JLabel codeLab = new JLabel();<br>code_panel.add(codeLab);
 22 JPanel credit_panel = new JPanel();<br>23 final JLabel creditLab = new JLabel();<br>credit_panel.add(creditLab);
 35<br>
1993 JPanel message_panel = new JPanel();<br>
final JLabel messageLab = new JLabel("Are you sure to proceed Updating ");<br>
27 message_panel.add(messageLab);<br>
1994 JPanel write_panel = new JPanel();<br>
1994 JPanel write_panel
\frac{31}{32}\begin{array}{c|ccccc}\n 32 & // Add Action to the button \\
 33 & \text{write.addActionListener( new)}\n\end{array}33 write.addActionListener(new ActionListener(){<br>34 public void actionPerformed(ActionEvent e){
\begin{array}{c|c|c|c|c} \hline \text{51} & \text{62} & \text{63} & \text{64} \\ \hline \end{array} boolean record Updated = false;
\begin{array}{c} 36 \\ 37 \\ 38 \end{array}37 // Declare Connection properties<br>38 Connection conn = null;
 \begin{array}{c|c|c} 39 & \textbf{Statement} & \text{stmt} & = \textbf{null}; \\ 40 & \textbf{int} & \textbf{recAffected} & = 0; \end{array}\frac{41}{42}\begin{array}{c|c|c|c} \hline 42 & & \end{array} // Declare Connection properties
 43 degree try {<br>
44 Class .forName("com . mysql .jdbc . Driver") . newInstance ();
 45<br>
\textbf{String connection} \text{Un} = \text{"jobc:mysql://localhost:3306/testdb1";}<br>
\textbf{String connection} \text{Conserved} = \text{"roct"};<br>
\textbf{String connection} \text{Cannection} \text{Assword} = \text{"v};<br>
\text{conn} = \textbf{DiriverManager}. \text{getConnection} (\text{connectionUn}, \text{connectionUn}, \text{connection}, \text{connection}, \text{con.} \text{con.} \text{correction} \text{Con.} \text\frac{51}{52}\begin{array}{c|c|c|c|c|c|c|c} \hline 52 & // \hline \end{array} Execute a query statement or a stored procedure<br>CallableStatement processes \hline \end{array}53 CallableStatement procnone =<br>
54 conn.prepareCall("{call writeModule(?, ?)}");
^{\rm 55}_{\rm 56}\begin{array}{c|c|c|c|c|c} \hline 56 & & \\ 57 & & \\ 57 & & \\ \hline \end{array} \begin{array}{c} \textbf{56} & \textbf{66} & \textbf{57 procnone.setInt(1, Integer.parseInt(code));<br>58 procnone.setInt(3, Integer.parseInt(credit));
\begin{array}{c} 59 \\ 60 \end{array}60 // Registers the out parameters<br>61 procnone.registerOutParameter(3, Types.INTEGER);
62<br>63\text{rec}A f f e c t e d = pro cnone . execute Update ( ) ;
64<br>6565 // Check if there is an updated record<br>66 if (recAffected > 0)<br>67 recordUpdated = true;
\begin{array}{c|c} 0 & 0 \\ 68 & 0 \end{array} else
\begin{array}{c|c}\n\hline\n\text{69} \\
\hline\n\text{69}\n\end{array} record Updated = false;
^{70}_{70}\begin{array}{c|c|c|c} \hline \text{71} & \text{72} \\ \hline \text{72} & \text{81} \\ \text{82} & \text{9} \\ \text{73} & \text{9} \\ \hline \end{array} \quad \begin{array}{c} \text{73} \\ \text{84} \\ \text{85} \\ \text{96} \\ \text{187} \\ \text{197} \\ \text{108} \\ \text{1198} \\ \text{1208} \\ \text{1309} \\ \text{1400} \\ \text{1500} \\ \text{1600} \\ \text{171} \\ \text{1820} \\ \text{199} \\The state of the contribution of the contr
 75 { e1.printStackTrace(); }<br>
76 try { if (conn != null) conn.close(); } catch (SQLException e1)<br>
77 { e1.printStackTrace(); }
 78 }
\begin{array}{c} 78 \\ 79 \\ 80 \\ 80 \end{array} if (recordUpdated) {
 80 System.out.println(" YES YES YES YES YES YES");<br>
JOptionPane.showMessageDialog(null,
 82 "The business task is completed successfully.");<br>83 System.exit(0);
 84 }
\begin{array}{c|c}\n 85 \\
 86\n \end{array} else {
 86<br>
86<br>
87 System.out.println("NONONONONONONONO");<br>
Write_Module_Ready_Error nextErrWindow = new<br>
Write_Module_Ready_Error ("Write_Module_Ready_Error",<br>
"Can not be Updated. Please Try Again");
```

```
90 nextErrWindow . s e t S i z e ( 4 0 0 , 3 0 0 ) ;
91 nextErrWindow . s e t D e f a u l t C l o s eO p e r a t i o n (JFrame . EXIT ON CLOSE ) ;
 92 nextErrWindow .setLocationRelativeTo(null);<br>93 nextErrWindow .setVisible(true);
 \begin{array}{c|c} 94 & \text{dispose ()}; \\ 95 & \text{ } \end{array}\begin{array}{c} 95 \ 96 \ 97 \end{array} } }
 98
\begin{array}{c|c} 99 & // Set \; Layout \; Manager \end{array}<br>100 setLayout (new GridLayout (4, 0));
\frac{101}{102}\begin{array}{c|cc} 102 & // \textit{Add Panels to the Window} \\ \textbf{this.add(code.panel)} \end{array}103 this.add(code_panel);<br>104 this.add(credit_panel);
\left| \begin{array}{c} \text{time term} \\ \text{this add (message panel)} \end{array} \right|\begin{array}{c|c}\n 106 & \text{this add (write-panel)}; \\
 \hline\n 107 & \end{array}\begin{array}{c} 107 \\ 108 \end{array} } // END OF CLASS
```
Write Module Ready


```
\begin{array}{c|c} 1 & \text{import java.awt} \ . \ \text{GridLayout} \\ \text{import java.awt.event} \ . \ \text{Actio} \end{array}2 import java.awt.event.ActionEvent;<br>3 import java.awt.event.ActionListener;<br>4 import javax.swing.*;
    5 import java.sql.CallableStatement;<br>6 import java.sql.Connection;
    7 | import java.sql.DriverManager;<br>8 | import java.sql.ResultSet;
 9 import java.sql.SQLException;<br>
10 import java.sql.Statement;<br>
11 import java.sql.Types;
\frac{12}{13}13 | public class Write_Module_Ready_Error extends JFrame {<br>14 |       public Write_Module_Ready_Error(String winTitle, String err) {
15 super ( win Title );
\frac{16}{17}\begin{tabular}{ll} $17$ & // Create Swing Components\\  \textbf{JPanel ()};\\ 19$ & \textbf{final JLabel write_model-readv_error-warning\_panel = new JPanel();}\\ 19$ & \textbf{final JLabel write-module.ready_error-warningLab = new JLabel();}\\ 20$ & write_model_ready_error-warning\_panel\\ & .add(write_model=reacy_error-warningLab);\\ 19$ & \textbf{JPanel ()};\\ 21$ & \textbf{JPanel message\_panel = new JPanel();}\\ 23$ & \textbf{final JLabel message\_Label = new JPabel (err);} \end{tabular}24 message_panel.add (messageLab);<br>25 JPanel exit_panel = new JPane
 25 JPanel exit_panel = new JPanel();<br>26 JButton exit = new JButton("exit");<br>27
\frac{28}{29}\begin{tabular}{ll} $29$ & // Add Action to the button \\ \hspace*{1.5cm} \textbf{exit.addActionListener} {\small \textbf{(} } \texttt{with} \texttt{user} {\small \textbf{(} } \texttt{actionListener}) {\small \textbf{(} } \texttt{with} \texttt{user} {\small \textbf{(} } \texttt{in} \texttt{public void actionPerformed (ActionEvent e) } {\small \textbf{(} } \texttt{input_code.createdit_Waiting backWindow = new \\ & \hspace*{1.5cm} \texttt{Input-code.createdit_Waiting backWindow = new \\ & \hspace*{1.5cm} \texttt{Input-code.createdit_Waiting} {\small \textbf{(} ''1nput-code.createdit_Waiting'' } {\small \textbf\begin{array}{c} 39 \\ 40 \end{array} } }
\frac{41}{42}42 // Set Layout Manager<br>43 setLayout (new GridLayout (3, 0));
\frac{44}{45}45 // Add Panels to the Window<br>46 this. add (write_module_ready_error_warning_panel);<br>this. add (message_panel);
48 this . add ( exit panel );<br>49 }
 \begin{array}{c} 49 \\ 50 \end{array} } // END OF CLASS
```
B Appendix: Models and Java Code of Experiment (2)

B.1 Complete Models and Full results of Experiment (2)

This Appendix presents the XML representations of all μ ML models of the Student Enrollment Sub-system case study, and the completed Java and MySQL code generated from these models via BUILD. The experiment presented previously in chapter 10.

Experiment (2): Requirement Models Construction

Listing B.1: CaseStudy2.java


```
38 | import mde taskimpact2dataflow rule DiagramToDfDiagram;<br>39 | import org jast ast ASTWriter;
 \frac{40}{41}41 public class CaseStudy2 {<br>
42 public static void main (String [] args) throws IOException,<br>
TreeException , mde.gui.gen.TreeException {<br>
// TODO Auto−generated method stub
 45
  46<br>
\begin{array}{ll} & // Construct the Task Model & \textbf{Diggram} (); \ & \textbf{Diggram task Model} & = \textbf{new Diqundary} (); \ & \textbf{Boundary boundary} = \textbf{new Boundary("Enrol a Student");} \ & \textbf{Action } \textbf{actor} \text{ action} \textbf{1} = \textbf{new} \textbf{ Actor} \text{ ("Student");} \ & \textbf{Task} \textbf{ enrolStd} = \textbf{new} \textbf{ Fast} \text{ ("Enrol");} \ & \textbf{Participation link 4 = new \textbf{Participation}(); \ & \text{link 4.addRole (new mde. task. model. Role(" such", \textbf{ endStudent}'), \ & \54 Participation link3 = new Participation ();<br>55      link3.addRole(new mde.task.model.Role("student",actor1));<br>56     link3.addRole(new mde.task.model.Role("enrol", enrolStd));
  57 boundary . addActor ( a c t o r 1 ) ;
58 boundary . a d d P a r t i c i p a t i o n ( l i n k 4 ) ;
 59 boundary . add Participation (link 3)<br>
task Model . add Boundary (boundary) ;
                             boundary . add1 arricrpation (rinks)<br>taskModel . addBoundary (boundary);
 \frac{61}{62}62 ASTWriter w r i t e r = new ASTWriter(new F i l e ( " S t u d E n r o l t a s kM o d e l . xml" ) ) ;
63 w r i t e r . us ePa ckag e ( "mde . t a s k . model " , " xmlns : t a s k " ) ;
64 w r i t e r . wr iteDocument ( taskMod e l ) ;
 65 writer, close()66<br>6767 System.out.println("(1) Task Model is Created by user.");<br>
\frac{68}{ }68 // −−−−−−−−−−−−−−−−−−−−−−−−−−−−−−−−−−−−−− //
 69
 71<br>
72<br>
73<br>
74<br>
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79<br>
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171 ImpDiagram ImpactModel = new ImpDiagram ();<br>72 ImpBoundary impboundary = new ImpBoundary("Enrol a Student");
  73<br>
ImpTask impEnrolStd = new mde.impact.model.ImpTask("Enrol");<br>
ImpObject impObj3 = new ImpObject("Enrollment");<br>
ImpCreateFlow cf = new ImpCreateFlow();<br>
ImpRole impcf2 = new ImpRole("enrol", impEnrolStd);<br>
ImpRole imp
  78 cf.addImpRole(impcf2);<br>79 cf.addImpRole(impcf1);
 80 impboundary . addImpObject (impObj3);<br>81 impboundary . addImpCreateFlow(cf);
  81 impboundary . addImpCreateF low ( c f ) ;
82 ImpactModel . addImpBoundary ( impboundary ) ;
 \frac{83}{84}84 ASTWriter writer1 = new ASTWriter(new File("Stud_Enrol_impactModel.xml"));<br>85 writer1.usePackage("mde.impact.model", "xmlns:imp");<br>86 writer1.writeDocument(ImpactModel);
 87 writer 1. close ():
 rac{88}{89}System.out.println (" ( 2) Impact Model is Created by user.");
  90 // −−−−−−−−−−−−−−−−−−−−−−−−−−−−−−−−−−−−−−−−−−−− //
 \frac{90}{91}92<br>
\frac{1}{2} mde. information model Diagram informationModel =<br>
194<br>
mew mde. information nodel Diagram ();<br>
Entity studentEntity = new Entity ("Account");<br>
Entity studentEntity = new Attribute ("regNumber", new Type
104 s t u d e n t E n t i t y . a d dA t t r i b u t e ( a t t r 3 ) ;
105 s t u d e n t E n t i t y . a d dA t t r i b u t e ( a t t r 4 ) ;
106 studentEntity.addAttribute(attr7);<br>107 studentEntity.addAttribute(attr8);
\frac{108}{109}Entity enrolEntity = new Entity("Enrollment");<br>
100 Entity enrolEntity = new Attribute("regNumber", new Type("Integer")).setIdentifier(true);<br>
1111 Attribute attr11 = new Attribute("code", new Type("String")).setIdentifier
114
Entity moduleEntity = new Entity("Module");<br>
116 Entity moduleEntity = new Attribute ("code", new Type("Integer")).setIdentifier(true);<br>
117 Attribute attr13 = new Attribute ("credit", new Type("String"));<br>
Attribute attr1
120 moduleEntity . addAttribute(attr12);<br>121 moduleEntity . addAttribute(attr13);<br>moduleEntity . addAttribute(attr14);
123 moduleEntity . addAttribute (attr15)
124informationModel addEntity (studentEntity);
126 informationModel.addEntity(moduleEntity);<br>
127 informationModel.addEntity(enrolEntity);<br>
128 informationModel.addAssociation(new Association()<br>
.addRole(new Role("module", informationModel<br>
.addRole(new Role("module"))
```
132 . getEntity ("Enrollment")). setMultiple(**true**)));
133 informationModel.addAssociation (**new Association** () 134 . addRo le (new Role (" s t u d e n t " , in f o rm a t i onM od e l 135 . g e t E n t i t y ("Account ")) . s e tM u l t i p l e (f a l s e)) 136 . addRo le (new Role (" e n r o l " , in f o rm a t i onM od e l 136 . addRole (new Role ("enrol", informationModel
137 . addRole (new Role ("enrol", informationModel
137 . et Multiple (true)) ; 138 139 ASTWriter w r i t e r 2 = new ASTWriter(new F i l e (" S t u d E n r o l i n f o rm a t i o nM o d e l . xml")) ; 140 w r i t e r 2 . us ePa ckag e ("mde . i n f o rm a t i o n . model " , " xmlns : i n f ") ; 141 w r i t e r 2 . wr iteDocument (in f o rm a t i onM od e l) ; 142 writer 2. close (); $\frac{143}{144}$ 144 System . out . p r i n t l n (" (3) I n f . Model i s Cr eat ed by u s e r . ") ; 145 // −−− // $\frac{146}{147}$ 147 // G en e r a t e Da taF low Mode l 148 DiagramToDfDiagram topRu l e = new DiagramToDfDiagram () ; 149 DfDiagram data f lowMod e l = topRu l e . t r a n s l a t e (taskMode l , ImpactModel) ; $\frac{149}{150}$ 151 ASTWriter w r i t e r 3 = new ASTWriter(new F i l e (" Stud Enro l DataF lowMod e l . xml")) ; 152 w r i t e r 3 . us ePa ckag e ("mde . d a t a f l ow . model " , " xmlns : d fd ") ; 153 w r i t e r 3 . wr iteDocument (data f lowMod e l) ; 154 writer 3. close (): 155
156 156 System . out . println ("(4) DataFlow Model is Created"+
157 "(Task + Impact Model --> DataFlow Model)."); 157 " (Task + Impact Model −−> DataFlow Model) . ") ; 158 // −−− // 159
160 $\begin{array}{lll} 160 & // Adding some data on flows in dataflow model. \\ 161 & // input to login task \\ 162 & dataflow Model.getDfBoundaries(), get 0), get DfIn \end{array}$ $\begin{tabular}{ll} 161 & // input to login task \\ 162 & dataflow Model.\ getDfBoundaries().get(0).getDfInputFlows() \\ 163 & .get(0).addDataonflow("regNumber"); \\ 164 & .get(1).addDataonflow("regNumber"); \\ 165 & .get(1).addDataonflow("code"); \\ 166 & // create assignment \end{tabular}$ 167 data f lowMod e l . g e tD fB ound a r i e s () . g e t (0) . g e tD fC r e a t eF l ow s () 168 . g e t (0) . addAss ignment ("@regNumber = regNumber ") ; 169 data f lowMod e l . g e tD fB ound a r i e s () . g e t (0) . g e tD fC r e a t eF l ow s () 170 . g e t (0) . addAss ignment ("@code = cod e ") ; 171 // create
172 dataflowModel.getDfBoundaries ().get(0).getDfCreateFlows ()
173 .eet (0).addDataonflow ("regNumber");
174 dataflowModel.getDfBoundaries ().get(0).getDfCreateFlows ()
175 .get (0).addDataonflow ("code"); 176 177 System .out . println ("(4) Data on flows are added"+
178 "by user in the DataFlow model."); 179 // −−− // 180 181 // G en e r a t e D e t a i l e d Da tF low Mode l 182 DfDiagramToDfDiagram topRu l e6 = new DfDiagramToDfDiagram () ; 183 DfDiagram d e t a i l e d d a t a f l o wM o d e l = topRu l e6 . t r a n s l a t e (data f lowMod e l) ; 184 185 ASTWriter w r i t e r 6 = new ASTWriter(new F i l e (" S tud En r o l D e t a i l ed D a t a F l owM od e l . xml")) ; 186 w r i t e r 6 . us ePa ckag e ("mde . d a t a f l ow . model " , " xmlns : d fd ") ; 187 w r i t e r 6 . wr iteDocument (d e t a i l e d d a t a f l o wM o d e l) ; 188 writer6.close() 189 190 System . out . p r i n t l n (" (5) D e t a i l e d DataFlow Model i s Cr eat ed "+ 191 " (DataFlow Model −−> DataFlow Model) . ") ; 192 // −− // $\frac{152}{193}$ 194 // Generate Data Model
195 **InfDiagramToDDiagram** 195 **InfDiagramToDDiagram** topRule4 = **new InfDiagramToDDiagram**();
196 **DDiagram** dataModel = topRule4.translate(informationModel); $\frac{197}{198}$ 198

 ASTWriter writer4 = **new ASTWriter** (new File ("Stud_Enrol_DataModel.xml"));

writer4.usePackage("mde.data.model", "xmlns:data");

201 writer4.writeDocument(dataModel);

writer4.writeDocument(dataModel); 202 writer 4. close (); $\frac{203}{204}$ 204 System .out .println("(5) Data Model is Created"+
205 "(Information Model −−> Data Model)."); 206 // −− // $\frac{207}{208}$ 208 //Generate State Model for Screens Navigation
209 // (DataFlow Model ---> State Model)
210 **DidDiagramToStDiagram** topRule7 = **new DidDiagramToStDiagram**();
211 **StDiagram stateModel = topRule7.translate(detailed_dataf** $\frac{211}{212}$ 213
213 **ASTWriter** writer7 = **new ASTWriter** (new File ("Stud_Enrol_StateModel.xml"));
214 writer7.wsePackage ("mde.state.model", "xmlns:state");
215 writer7.chiteDocument (stateModel);
writer7.close (); $\frac{217}{218}$ 218 System . out . println ("(6) State Model is Created (Detailed DataFlow Model --> State Model). "
210 // −−− // ; 219 // −−− // $\frac{220}{221}$ 221 //*Generate Gui Description Model*
222 //(*State Model ----> Gui Specification Model*)
223 **StDiagramToGuiDiagram** topRule9 = **new StDiagramToGuiDiagram**();
224 **GuiDiagram GuiModel = topRule9.translate(stateModel);** 225

```
226 ASTWriter writer9 = new ASTWriter(new File("Stud_Enrol_GuiModel.xml"));<br>227 writer9.usePackage("mde.gui.model", "xmlns:gui");
228 writer9.writeDocument (GuiModel);<br>229 writer9.close ();
\frac{230}{231}231 System . out . p r i n t l n ( " ( 7 ) Gui S p e c i f i c a t i o n Model i s Cr eat ed ( S t a t e Model −−> Gui Model ) . " ) ;
232 // −−−−−−−−−−−−−−−−−−−−−−−−−−−−−−−−−−−−−−−−−−−− //
233
234 // generate Database Schema Model<br>235 // (Detailed_DFD + Data Model ----
235 // ( D e t a i l e d DFD + Data Mode l −−−> Schema )
236 DDiagramToSchema2 topRu l e5 = new DDiagramToSchema2 ( ) ;
237 Schema schemaModel = topRu l e5 . t r a n s l a t e ( d e t a i l e d d a t a f l o wM o d e l , dataMode l ) ;
238
                        System . out . println (schemaModel . getName ());
\frac{239}{240}241 ASTWriter writer5 = new ASTWriter (new File ("Stud_Enrol_SchemaModel.xml"));<br>242 writer5.usePackage("mde.dbs.model", "xmlns:dbs");<br>343 writer5.usePackage("mde.model", "xmlns:mod");<br>writer5.writeDocument(schemaMod
\begin{array}{c} 244 \overline{\smash{\big)} } \\ 245 \end{array} writer 5. close ():
\frac{246}{247}247 System.out.println("(9) Database Schema Model is Created"+<br>248 " (DataFlow + Data Model −−> Database Schema Model).");
249 // −−−−−−−−−−−−−−−−−−−−−−−−−−−−−−−−−−−−−−−−−−−−− //
\frac{250}{251}251 //Generate Code Model<br>252 // (Database Schema + DataFlow Model ---> Code Model)<br>253 DDiagramToCDiagram2x topRule8 = new DDiagramToCDiagram2x();<br>254 Diagram codeModel = topRule8
255 . The contract of the set of the contract 
\frac{256}{257}257 ASTWriter writer8 = new ASTWriter(new File("CodeModel.xml"));<br>258 writer8.usePackage("mde.code.model", "xmlns:code");
259 writer8.usePackage("mde.model", "xmlns:mod");<br>260 writer8.writeDocument(codeModel);
261 | writer 8. close ();
\frac{262}{263}263 System . out . p r i n t l n ( " ( 1 0 ) Code Model i s Cr eat ed "+
264 " ( Database Scehma Model −−> Code Model ) . " ) ;
265 System . out . p r i n t l n ( "\nModel T r an s f o rm a t i on System i s Completed . . . " ) ;
\frac{200}{266}267<br>268 // Code Generation<br>DumpFileGenerator
268 DumpFileGenerator MySQLGenerator = new MySQLDumpFileGenerator( schemaModel ) ;
269 System . out . p r i n t l n ( " c a l l i n g g e n e r a t e i n mysq l pa ckag e g e n e r a t o r " ) ;
270 MySQLGenerator . g e n e r a t e ( ) ;
271 System.out.println ("Finished generating MySQL Schema OK");
\frac{272}{273}\textbf{for}(\textbf{GuiBoundary} \text{ bound : } \textbf{GuiModel.getGuiBoundaries}())274275 CodeFileGenerator JavaGenerator = new JavaCodeFileGenerator (bound);<br>276 System out println ("calling generate in java code file generator");
276 System .out .println("calling generate in java code file generator");<br>277 JavaGenerator .generate();
278 System . out . println ("Finished generating Java Gui code");<br>279 }
\begin{array}{c|c} 279 & & \\ 280 & & \end{array}280281 }
```
Task model: Student Registration System (Enrol a Student)

Listing B.2: Stud_Enrol_taskModel.xml

```
\begin{array}{c|c} 1 < ? \ \textbf{xml} & \textbf{version} = "1.0" & \textbf{encoding} = "UTF-8" ? > \ \textbf{task:} \textbf{Diagram} & \textbf{whn:} \ \textbf{task:} = "mde, task, move, \end{array}2 \left\{\text{task: Diagram } xmlns: task="mde. task.\ model" id="0" \right\} \right. \left\{\text{task:Boundary } id="1" name="Enrol a Student" > \text{task:Boundary } id="2" name="Enrol a Student" > \text{task:Partition } id="3" > \text{task:Partition } id="3" > \text{task:Partition } id="3" > \text{task:Matrix } id="4" name="student" > \text{task:Value } id="5" > \text{task:Value } id="6" > \text{task:Value } id="7" \right\}.\begin{array}{c}\n 8 \\
 8 \\
 \hline\n \end{array} \begin{array}{c}\n \langle \text{task:Role}\rangle \\
 \langle \text{task:Role}\rangle \\
 \end{array}9 <task:Role id="5" name="enrol"><br>10 </task:Task id="6" name="Enrol"/><br>11 </task:Role>
12 </ ta s k :Pa r t ic ipa t ion>
13 <task:Participation id="7"><br>
14 <task:Role id="8" name=" student"><br>
15 </task:Rotor ref="2"/><br>
16 </task:Role><br>
17 <task:Role id="9" name=" enrol">
18 <task:Task r e f="6 "/>
19 </ task:Role>
\begin{array}{c} 20 \\ 20 \end{array} </task:Participation>
\begin{array}{c|c} 21 & \angle/\texttt{task:Boundary}> \\ 22 & \angle/\texttt{task:Diagram}> \end{array}\langle/task:Diagram>
```
Impact model: Student Registration System (Enrol a Student)

Listing B.3: Stud Enrol impactModel.xml

```
1 <imp:ImpDiagram xm lns : imp="mde . impact . model " i d="0 ">
  2 <imp:ImpBoundary id="1" name="Enrol a Student"><br>
3 <imp:ImpObject id="2" name="Enrollment"/><br>
4 <imp:ImpTcreateFlow id="3"><br>
4 <imp:ImpTcreateFlow id="4" name="enrol"><br>
4 <imp:ImpTask id="5" name="Enrol"/>
  7 </imp:ImpRole>
8 <imp:ImpRole i d="6" name=" e n r o l lm e n t ">
9 <imp:ImpObject r e f="2 "/>
10 </imp:ImpRole>
\begin{array}{c|c}\n 11 & & \langle / \text{imp:ImpCreateFlow} \rangle \\
 & & \langle / \text{imp:ImpBoundar} \mathbf{v} \rangle\n\end{array}\begin{array}{c|c}\n 12 & \text{(imp:ImpBoundary)} \\
 \hline\n 13 & \text{(imp:ImpDiagram)}\n \end{array}\langle / \text{imp:ImpDiagram} \rangle
```
Inofrmation model: Student Registration System (Enrol a Student)

Listing B.4: Stud Enrol InformationModel.xml

```
1 <inf:Diagram xmlns:inf="mde.information.model" id="0"><br>
2 <inf:Entity id="23" name="Enrollment"><br>
<inf:Attribute id="24" name="regNumber" identifier="true" size="0"><br>
4 <Type id="25" name="Integer"/><br>
</inf:Attribute>
   6 <in f :A t t r ibu te i d=" 26 " name=" cod e " i d e n t i f i e r=" t r u e " s i z e="0 ">
7 <Type i d=" 27 " name=" S t r i n g "/>
8 </ in f :A t t r ibu te>
\begin{array}{c|c} 9 & \sqrt{\text{inf}}: \text{Entity} > 10 & \sqrt{\text{inf}}: \text{Diagram} > \end{array}\langleinf:Diagram>
```
DataFlow model (initial): Student Registration System (Enrol a Student)

Listing B.5: Stud Enrol DataFlowModel.xml

$\mathbf{1}$	\langle dfd:DfDiagram xmlns:dfd="mde.dataflow.model" id="0">
$\mathbf{2}$	<dfd:dfboundary id="1" name="Enrol a Student"></dfd:dfboundary>
3	<dfd:dfactor id="2" name="Student"></dfd:dfactor>
$\overline{4}$	\langle dfd:DfObject id="3" name="Enrollment"/>
5	\langle dfd:DfInputFlow id="4">
6	<dfd:dfrole id="5" name="student"></dfd:dfrole>
$\overline{7}$	\langle dfd:DfActor ref="2"/>
$\,$ 8 $\,$	\langle dfd:DfRole>
9	<dfd:dfrole id="6" name="enrol"></dfd:dfrole>
10	<dfd:dftask id="7" name="Enrol"></dfd:dftask>
11	\langle dfd:DfRole>
12	
13	\langle dfd:DfInputFlow id="8">
14	<dfd:dfrole id="9" name="student"></dfd:dfrole>
15	\langle dfd:DfActor ref="2"/>
16	\langle dfd:DfRole>
17	<dfd:dfrole id="10" name="enrol"></dfd:dfrole>
18	\langle dfd:DfTask ref="7"/>
19	\langle dfd:DfRole>
20	\langle dfd:DfInputFlow>
21	<dfd:dfcreateflow id="11"></dfd:dfcreateflow>
22	<dfd:dfrole id="12" name="enrol"></dfd:dfrole>
23	\langle dfd:DfTask id="13" name="Enrol"/>
24	\langle dfd:DfRole>
25	\langle dfd:DfRole id="14" name="enrollment">
26	<dfd:dftask id="15" name="Enrollment"></dfd:dftask>
27	\langle dfd:DfRole>
28	\langle dfd:DfCreateFlow>
29	\langle dfd:DfBoundary>
30	\langle dfd:DfDiagram>

DataFlow model (detailed): Student Registration System (Enrol a Student)

Data Dependency model: Student Registration System (Enrol a Student)

Listing B.7: Stud Enrol DataModel.xml

```
data:DDiagram xmlns:data="mde.data.model" id="0"><br>
<data:DEntity id="23" name="Enrollment"><br>
<data:DAttribute id="24" name="Enrollment"><br>
<data:DAttribute id="24" name="regNumber" identifier="true" size="10"><br>
<data:DAttri
\left| \right| </data:DDiagram>
```
Screen State model: Student Registration System (Enrol a Student)

GUI model: Student Enrollment System (Enrol a Student)

Listing B.9: Stud Enrol GuiModel.xml

Database and Query model: Student Enrollment System (Enrol a Student)

Listing B.10: Manage_Module_SchemaModel.xml

OOP Code model: Student Enrollment System (Enrol a Student)

Listing B.11: Stud Enrol CodeModel.xml

MySQL Script: Student Registration System (Enrol a Student)

Listing B.12: database MySQL.sql

Java Swing Source Code: Student Enrollment System

Input_code_Waiting

Listing B.13: Input_code_Waiting.java

¹ import java.awt.GridLayout;
2 import java.awt.event.ActionEvent;

³ import java.awt.event. ActionListener;
4 import javax.swing.*;

```
5 import java.sql.CallableStatement;<br>6 import java.sql.Connection;
   7 import java.sql.DriverManager;<br>8 import java.sql.ResultSet;
 9 import java.sql. SQLException;<br>10 import java.sql. Statement;
11 | import java.sql.Types;
\frac{12}{13}13 public class Input_code_Waiting extends JFrame {<br>14 public Input_code_Waiting(String winTitle, final String regNumber) {
15 super (win Title);
\frac{16}{17}17 17 // Create Swing Components<br>18 19 - JPanel message-panel = new JPanel();<br>19 11 - final JLabel messageLab = new JLabel("Input your Data");
20<br>
message_panel . add (messageLab);<br>
21<br>
JPanel code_panel = new JPanel
 21 JPanel code<sub>-P</sub>anel = new JPanel();<br>22 JLabel code<sub>-</sub>label = new JLabel("code");<br>23 final JTextField codeTxt = new JTextField(30);
 \begin{array}{ll} 24 & {\rm code\_panel.\,add(\,code\_label\,);\ & \\ 25 & {\rm code\_panel.\,add(\,code{\rm Txt}\,);\ } \end{array}<br>
\begin{array}{ll} 26 & {\rm J\,Panel\, \, in\,put\_panel\,.\ & \\ {\rm J\,Button\, \, in\,put\, =\, new\, \,J\,Button("Input");\ } \end{array}<br>
\begin{array}{ll} 27 & {\rm J\,Button\, \, in\,put\, =\, new\, \,J\,Button("Input");\ } \end{array}\frac{29}{30}<br>30<br>31
 30 // Add Action to the button<br>
32 input addActionListener (new ActionListener () {<br>
public void actionPerformed (ActionEvent e) {<br>
String code = codeTxt .getText ();
\frac{34}{35}\begin{array}{c} \texttt{35} \\ \texttt{46} \\ \texttt{37} \\ \texttt{38} \\ \texttt{39} \\ \texttt{40} \\ \texttt{41} \\ \texttt{42} \\ \texttt{43} \\ \texttt{44} \\ \texttt{45} \\ \texttt{46} \\ \texttt{47} \\ \texttt{48} \\ \texttt{49} \\ \texttt{40} \\ \texttt{41} \\ \texttt{42} \\ \texttt{43} \\ \texttt{44} \\ \texttt{45} \\ \texttt{49} \\ \texttt{40} \\ \texttt{41} \\ \texttt{42} \\ \texttt{43} \\ \texttt{444 }
\begin{array}{c|c}\n 45 \\
 46\n\end{array} else
45 else {<br>Create_Enrollment_Ready nextWindow = new<br>47 Create_Enrollment_Ready ("Create_Enroll
 47 C r e a t e En r o l lm en t R e ad y ( " C r e a t e En r o l lm en t R e ad y " , regNumber , cod e ) ;
48 nextWindow . s e t S i z e ( 4 0 0 , 3 0 0 ) ;
49 nextWindow . s e t D e f a u l t C l o s e O p e r a t i o n (JFrame . EXIT ON CLOSE ) ;
 50 nextWindow .setLocationRelativeTo(null);<br>51 nextWindow .setVisible(true);
52 d i s p o s e ( ) ;
 53 }
 54 }<br>55 });<br>56 // Set Layout Manager
                              \texttt{SetLayout}(\textbf{new GridLayout}(3, 0));\frac{57}{58}\begin{array}{c|cc}\n 59 & // Add Panels to the Window \n 60 & this.add(message\_panel):\n \end{array}60 this . add ( message_panel );<br>
62 this . add ( code_panel );<br>
this . add ( input_panel );
 \begin{array}{c} 63 \\ 64 \end{array} } // END OF CLASS
```
Input code Waiting Error

Listing B.14: $Input_{c}ode_Waiting_Error.java$

$\mathbf{1}$	import java.awt.GridLayout;
$\boldsymbol{2}$	import java.awt.event.ActionEvent;
3	import java.awt.event.ActionListener;
$\overline{4}$	import javax.swing.*;
5	import java.sql. CallableStatement;
6	import java.sql. Connection;
$\overline{7}$	import java.sql.DriverManager;
8	import java.sql.ResultSet;
9	import java.sql.SQLException;
10	import java.sql.Statement;
11	import java.sql.Types;
12	
13	public class Input_code_Waiting_Error extends JFrame {
14	public Input_code_Waiting_Error(String winTitle, String err) {
$15\,$	super(winTitle);
16	
17	// Create Swing Components
18	JPanel input_code_waiting_error_warning_panel = new JPanel ();
19	final JLabel input_code_waiting_error_warningLab = new JLabel();
20	input_code_waiting_error_warning_panel
21	.add(input_code_waiting_error_warningLab);
22	JPanel message_panel = new JPanel();
23	final JLabel messageLab = new JLabel(err);

Input regNumber Waiting


```
\begin{array}{c|c} 1 & \text{import java.awt.GridLayout}; \\ \hline 2 & \text{import java.awt.event.Actio} \end{array}2 import java.awt.event.ActionEvent;<br>3 import java.awt.event.ActionListen
  3 import java.awt.event.ActionListener;<br>4 import iavax.swing.*:
   4 import javax.swing.*;<br>5 import java.sql.CallableStatement;<br>6 import java.sql.Connection;
   7 import java.sql.DriverManager;<br>8 import java.sql.ResultSet;<br>9 import java.sql.SQLException;
10 import java.sql. Statement;<br>11 import java.sql. Types:
          import java.sql.Types;
\frac{12}{13}13 public class Input regNumber Waiting extends JFrame {<br>14 public Input regNumber Waiting (String win Title) {<br>15 super (win Title);
\begin{array}{c} 16 \\ 17 \\ 18 \end{array}\begin{tabular}{ll} $17$ & // Create Swing Components \\ 18 & JPanel message.panel = new JPanel(); \\ 19 & final JLabel messageLabel = new JLabel("Input your Data"); \\ 20 & message.panel.add(messageLab);\quad JPanel (p) \\ 21 & JPanel (pgNumber.panel = new JPanel(); \\ 22 & JLabel regNumber-label = new JLabel("regNumber"); \\ 23 & final JTextField regNumberText = new JTextField(30);\n    regNumber.panel.add(regNumberText);\n    regNumber.panel.add(regNumberText);\n    JPanel input.panel = new JPanel(); \\ 24 & JButton("Input"); \\ 25 & input.panel.add(input);\n    }\end{tabular}\frac{28}{29}<br>30
 30 30 // Add Action to the button<br>32 input addActionListener (new ActionListener () {<br>public void actionPerformed (ActionEvent e) {
33
                                                                   String \; regNumber = \; regNumberTxt.\;getText();
\begin{array}{c} 34 \\ 35 \\ 36 \end{array}\begin{array}{rcl} \textbf{36} & \textbf{if } ((\texttt{regNumberTxt}.\texttt{getText}().\texttt{isEmpty}))) & \{ \\ & \texttt{Input-regNumber}.\texttt{Waiting}.\texttt{Error nextString} & \texttt{PortWindow} = \texttt{new} \\ \textbf{38} & \texttt{Input-regNumber}.\texttt{Waiting}.\texttt{Error "Input-regNumber}.\texttt{Waiting}.\texttt{Error "} \\ \textbf{39} & \texttt{Input}.\texttt{Please (100, 300)}; \\ & \texttt{nextErrWindow}.\texttt{setDefaultCloseOperator(JFrame.EXT-ON.CLOSE}); \end{array}42 nextErrWindow .setLocationRelativeTo(null);<br>nextErrWindow .setVisible(true);
\begin{array}{c|c}\n 44 \\
 45 \\
 \hline\n 45\n \end{array}45 }
\begin{array}{c|c} 13 & \text{else} \end{array} else
 47 Input_code_Waiting nextWindow = new<br>48 Input_code_Waiting("Input_code_Waiting", regNumber);<br>nextWindow.setSize(400, 300);
 50 nextWindow .setDefaultCloseOperation (JFrame .EXIT_ON_CLOSE);<br>52 nextWindow .setUocationRelativeTo (null );<br>nextWindow .setVisible (true );
\begin{bmatrix} 53 \\ 53 \end{bmatrix} dispose ();
 54 }
\begin{array}{c|c}\n 54 \\
 55 \\
 56 \\
 57\n \end{array} } ;<br>
\begin{array}{c}\n 3 \\
 \end{array}56 } ) ;
57 // S e t L a y ou t Manager
```
 58 set Layout (new GridLayout $(3, 0)$); $\frac{59}{60}$ 60 $\begin{array}{ll}\n\end{array}$ // Add Panels to the Window

62 this .add (regnumber_panel);

63 this .add (regnumber_panel); $64\atop 65$ END OF CLASS

Input regNumber Waiting Error


```
1 import java.awt.GridLayout;<br>
2 import java.awt.event.Actio<br>
3 import java.awt.event.Actio
   2 import java.awt.event.ActionEvent;<br>3 import java.awt.event.ActionListener;
   4 import javax.swing.*;<br>5 import java.sql.CallableStatement;<br>6 import java.sql.Connection;
  <sup>7</sup> import java.sql. DriverManager ;<br>8 import java.sql. ResultSet :
   8 import java.sql.ResultSet;<br>9 import java.sql.SQLException;
10 import java.sql.Statement;<br>11 import java.sql.Types;
         import java.sql.Types;
\begin{array}{c} 12 \\ 13 \end{array}<sup>13</sup> public class Input regNumber Waiting Error extends JFrame {<br>14<br>14 public Input regNumber Waiting Error (String winTitle
 14 public Input_regNumber_Waiting_Error(String winTitle, String err) {<br>15 super(winTitle);
\frac{16}{17}17 // Create Swing Components<br>18 JPanel input-regnumber-wa
 18 JPanel i n p u t r e g n um b e r w a i t i n g e r r o r w a r n i n g p a n e l = new JPanel ( ) ;
19 f i n a l JLabel i n p u t r e g n um b e r w a i t i n g e r r o r w a r n i n g L a b = new JLabel ( ) ;
 20 i n p u t r e g n um b e r w a i t i n g e r r o r w a r n i n g p a n e l
21 . add ( i n p u t r e g n um b e r w a i t i n g e r r o r w a r n i n g L a b ) ;
22 JPanel m e s s a g e p a n e l = new JPanel ( ) ;
23 f i n a l JLabel messageLab = new JLabel ( e r r ) ;
24 m e s s a g e p a n e l . add ( messageLab ) ;
25 JPanel i n i t i a l i s e p a n e l = new JPanel ( ) ;
26 JButton i n i t i a l i s e = new JButton( " i n i t i a l i s e " ) ;
                              initialise\_panel.add(imitialise);\frac{27}{28}<br>28
\begin{array}{c|c|c|c} \hline 29 & // \hline \end{array} Add Action to the button<br>30 initialise.addActionListener
30 initialise.addActionListener (new ActionListener () {<br>public void actionPerformed (ActionEvent e) {<br>Input_regNumber_Waiting backWindow = ne
 31<br>
Public void actionPerformed(ActionEvent e) {<br>
Input_regNumber_Waiting backWindow = new<br>
\frac{1}{33}<br>
34<br>
backWindow setSize(400, 300);<br>
backWindow setDefaultCloseOperation(JFrame EXIT_ON_CLOSE);<br>
backWindow setLocation
37<br>37 backWindow . s et V i s i ble ( true ) ;<br>38 dispose ( ) ;
\begin{array}{c|c}\n 38 & \text{dispose} (); \\
 & & \text{ } \\
\end{array}39 }
\begin{array}{c|c} 30 & & \end{array} } );
\begin{array}{c} 41 \\ 42 \end{array}\begin{array}{c|c} 42 & // Set \; Layout \; Manager \end{array}<br>
\begin{array}{c} 43 \end{array} setLayout (new GridLayout (3, 0));
\frac{44}{45}45 // Add Panels to the Window<br>46 this. add(input_regnumber_w
46<br>
this . add (input regnumber waiting error warning panel);<br>
this . add (message panel) ;
 47 this . add (message_panel);<br>
48 this . add (initialise_panel);
\frac{49}{50}\frac{1}{2} // END OF CLASS
```
Create_Enrollment_Ready

280

```
14 public Create_Enrollment_Ready (String winTitle, final String regNumber, final String code) {
                                 final String code) {
                              super(winTitle);\frac{16}{17}\begin{tabular}{ll} $18$ & // Create Swing Components\\  \textbf{19} & \textbf{JPanel} = [1] & \textbf{ILabel} = \textbf{new JPanel}();\\  \textbf{20} & \textbf{final JLabel}~ \textbf{regular} = \textbf{new JLabel}();\\  \textbf{21} & \textbf{regular}~ \textbf{regular}~ \textbf{1}~ \text29 \n\begin{array}{r}\n 29 \n \end{array} JPanel create_panel = new JPanel();<br>
30 JButton create = new JButton("Create");
                              \frac{31}{2} create_panel.add (create);
 31<br>32<br>3333 33 // Add Action to the button<br>34 create.addActionListener(new ActionListener(){<br>35 public void actionPerformed(ActionEvent e){
 36 boolean record Inserted = false;
 \frac{37}{38}\begin{array}{c|c|c|c} \hline 38 & & \\ \hline 39 & & \textbf{Connection conn = null}; \\ \hline 40 & & \textbf{Statement sum = null}; \\ \hline \end{array}39 Connection conn = null;<br>
Statement stmt = null;<br>
int recAffected = 0;
 \frac{42}{43}43 // Declare Connection properties<br>44 try {
  \begin{array}{ll} \textbf{try} \; \{ \\ 45 & \textbf{Class. forName("com. mysq1.jdbc. Drive" ) . newInstance } \textbf{();} \\ 46 & \textbf{String connection Ur1 = "jdbc: mysq1://localhost:3306/testdb1";} \\ 47 & \textbf{String connectionUser = "root";} \\ 48 & \textbf{String connectionParser = "s}; \\ 49 & \textbf{conn = DriverManager.getConnection (connectionUr1, & ...)} \\ \end{array}50 connectionUser, connectionPassword);<br>
51 stmt = conn.createStatement();
 \frac{52}{53}\left| \begin{array}{rcl} \end{array} \right| / \left| \begin{array}{rcl} \textit{Execute a query statement or a stored procedure} \ \end{array} \right| \left| \begin{array}{rcl} \end{array} \right| \left| \begin{array}{rcl} \textit{if} \ \textit{\begin{array}{c} 55 \\ 56 \end{array}56 // S e t s t h e i n p u t p a r am e t e r
57 pro cnon e . s e t I n t ( 1 , I n t e g e r . p a r s e I n t ( regNumber ) ) ;
58 pro cnon e . s e t S t r i n g ( 2 , cod e ) ;
 \begin{array}{c} 59 \\ 60 \end{array}60 \begin{array}{ccc} \text{60} & \text{7} & \text{References} \ \text{61} & \text{63} & \text{55} \end{array} recAffected = procnone.executeUpdate ();
 6263 // Check if there is an updated record<br>64 if (recAffected > 0)<br>65 cordInserted = true;
 \begin{array}{c|c}\n66 & \text{else}\n\end{array}recordInserted = false;68<br>6969 } catch (Exception e1)<br>
70 } e1.printStackTrace<br>
71 } finally {
  ro<br>
anally {<br>
if is the set of the 
  76 }
 76<br>77<br>7878 if (recordInserted) {<br>
30 ptionPane . showMessageDialog(null, "The business task"+<br>
80 m is completed successfully .");<br>
81 System . exit (0);
  82 }
 \begin{array}{c|c}\n 82 & & & \n 83 & & \n 83 & & \n 84 & & \n\end{array} else {
 84 Create_Enrollment_Ready_Error nextErrWindow = new<br>
85 Create_Enrollment_Ready_Error ("Create_Enrollment
  85 C r e a t e E n r o l lm e n t R e a d y E r r o r ( " C r e a t e E n r o l lm e n t R e a d y E r r o r " ,
86 "Can not be Cr eat ed . P l e a s e Try Again " ) ;
87 nextErrWindow . s e t S i z e ( 4 0 0 , 3 0 0 ) ;
88 nextErrWindow . s e t D e f a u l t C l o s eO p e r a t i o n (JFrame . EXIT ON CLOSE ) ;
 89 nextErrWindow . setLocationRelativeTo (null);<br>90 nextErrWindow . setVisible (true):
 90 \begin{array}{c|c} \hline \text{nextErrWindow.} \text{setVisible (true)}; \\\hline \text{dispose} \hline \end{array}\begin{array}{c|c} 91 & \text{dispose } ( ) \, ; \end{array}92 }
 93 }
  94 } ) ;
95 // S e t L a y ou t Manager
 96 \Big\{\n \begin{array}{c}\n \text{setLayout (new GridLayout (4, 0));}\n \end{array}\n \right.\frac{97}{98}\begin{array}{c|c|c|c|c} \text{98} & // Add Panels to the Window \\ \text{199} & \text{199} & \text{190} & \text{190} \\ \end{array}99 this.add (regnumber_panel);<br>
100 this.add (code_panel);<br>
101 this.add (message_panel);
\begin{array}{c|c} 102 & \text{this.add}(\text{create\_panel}); \\ 103 & \end{array}\begin{array}{c} 103 \\ 104 \end{array} } // END OF CLASS
```
Create Enrollment Ready Error

Listing B.18: Create Enrollment Ready Error.java

```
\begin{array}{c|c} 1 & \text{import java.awt.GridLayout}; \\ \hline 2 & \text{import java.awt.event.Action} \end{array}2 import java.awt.event.ActionEvent;<br>3 import java.awt.event.ActionListener;
    4 import javax.swing.*;<br>5 import java.sql.CallableStatement;<br>7 import java.sql.Connection;<br>7 import java.sql.DriverManager;
 8 import java.sql.ResultSet;<br>9 import java.sql.SQLException;<br>10 import java.sql.Statement;
             import java.sql.Types;
\begin{array}{c} 11 \\ 12 \\ 13 \end{array}public class Create_Enrollment_Ready_Error extends JFrame {
public Create_Enrollment_Ready_Error(String winTitle, String err) {
                                           super(winTitle);\frac{15}{16}<br>\frac{1}{17}17<br>
18 // Create Swing Components<br>
19 JPanel create_enrollment_ready_error_warning_panel = new JPanel();<br>
19 final JLabel create_enrollment_ready_error_warningLab = new JLabel();<br>
create_enrollment_ready_error_warn
 add(create_enrollment_ready_error_warningLab);<br>
1922 JPanel message_panel = new JPanel();<br>
1923 final JLabel messageLab = new JRahel(err);<br>
23 message_panel.add(messageLab);<br>
1926 JPanel exit_panel = new JPanel();<br>
1926 JB
\frac{27}{28}<br>28
 \begin{tabular}{ll} 29 & // Add Action to the button \\ \hspace*{1.5cm} \textbf{exit.} addAction Listener() { \\ 30 & public void actionPerformed (ActionEvent e) { \\ 31 & public void actionPerformed (ActionEvent e) { \\ 32 & Input-regNumber.Waiting backWindow = new \\ & Input-regNumber.Waiting("Input-regNumber in the text) { \\ 33 & backWindow. setSize(400, 300); \\ & backWindow. setDefaultCloseOperation (JFrame. EXIT-ON\_CLOSE) { \\ & backWindow. setLocalInRelativeTo (null) { \\ & backWindow. setVisual (Cluse) {\end{tabular}} \end{tabular}\begin{array}{c} 38 \ 39 \ 40 \ 41 \ 42 \end{array} \qquad \qquad \begin{array}{c} \text{dispose }(\;); \ \text{d} \text{ispose }(\;); \ \text{d} \text{is positive } \42<br>43<br>44<br>4544 // Add Panels to the Window<br>45 this.add(create_enrollment_ready_error_warning_panel);<br>this.add(message_panel);
\begin{array}{c|c}\n 47 \\
 48\n \end{array} this . add (exit_panel);
 \begin{array}{c} 48 \\ 49 \end{array} } // END OF CLASS
```
\bigcup_{\sim} Appendix: Models and MySQL Scripts of Experiment (3)

C.1 Complete Models and Full results of Experiment (3)

This Appendix presents the XML representations of all μ ML models In The University Administration System case study, discussed previously in chapter 10.

Experiment (3a): Information Model Construction

Listing C.1: CaseStudy3.java

```
1 \mid package mde. example;
 \frac{2}{3}\begin{array}{c|c} \text{z} \\ \text{3} \\ \text{4} \end{array} import java.io.IOEx
  4 import java.io.IOException;<br>5 import mde.data.model.DDiagram;
  6 import mde . d a t a b a s e . gen . DumpFileGenerator ;<br>7 import mde . mysql . gen . MySQLDumpFileGenerator ;
  8 import mde .database .gen .TreeException;<br>9 import mde .dbs .model .Schema;
10 import mde.dm2schem.rule.DDiagramToSchemaOnly;<br>11 import mde.inf2dm.rule.InfDiagramToDDiagram;
12 import mde. information. model. Association;<br>13 import mde. information. model. Composition;
14 import mde.information.model.Entity;<br>15 import mde.information.model.Attribute;<br>16 import mde.information.model.Generalisation;<br>17 import mde.information.model.Role;
18 import mde. model. Type;<br>19 import org. jast. ast. ASTWriter;
\begin{array}{c} 20 \\ 21 \end{array}public class CaseStudy4 {
\frac{22}{23}public static void main (String [] args) throws IOException, TreeException, <br>
24 \over 25 \over 7/7000 \ \text{Auto}-generated \ method \ study}24 mde.gui.gen.TreeException {<br>25 // TODO Auto−generated method stub
\frac{26}{27}\begin{array}{c|c|c|c|c} 27 & // Constructor the Information Model. \end{array} Diagram information.
28 \text{ } mde. information . model . Diagram informationModel = new<br>29 mde. information . model . Diagram ();
\frac{29}{30}<br>\frac{30}{31}31 Entity p e r s o n E n t i t y = new Entity ( " P erson " ) ;
32 Attribute a t t r 2 = new Attribute ( "name" , new Type( " S t r i n g " ) ) ;
33 Attribute a t t r 3 = new Attribute ( " age " , new Type( " S t r i n g " ) )
34 . setLowerbound ( 1 8 ) . setUpperbound ( 3 5 ) ;
35 Attribute a t t r 4 = new Attribute ( " g end e r " , new Type( " S t r i n g " ) ) ;
36 personEntity.addAttribute(attr2);<br>37 personEntity.addAttribute(attr3);
38 personEntity.addAttribute(attr4);<br>39 Entity studentEntity = new Entity("Student");
```


```
134 // Generate Data Model<br>135 InfDiagramToDDiagram
135 InfDiagramToDDiagram topRule4 = new InfDiagramToDDiagram ();<br>136 InfDiagramToDDiagram dataModel = topRule4 translate (informationModel);
                          \textbf{DDiagram} \quad \text{dataModel} = \text{topRule4} \cdot \text{transfer}(\overline{\text{informationModel}});137<br>138
138 <br>
ASTWriter writer4 = new ASTWriter (new File ("Uni_DataModel.xml"));<br>
writer4.usePackage ("mde.data.model", "xmlns:data");<br>
writer4.usePackage ("mde.model", "xmlns:mod");<br>
writer4.writeDocument (dataModel);<br>
unite
143<br>144
144 System.out.println("(5) Data Model is Created (Information"+<br>145 "Model −−> Data Model).");<br>// −−−−−−−−−−−−−−−−−−−−−− //
147
148 // g e n e r a t e D a t a b a s e Schema Mode l ( D e t a i l e d DFD"+
149 " + Data Model −−−> Schema )
150 DDiagramToSchemaOnly topRule5 = new DDiagramToSchemaOnly();<br>151 Schema schemaModel = topRule5.translate(null, dataModel);
151<br>152<br>153System.out.println(schemaModel.getName());
154
155 ASTWriter w r i t e r 5 = new ASTWriter ( new F i l e ( "Uni SchemaModel . xml" ) ) ;
156 w r i t e r 5 . us ePa ckag e ( "mde . dbs . model " , " xmlns : dbs " ) ;
157 w r i t e r 5 . us ePa ckag e ( "mde . model " , " xmlns : mod" ) ;
158 w r i t e r 5 . wr iteDocument ( schemaModel ) ;
                          write r 5 . close ();
\frac{159}{160}161 System .out .println("(9) Database Schema Model is Created"+<br>162 " (DataFlow + Data Model −−> Database Schema Model).");
\frac{163}{164}164 // cod e g e n e r a t i o n
165 DumpF i leGenerator MySQLGenerator = new
166 MySQLDumpFileGenerator (schemaModel);<br>166 MySQLDumpFileGenerator (schemaModel);<br>167
167 System out println("calling generate in mysql"+<br>168 Package generator");
169 MySQLGenerator . g e n e r a t e ( ) ;
170 System . out . p r i n t l n ( " F i n i s h e d g e n e r a t i n g MySQL Schema OK" ) ;
171 }
172 }
```
Information Model Representation (3a)


```
\begin{array}{c|c|c|c|c} 1 < 2 < \text{cml version="1.0" encoding="UTF-8"?} \\ \hline 2 < \text{inf.Diagram } \text{x} < \text{mls:inf="mde. information} \end{array}2 \langle inf:Diagram x mlns:inf="mde.information.model" id="0"><br>3 \langle inf:Entity id="1" name="Student">
   3 <i n f :E n t i t y i d="1" name=" Stud ent ">
4 <in f :A t t r ibu te i d="2" name=" regNumber " i d e n t i f i e r=" t r u e " s i z e=" 0">
5 <Type i d="3" name=" I n t e g e r " />
  4 \langle inf:Attribute i<br>5 \langle Type id="3" n<br>6 \langle inf:Attribute><br>7 \langle inf:Attribute i
   7 <inf:Attribute id="4" name="username" size="0"><br>8 <Type id="5" name="String"/><br>9 </inf:Attribute>
 10 <inf:Attribute id="6" name="password" size="0"><br>
11 <Type id="7" name="String"/><br>
</inf:Attribute>
\begin{array}{c|c}\n 13 & & \langle / \text{inf:Entity} \rangle \\
 14 & & \langle \text{inf:Entity} \rangle\n \end{array}14 <inf:Entity id="8" name="Staff"><br>15 </inf:Attribute id="9" name="empNumber" identifier="true" size="0"><br>16 </rype id="10" name="Integer"/>
 17 </inf:Attribute><br>
18 <inf:Attribute id="11" name="salary" size="0"><br>
19 <Type id="12" name="Double"/><br>
<//> </reduction="Double"/>
 \begin{array}{ll} 21 & </\texttt{inf:Entity>} \ 22 & <\texttt{inf:Entity} \texttt{ id} = "13" \texttt{ name} = "Person" > \end{array}23 <inf:Attribute id="14" name="name" size="0"><br>24 <Type id="15" name="String"/><br>25 </inf:Attribute>
 26 <in f :A t t r ibu te i d=" 16 " name=" age " upperbound=" 35 " lowerbound=" 18 " s i z e="0">
27 <Type i d=" 17 " name=" S t r i n g "/>
28 </ in f :A t t r ibu te>
 29 <inf:Attribute id="18" name="gender" size="0"><br>30 <Type id="19" name="String"/>
\begin{array}{c|c}\n 31 \\
 32 \\
 \end{array} \begin{array}{c}\n \begin{array}{c}\n \begin{array}{c}\n \begin{array}{c}\n \begin{array}{c}\n \end{array} \\
 \begin{array}{c}\n\begin{array}{c|c} 32 \overline{33} & \overline{51} & \overline{11} & \overline{11} & \overline{11} \\ 33 & \overline{11} & \overline{11} & \overline{11} \\ 34 & \overline{11} & \overline{11} & \overline{11} \\ \end{array}33 <inf:Entity id="20" name="Module"><br>
34 <inf:Attribute id="21" name="code" identifier="true" size="0"><br>
35 </Iype id="22" name="Integer"/><br>
</inf:Attribute>
 37 <inf:Attribute id="23" name="title" size="0"><br>38 <Type id="24" name="String"/><br></inf:Attribute>
 40 <inf:Attribute id="25" name="credit" size="0"><br>41 <\langleType id="26" name="Integer"/><br>42 </inf:Attribute>
 43 <inf:Attribute id="27" name="desc" size="0"><br>44 <Type id="28" name="String"/>
```


```
139 </inf:Role><br>140 </inf:Association><br>141 </inf:Diagram>
```
Data Dependency Model Representation (3a)

Listing C.3: Uni DataModel.xml

Database and Query (DBQ) Model Representation (3a)

Listing C.4: Uni SchemaModel.xml

```
1 <?xml vers ion=" 1 . 0 " en c od in g="UTF−8" ?>
2 <dbs:Schema xm ln s :db s="mde . dbs . model " i d="0" name=" d a t ab a s e ">
3 <dbs:Table i d="1" name=" P e r s on S tud en t ">
4 <dbs:Column i d="2" name=" i d e n t i t y " au t om a t i c=" t r u e " s i z e=" 10 ">
5 <mod:Type xmlns:mod="mde . model " i d="3" name="INTEGER"/>
6 </dbs:Column>
7 <dbs:Column i d="4" name="name" s i z e=" 30 ">
8 <mod:Type i d="5" name="VARCHAR"/>
9 </dbs:Column>
10 <dbs:Column i d="6" name=" age " upperbound=" 35 " lowerbound=" 18 " s i z e=" 30 ">
11 <mod:Type i d="7" name="VARCHAR"/>
12 </dbs:Column>
13 <dbs:Column i d="8" name=" g end e r " s i z e=" 30 ">
14 <mod:Type i d="9" name="VARCHAR"/>
15 </dbs:Column>
16 <dbs:Column i d=" 10 " name=" regNumber " s i z e=" 10 ">
```
 ${17} \begin{array}{rcl} & & <\hspace{-10.1cm}\textbf{mod:Type} \text{ id} = "11" \text{ name="INTEGR" } / \hspace{-10.1cm} > \\ & & <\hspace{-10.1cm}\textbf{mod:Column}\hspace{-10.1cm} > \end{array}$ 18 </dbs:Column>

19 <dbs:Column id="12" name="username" size="30">

20 <mod:Type id="13" name="VARCHAR"/>

21 </dbs:Column> 22 <dbs:Column id="14" name="password" size="30">

<mod:Type id="15" name="VARCHAR"/> 24 </dbs:Column>
25 <dbs:PrimaryKey id="16"> 25

26

26 <dbs:Column r e f="2"/>

27

27 </dbs:PrimaryKey> $\begin{array}{ll} \vspace{2.0ex} \begin{array}{ll} & \vspace{2.0ex} \vspace{2.0$ $\begin{array}{c|c} 34 & <\textbf{mod:Type} \text{ref} = "5" / > \ \hline & <\textbf{dbs:Column}> \end{array}$ 36 <**dbs:Column** id="21" name="age" upperbound="35" lowerbound="18" size="30">
37 <m**od:Type** ref="7"/>
38 </**dbs:Column**> 39 <dbs:Column id="22" name="gender" size="30">
40 <mod:Type ref="9"/>
41 </dbs:Column> 42 <dbs:Column i d=" 23 " name="empNumber" s i z e=" 10 "> 43 <mod:Type i d=" 24 " name="INTEGER"/> 44 </dbs:Column> 45 <dbs:Column i d=" 25 " name=" s a l a r y " s i z e="5 "> 46 <mod:Type i d=" 26 " name="DOUBLE" /> 47 </dbs:Column> 48 <dbs:Column i d=" 27 " name=" Addr ess . i d e n t i t y " s i z e=" 10 "> 49 <mod:Type i d=" 28 " name="INTEGER"/> 50 </dbs:Column> 51
 \langle dbs:ForeignKey id="29">
 \langle dbs:Column ref="27"/>
 \langle dbs:Table id="30" nan 52 <dbs:Column ref="27"/>

<dbs:Table id="30" name="Address">

<dbs:Table id="30" name="identity" size="10">

55 <mod:Type ref="28"/> 56 </dbs:Column>
57 </dbs:Column io 57 <dbs:Column id="32" name="postcode" size="30">

<mod:Type id="33" name="VARCHAR"/> 59 </dbs:Column> 60 <dbs:Column i d=" 34 " name=" s t r e e t " s i z e=" 30 "> 61 <mod:Type i d=" 35 " name="VARCHAR"/> $\begin{array}{c|c}\n 62 & & \leq d$ bs:Column>

63 <dbs:Column id 63 <dbs:Column i d=" 36 " name=" c i t y " s i z e=" 30 "> 64 <mod:Type i d=" 37 " name="VARCHAR"/> 65 \angle /dbs:Column> 66 <dbs:PrimaryKey id="38">

67 <dbs:Column ref="31"/>

68 </dbs:PrimaryKey> 69 </dbs:Table> 70 </dbs:ForeignKey> 71 <dbs:PrimaryKey id="39">

72 <dbs:Column ref="18"/> 73 </dbs:PrimaryKey> 74 </dbs:Table> 75 <dbs:Table i d=" 40 " name="Assessment Exam "> 76 <dbs:Column i d=" 41 " name=" i d e n t i t y " au t om a t i c=" t r u e " s i z e=" 10 "> 77 <mod:Type i d=" 42 " name="INTEGER"/> 78 </dbs:Column> 79 <dbs:Column i d=" 43 " name=" i d " s i z e=" 30 "> 80 <mod:Type i d=" 44 " name="VARCHAR"/> 81 </dbs:Column> 82 <dbs:Column i d=" 45 " name=" t i t l e " s i z e=" 30 "> 83 <mod:Type i d=" 46 " name="VARCHAR"/> 84 </dbs:Column> 85 <dbs:Column id="47" name="group" size="5">
86 <mod:Type id="48" name="BOOLEAN"/>
87 </dbs:Column> 88 <dbs:Column id="49" name="date" size="5">
89 <mod:Type id="50" name="DATE"/> 90 </dbs:Column>

91 cdbs:Column id="51" name="Module .code" size="10">

92 <mod:Type id="52" name="INTEGER"/>

33 </dbs:Column> 94
 \langle dbs:ForeignKey id="53">
 \langle dbs:Column ref="51"/> 95 <dbs:Column ref="51"/>

96 <dbs:Table id="54" name="Module">

97 <dbs:Table id="54" name="code" size="10">

98 <mod:Type ref="52"/> 99 </dbs:Column>

100 </dbs:Column io

101 </additions\> 100 <dbs:Column i d=" 56 " name=" t i t l e " s i z e=" 30 "> 101 <mod:Type i d=" 57 " name="VARCHAR"/> 102 </**dbs:Column>**
103 <**dbs:Column** id="58" name="credit" size="10">
104 <**mod:Type** id="59" name="NTEGER"/> $\begin{array}{c|c}\n 105 \\
 106\n \end{array}$ $\begin{array}{c}\n 3.16 \\
 4 \text{bs:Column} \\
 5.16\n \end{array}$ 106 <dbs:Column i d=" 60 " name=" d e s c " s i z e=" 30 "> 107 <mod:Type i d=" 61 " name="VARCHAR"/> 108 </dbs:Column> 109 <dbs:PrimaryKey i d=" 62 "> 110 <dbs:Column r e f=" 55 "/>

Executable MySQL Code (3a)

Listing C.5: database MySQL.sql

 $--\ Data base\ Create$ ${\bf C}\textbf{REATE} \text{ DATABASE } \text{sysDatabase} \; ;$ USE sysDatabase; −− Structure for table 'Person_Student' 8 **CREATE TABLE** Person_Student (10) MOT NULL,

10 identity INT (10) NOT NULL,

10 name VARCHAR(30),

2 gender VARCHAR(30),

2 gender VARCHAR(30),

13 regNumber INT (10),

15 password VARCHAR(30),

15 password VARCHAR(30), $\begin{bmatrix} 13 \\ 14 \\ 15 \\ 16 \\ 17 \\ 18 \end{bmatrix}$ $- -$ Structure for table 'Person_Staff'

```
\frac{19}{20}\begin{array}{c} 13 \\ 20 \\ 21 \end{array} CREATE TABLE Person-Staff<br>21 identity INT(10) NOT NU.
  21 identity INT(10) NOT NULL,<br>
22 name VARCHAR(30),<br>
23 age VARCHAR(30),
 \begin{bmatrix} 24 \\ 25 \end{bmatrix} gender VARCHAR(30)<br>25 empNumber INT(10).
  25 empNumber INT(10),<br>26 salary DOUBLE(5),<br>27 Address.identity INT(10) NOT NULL,
  28 |    PRIMARY KEY(identity),<br>29 |    FOREIGN KEY(Address.identity) REFERENCES Address(identity));
 \frac{30}{31}31 −− S t r u c t u r e f o r t a b l e ' Assessmen t Exam '
 32
 \begin{array}{c} 33 \\ 33 \\ 34 \end{array} CREATE TABLE Assessment Exam (
  \begin{array}{c|c} 34 & \textbf{identity INT}(10) & \textbf{NOT NULL}, \ \hline 35 & \textbf{id VARCHAR}(30) \,, \end{array}36 title VARCHAR(30),<br>37 group BOOLEAN(5),
  38 date DATE(5),<br>39 Module.code INT(10),<br>40 PRIMARY KEY(identity),<br>41 FOREIGN KEY(Module.code) REFERENCES Module(code));
 \frac{42}{43}43 −− S t r u c t u r e f o r t a b l e ' A s s e s sm e n t P r o j e c t '
 \begin{array}{c} 44 \\ 45 \end{array}\begin{array}{c} 45 \\ 46 \end{array} CREATE TABLE Assessment_Project (
  46 id VARCHAR(30) NOT NULL,<br>47 title VARCHAR(30),
  48 group BOOLEAN(5),<br>
50 value INT(10),<br>
deadline DATE(5),
  51 Module.code INT(10),<br>52 PRIMARY KEY(id),
 53 FOREIGN KEY(Module . code) REFERENCES Module (code));
 \begin{array}{c} 54 \\ 55 \end{array}- Structure for table 'Enrollment '
 \frac{56}{57}57 CREATE TABLE Enrollment (
identity INT(10) NOT NULL,<br>59 code INT(10),<br>
ende INT(10),<br>
regNumber INT(10),
  61 Module.code INT(10),<br>
Person_Student.identity INT(10) NOT NULL,<br>
PRIMARY KEY(identity),<br>
FOREIGN KEY(Module.code) REFERENCES Module(code),<br>
FOREIGN KEY(Person_Student.identity) REFERENCES Person_Student(identity));
 66
          67 −− S t r u c t u r e f o r t a b l e ' Modu le '
 \frac{68}{69}69 CREATE TABLE Module (<br>70 code INT(10) NOT NULL,
  \begin{array}{c|c} 71 & \text{title } \textbf{VARCHAR}(30)\,, \ \hline 72 & \text{credit } \textbf{INT}(10)\,, \end{array}73 desc VARCHAR(30)<br>74 PRIMARY KEY(code
             PRIMARY KEY(code) ;
 rac{75}{76}76 −− S t r u c t u r e f o r t a b l e ' Add r ess '
 rac{77}{78}78 CREATE TABLE Address (<br>
79 identity INT(10) NOT NULL,<br>
80 postcode VARCHAR(30),
 \begin{array}{c|c|c|c} \hline \text{81} & \text{street VARCHAR}(30), \ \hline \text{61} & \text{city VARCHAR}(30), \ \hline \end{array}82 city VARCHAR(30)<br>83 PRIMARY KEY(iden
             PRIMARY KEY(identity):
 \frac{84}{85}-- Trigger: Applying Checking Constraints on table 'Person_Student'
 rac{86}{87}87 DELIMITER //
 88<br>89
  89 | DROP TRIGGER IF EXISTS 'person_studentCheck' //<br>90 | CREATE TRIGGER 'person_studentCheck' BEFORE INSERT ON Person_Student
 91 FOR EACH ROW<br>92 IF (NEW. ago
  92 IF (NEW. age < 18 OR NEW. age > 35) THEN<br>
93 SET msg = 'INVALID DATA IN age ';<br>
94 SIGNAL SQLSTATE '45000 ' SET MESSAGE_TEXT = msg
 \begin{array}{c|c}\n 95 & \text{END IF} ;\n\end{array}/97<br>98
          -− Trigger: Applying Checking Constraints on table 'Person_Staff'
\frac{99}{100}DELIMITER //
\begin{array}{l} 101 \\ 102 \\ 103 \end{array}102 | DROP TRIGGER IF EXISTS 'person_staffCheck' //<br>103 | CREATE TRIGGER 'person_staffCheck' BEFORE INSERT ON Person_Staff
104 FOR EACH ROW<br>105 IF (NEW. ago
105 IF (NEW. age < 18 OR NEW. age > 35) THEN<br>106 SET msg = 'INVALID DATA IN age';<br>107 SIGNAL SQLSTATE '45000 ' SET MESSAGE_TEXT = msg;
\begin{bmatrix} 108 \\ 108 \\ 109 \end{bmatrix} // END IF;
          /
```
Experiment (3b): Impact Model Consturction

Listing C.6: CaseStudy12.java

```
1 package mde example:
  \frac{2}{3}\begin{array}{c|c} 3 & \text{import java.io.File}; \\ 4 & \text{import java.io.IOEx} \end{array}4 import java.io.IOException;<br>5 import mde.data.model.DDiagram;
   6 import mde . database .gen . DumpFileGenerator ;<br>8 import mde . database .gen . TreeException ;<br>9 import mde . dm2schem . rule . DDiagramToSchemaOnly ;<br>9 import mde . dm2schem . rule . DDiagramToSchemaOnly ;
10 import mde.impact.model.ImpBoundary;<br>11 import mde.impact.model.ImpConjunction;
12 import mde.impact.model.ImpCreateFlow;<br>13 import mde.impact.model.ImpDiagram;
14 import mde . impact . model . ImpObject ;
15 import mde . impact . model . ImpReadFlow;
16 import mde .impact .model .ImpRole ;<br>17 import mde .impact .model .ImpTask ;<br>18 import mde .impact .model .ImpUpdateFlow ;<br>19 import mde .impact2information .ImpDiagramToDiagram;
20 import mde.inf2dm.rule.InfDiagramToDDiagram;<br>21 import mde.information.model.Diagram;
22 import mde.mysql.gen.MySQLDumpFileGenerator;<br>23 import org.jast.ast.ASTWriter;
\begin{array}{c} 24 \\ 25 \end{array}25 public class Full_example12 {<br>
public static void main(String [] args) throws IOException, TreeException,<br>
27 mde.gui.gen.TreeException {<br>
//TODO Auto-generated method stub
29
\begin{array}{ll}30 & // \textit{Construct the Impact Model} \\\textbf{31} & \textbf{ImpDiagram Inpatch} \\\textbf{32} & \textbf{ImpDiagram Inpatch} \\\textbf{33} & \textbf{ImpBoundary impboundary2 = new ImpBoundary("Enrol_Boundary");} \\\textbf{34} & \textbf{ImpBoundary impboundary2 = new ImpBoundary("SetAddress_Boundary");} \\\textbf{35} & \textbf{ImpBoundary impboundary3 = new ImpBoundary("ManageAssessment_Boundary");} \end{array}35 \left\{\begin{array}{lll} \text{ImpTask impTask1 = new ImpTask("Enrol");} \text{ImpTask impTask impTask2 = new ImpTask("Set Address");} \text{ImpTask impTask impTask("Set Answers");} \text{ImpTask impTask impTask("Set Answers");} \end{array}\right.39<br>
ImpConjunction conj = new ImpConjunction ("Set Address");<br>
ImpObject impObje = new ImpObject ("Student");<br>
142<br>
ImpObject impObj2 = new ImpObject ("Staff");<br>
ImpObject impObj3 = new ImpObject ("Enrollment");<br>
143<br>
Imp
\frac{46}{47}\begin{tabular}{ll} ${{\cal A}}$ & ${{\cal B}}$ & ${{\cal68<br>69<br>70<br>71\text{impboundary} . addImpTask ( \text{impTask1} )
71<br>
71 impboundary . addImpObject (impObj1);<br>
72 impboundary . addImpObject (impObj5);<br>
73 impboundary . addImpObject (impObj3);
72 impboundary . addImpObject ( impObj5 ) ;
73 impboundary . addImpObject ( impObj3 ) ;
74 impboundary . addImpReadFlow ( r f 1 ) ;
75 impboundary . addImpReadFlow ( r f 2 ) ;
76 impboundary . addImpCreateF low ( c f ) ;
rac{76}{77}// set address boundary contents\frac{78}{79}80<br>
ImpReadFlow rf11 = new ImpReadFlow();<br>
82<br>
RDROLe imprf11 = new ImpReadFlow("staddress.conj", conj);<br>
ImpReadFlow imprf21 = new ImpReadFlow("student", impObj1).setMultiple(false);<br>
83<br>
rf11.addImpRele (imprf21);<br>
86<br>

88 r f 21 . addImpRole (imprf41);<br>89 r f 21 . addImpRole (imprf31);
```

```
90 <br>
ImpReadFlow rf31 = new ImpReadFlow ();<br>
92 <br>
ImpRole imprf81 = new ImpRole ("set_address_conj", conj);<br>
193 <br>
ImpRole imprf71 = new ImpRole("set_address", impTask2);<br>
193 <br>
193 <br>
193 <br>
ImpRole imprf81 = new ImpCreate
 99 cf1.addImpRole(imprf61);
100<br>101
101 impboundary2.addImpTask(impTask2);<br>102 impboundary2.addImpObject(impObj1)
102 impboundary2 . addImpObject ( impObj1 ) ;
103 impboundary2 . addImpObject ( impObj2 ) ;
104 impboundary2 . addImpObject ( impObj4 ) ;
105 impboundary2.addImpConjunction (conj);<br>106 impboundary2.addImpReadFlow(rf11);<br>107 impboundary2.addImpReadFlow(rf21);<br>impboundary2.addImpReadFlow(rf31);
109 impboundary2.addImpCreateFlow(cf1);
\frac{110}{111}11 set assessment boundary contents
\frac{112}{113}113<br>
114 ImpRole impr<br/>f22 = newImpRole("set_assess", impTask3);<br>
115<br>
115 ImpRole impr<br/>f22 = newImpRole("module", impObj5);<br>
116 r f 22.addImpRole(impr<br/>f42);<br>
117 r f 22.addImpRole(imprf42);<br>
117 r f 22.add
133
134 impboundary3.addImpTask(impTask3);<br>135 impboundary3.addImpTask(impTask4);
135 impboundary3 . addImpTask ( impTask4 ) ;
136 impboundary3 . addImpObject ( impObj5 ) ;
137 impboundary3 . addImpObject ( impObj6 ) ;
138 impboundary3.addImpReadFlow ( rf22 );<br>139 impboundary3.addImpReadFlow (rf32 );<br>140 impboundary3.addImpCreateFlow ( cf11 );<br>impboundary3.addImpUpdateFlow ( uf );
142
143 ImpactModel.addImpBoundary (impboundary)<br>144 ImpactModel addImpBoundary (imphoundary)
144 ImpactModel . addImpBoundary ( impboundary2 ) ;
145 ImpactModel . addImpBoundary ( impboundary3 ) ;
146<br>147
147 ASTWriter w r i t e r 1 = new ASTWriter(new F i l e ( " impactMode l . xml" ) ) ;
148 w r i t e r 1 . us ePa ckag e ( "mde . impact . model " , " xmlns : imp" ) ;
149 w r i t e r 1 . wr iteDocument ( ImpactModel ) ;
\begin{array}{c} 1150 \\ 150 \end{array} writer 1 . close ();
\frac{151}{152}System.out.println ("(2) Impact Model is Created by user.");
\frac{153}{153}154 // Generate Information Model<br>155 ImpDiagramToDiagram topRule =
155 | ImpDiagramToDiagram topRule = new ImpDiagramToDiagram ();<br>156 | Diagram initial_dataModel = topRule.translate(ImpactModel);
\frac{157}{158}158 ASTWriter w r i t e r 2 = new ASTWriter(new F i l e ( " i n i t i a l I n f o M o d e l . xml" ) ) ;
159 w r i t e r 2 . us ePa ckag e ( "mde . i n f o rm a t i o n . model " , " xmlns : i n f " ) ;
160 w r i t e r 2 . wr iteDocument ( i n i t i a l d a t aM o d e l ) ;
\begin{array}{c|c} \n & \text{inter 2. close ()};\n\end{array}\frac{162}{163}System.out.println(" (2) Initial Information Model is Create by rules.");164 | // −−−−−−−−−−−−−−−−−−−−−−−−−− //
165<br>166
\begin{array}{c|c} 166 & // Genome Data Model \\ 167 & \textbf{InfDiagramToDDiagram} \end{array}167 167 InfDiagramToDDiagram topRule4 = new InfDiagramToDDiagram ()<br>168 DDiagram dataModel = topRule4.translate(initial_dataModel)
                           \text{DDiagram} dataModel = topRule4.translate (initial_dataModel);
\frac{169}{170}170 ASTWriter writer4 = new ASTWriter (new File ("Uni_DataModel.xml"));<br>171 writer4.usePackage ("mde.data.model", "xmlns:data");<br>172 writer4.usePackage ("mde.model", "xmlns:mod");<br>173 writer4.writeDocument (dataModel
175
176 System . out . p r i n t l n ( " ( 5 ) Data Model i s Cr eat ed " ( I n f o rm a t i o n Model −−> Data Model ) . " ) ;
177 // −−−−−−−−−−−−−−−−−−−− //
178
179 – //generate Database Schema Model (Detailed_DFD+Data Model ----> Schema)<br>180 – DDiagramToSchemaOnly topRule5 = new DDiagramToSchemaOnly();<br>181 – Schema schemaModel = topRule5.translate(null, dataModel);
\frac{182}{183}System . out . println (schemaModel . getName ());
```


Impact Model Representation (3b)

Generated Information Model Representation (3b)


```
\begin{array}{c|c}\n 39 \\
 40 \\
 \end{array} \begin{array}{c}\n \times \inf : \textbf{Entity} \text{ ref} = "2" / > \n \\ \n  & \times / \inf : \textbf{Role}\n \end{array}\begin{tabular}{ll} $40$ & $\sless\mbox{\small{in}}$ \emph{if (Association)}$ \\ $42$ & $\leq\mbox{\small{in}}$ \emph{f (Association id="21'')$} \\ $43$ & $\leq\mbox{\small{in}}$ \emph{f (Rose left)}$ \\ $44$ & $\leq\mbox{\small{in}}$ \emph{f (Role left)}$ \\ $45$ & $\leq\mbox{\small{in}}$ \emph{f (Role right) right} $p$ \\ $46$ & $\leq\mbox{\small{in}}$ \emph{f (Role)>} \\ $46$ & $\leq\mbox{\small{in}}$ \emph{f (Role)>}\begin{array}{c|c} 59 & </ \inf : \mathbf{Role} > \ 60 & <\inf : \mathbf{Role} \text{ refer} = "25" / > \ 61 & </ \inf : \mathbf{Association} > \end{array}62 \mid \langle / \inf : \mathbf{Diagram} \rangle
```
Generated Data Model Representation (3b)

60 <data:DRole ref="25" />

61 </data:DDependency>

62 </data:DDiagram>

Generated Database and Query Model Representation (3b)

Listing C.10: SchemaModel.xml


```
81 <dbs:Table id="36" name="Assessment"><br>
<dbs:Column id="37" name="identity" automatic="true" size="10"><br>
<dbs:Column id="33"/><br>
<dbs:Column>>
  85 <dbs:PrimaryKey id="38"><br>86 <dbs:Column ref="37"/><br>87 </dbs:PrimaryKey>
  88 </dbs:Table><br>89 </dbs:ForeignKey>
  90 <dbs:ForeignKey id="39"><br>
91 <dbs:Column ref="32"/><br>
92 <dbs:Table ref="36"/>
  93 </dbs:ForeignKey><br>
94 <dbs:PrimaryKey id="40"><br>
95 <dbs:Column ref="30"/>
 \begin{array}{c|c}\n 96 \\
 97 \\
 \hline\n 6/ \text{dbs:PrimaryKey}\n\end{array}97 </dbs:Table >ef="16"/><br>
98 <dbs:Table ref="16"/><br>
99 <dbs:Table ref="21"/><br>
dbs:Table ref="25"/><br>
(dbs:Table ref="36"/>
102 \leq \angledbs:Schema>
```
Executable MySQL Code (3b)


```
1 \vert -\vert Database Creation
  \begin{array}{c} 2 \\ 3 \\ 4 \end{array}CREATE DATABASE sysDatabase;
           USE sysDatabase;
  \begin{array}{c} 5 \\ 6 \end{array}6 −− S t r u c t u r e f o r t a b l e ' G en e ra l '
   \frac{1}{8}\begin{array}{c|c} \n\hline\n8 & \text{CREATE TABLE General} \\
9 & \text{identity INT}(10) & \text{N} \\
\end{array}9 identity INT(10) NOT NULL,<br>
11 Student.identity INT(10) NOT NULL,<br>
11 Staff.identity INT(10) NOT NULL,<br>
21 Address.identity INT(10) NOT NULL,<br>
22 Address.identity INT(10) NOT NULL,<br>
23 FOREIGN KEY(Student.identity) REFER
\begin{array}{c} 17 \\ 18 \end{array}- Structure for table 'Student'
\frac{19}{20}20 CREATE TABLE Student (<br>
11 identity INT(10) NOT NULL,<br>
Enrollment.identity INT(10) NOT NULL,<br>
23 PRIMARY KEY(identity), FOREIGN KEY(Enrollment.identity), REFERENCES Enrollment(identity));
\begin{array}{c} 23 \\ 24 \\ 25 \\ 26 \end{array}26 −− S t r u c t u r e f o r t a b l e ' Modu le '
\frac{27}{28}28 CREATE TABLE Module (<br>29 identity INT(10) NOT NULL,<br>30 Enrollment .identity INT(10) NOT NULL,<br>31 Assessment .identity INT(10) NOT NULL,
 32 FRIMARY KEY(identity),<br>
33 FOREIGN KEY(Enrollment.identity) REFERENCES Enrollment(identity),<br>
34 FOREIGN KEY(Assessment.identity) REFERENCES Assessment(identity),<br>
55 FOREIGN KEY(Assessment.identity) REFERENCES Assessme
\frac{36}{37}- Structure for table 'Enrollment
38
\begin{array}{c} 39 \ 39 \ 40 \ \end{array} CREATE TABLE Enrollment
 \begin{array}{ll} \text{40} & \text{identity INT}(10) \text{ NOT NULL}, \\ \text{PRIMARY KEY}(\text{identity})) \, ; \end{array}\begin{array}{c} 42 \\ 43 \end{array}- Structure for table 'Staff'
\frac{44}{45}\begin{array}{c|c} 45 & \textbf{CREATE TABLE } \text{Staff} \\ 46 & \textbf{identity } \textbf{INT}(10) \end{array}\begin{array}{lll} 46 & & \text{identity} & \text{INT (10)} & \text{NOT} & \text{NULL}, \end{array}\frac{48}{49}- Structure for table 'Address'
\begin{array}{c} 50 \\ 51 \end{array}\begin{array}{c} 50 \\ 51 \\ 52 \end{array} CREATE TABLE Address<br>52 identity INT(10) NC
 \begin{array}{c|c} 52 & i \text{dentity INT}(10) & \text{NOT NULL}, \end{array} PRIMARY KEY(identity));
\begin{array}{c} 54 \\ 55 \end{array}- Structure for table 'Assessment'
\frac{56}{57}57 CREATE TABLE Assessment (<br>58 identity INT(10) NOT NULL,<br>59 PRIMARY KEY(identity));
```
D End-User Evaluation Experiment of ML Notation

D.1 End-User Evaluation Experiment

D.1.1 Description of the Online System

This section provides a description of business activity structure of an Online Hospital Booking System (OHBS). An admin can manage doctors in the system by either: adding a new doctor, updating a profile of an existing doctor, removing an existing doctor from the system. The admin is able to add a new doctor by first adding the personal detail of a new doctor to the system. Then, adding the login detail of the new doctor to the system.

- Draw an ellipse for each business activity (process) with a suitable name..
- Draw a white diamond and branches to represent options part processes (sub-processes) of a whole business process.
- Draw a white diamond and branches to represent all parts processes (sub-processes) of a whole business process.

Time now (before drawing):

Time now (After drawing):

D.1.2 Description of Business Items (Entities)

This part describes information (business data) that is used within the online system to complete business activities. The system needs information about each doctor to be stored in the system by Adding personal detail of the new doctor process. There is a login information that has to be stored for each new doctor in the system by adding login detail process.

In addition, the system modifies the information about doctor profile by Update doctor profile process, and removes an existing doctor from the system by Remove doctor process.

- Draw a rectangle for each business entity and write a suitable name for it.
- Draw a line between each two entities that have a kind of relationship.

Time now (before drawing):

Time now (After drawing):

D.1.3 Description of How the System Data interact with system

(Samilar text as the previous section) This part describes the interaction between business information (data) and business process within the online system. The system needs information about each doctor to be stored in the system by Adding personal detail of the new doctor process. There is a login detail that has to be stored for each new doctor in the system by adding login detail process.

In addition, the system modifies doctor profile by Update doctor profile process, and remove an existing doctor from the system by Removes doctor process.

- Draw an ellipse for each business activity (process) with a suitable name.
- Draw a rectangle for each business entity with a suitable name.
- Draw a suitable arrow between tasks and related entities (as explained).

Time now (before drawing):

Time now (After drawing):